

FINAL REPORT
Grant Number NAG-1-1378

**EFFECT OF DESIGN SELECTION ON
RESPONSE SURFACE PERFORMANCE**

from

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Table of Contents

1.	Introduction	1
	<u>1.1 Quality of Fit</u>	2
	<u>1.1.1 Fit at the designs</u>	2
	<u>1.1.2 Overall fit</u>	3
	<u>1.2. Polynomial Approximations</u>	4
	<u>1.2.1 Exactly-determined approximation</u>	5
	<u>1.2.2 Over-determined approximation</u>	5
	<u>1.2.3 Under-determined approximation</u>	6
	<u>1.3 Artificial Neural Nets</u>	7
2.	Levels of Designs	10
	<u>2.1 Taylor Series Approximation</u>	10
	<u>2.2 Example</u>	13
	<u>2.3 Conclusion</u>	18
3.	Standard Designs	19
	<u>3.1 Underlying Principle</u>	19
	<u>3.2 Statistical Concepts</u>	20
	<u>3.3 Orthogonal Designs</u>	23
	<u>3.3.1 Scaling</u>	24
	3.3.1.1 Example of Scaled Designs:	25

3.3.2 Bias	26
3.3.2.1 A bias example--linear approximating polynomial but the exact function contains linear terms and cross- product terms:	27
3.3.2.2 A bias example--linear approximating function but the exact function is a complete quadratic polynomial:	28
3.3.3 <u>Orthogonal Designs for Linear Approximations</u>	29
3.3.4 <u>Orthogonal Designs for 2nd Order Polynomial Approximations</u>	30
3.3.5 <u>General Discussion of Orthogonal Designs</u>	31
<u>3.4 Central Composite Designs--Designs for Fitting Second Order Models</u> ..	31
<u>3.4.1 Format of the central composite design</u>	31
3.4.1.1 Single center point rotatable second order experimental designs:	32
3.4.1.2 Multiple center point rotatable uniform precision designs:	33
3.4.1.3 Single center point orthogonal central composite designs:	34
3.4.1.4 Rotatable orthogonal designs:	35
<u>3.4.2 Discussion of the central composite design</u>	36
<u>3.4.3 Example -- Fox's Banana Function</u>	37
<u>3.4.4 Conclusion</u>	44

4.	Optimality Criteria	46
	<u>4.1 D, A, E, G, and V Optimality Criteria</u>	46
	<u>4.1.1 Criteria Applied to a One Dimensional Example</u>	47
	<u>4.2 S and Q Optimality Criteria</u>	49
	<u>4.2.2 Criteria Applied to a One Dimensional Example</u>	51
	<u>4.3 An Alternate Approach--Random Selection of Designs</u>	52
	<u>4.3.1 Random Selection of Designs Applied to a One Dimensional Example</u>	52
	<u>4.4 Larger Problems</u>	53
	<u>4.5 Optimality Criteria Based on Minimizing Uncertainty</u>	55
	<u>4.6 Conclusion</u>	55
5.	Significance Testing of Coefficients	57
	<u>5.1 Introduction</u>	57
	<u>5.2 t-test</u>	57
	<u>5.3 Example 1 -- Fox's Banana Function</u>	59
	<u>5.4 Example 2</u>	59
	<u>5.5 Conclusion</u>	60
6.	Applicability of the Response Surface Technique	61
	<u>6.1 Introduction</u>	61
	<u>6.2 Conclusion</u>	66

7.	Additional Examples	67
	<u>7.1 Introduction</u>	67
	<u>7.2 The 35 Bar Truss with 4 Design Variables</u>	67
	<u>7.2 The 35 bar truss with 15 design variables</u>	71
	<u>7.3 Analytical function--20 design variables</u>	74
	<u>7.4 Conclusion</u>	76
8.	Augmented Minimum Point Designs	77
	<u>8.1 Introduction</u>	77
	<u>8.2 Augmented Minimum Point Designs for 2nd Order Approximations</u>	79
	<u>8.2.1 Specifics of program DESIGNS</u>	80
	<u>8.3 Augmented Minimum Point Design for 3rd and 4th Order</u>	
	<u>Approximation</u>	81
	<u>8.3.1 Specifics of program DESIGN4</u>	82
	<u>8.4 Conclusion</u>	83
9.	Solution Algorithm	84
	<u>9.1 Introduction</u>	84
	<u>9.2 Program Specifics</u>	84
10.	Conclusion	86
11.	References	89

1. Introduction

The mathematical formulation of the engineering optimization problem is

$$\begin{aligned} \min \quad & f(\mathbf{x}) \\ \text{subject to} \quad & g_i(\mathbf{x}) \leq 0, \quad i=1,q \end{aligned} \quad (1)$$

where

\mathbf{x} is an $n \times 1$ matrix of design variables,

$f(\mathbf{x})$ is the objective function, and

$g_i(\mathbf{x})$ are constraint equations.

Evaluation of the objective function and constraint equations in Equation (1) can be very expensive in a computational sense. Thus, it is desirable to use as few evaluations as possible in obtaining its solution. In solving Equation (1), one approach is to develop approximations to the objective function and/or restraint equations and then to solve Equation (1) using these approximations in place of the original functions. These approximations are referred to as response surfaces.

The desirability of using response surfaces depends upon the number of functional evaluations required to build the response surfaces compared to the number required in the direct solution of Equation (1) without approximations. The present study is concerned with evaluating the performance of response surfaces so that a decision can be made as to their effectiveness in optimization applications. In particular, this study focuses on how the

quality of approximations is effected by design selection. Polynomial approximations and neural net approximations are considered.

To provide the groundwork for future discussion, this introductory section discusses:

1. measures of quality of fit at the designs and measures of quality of fit over a region of interest and
2. the methodology used to build the approximations.

1.1 Quality of Fit

Let us consider a problem with n design variables, the components of the vector $\{\mathbf{x}\} = \{x_1, x_2, \dots, x_n\}^t$. A total of N designs will be considered: $\{\mathbf{x}\}_j, j=1, N$. At the designs $\{\mathbf{x}\}_j$, let

y_j = the value of the function to be approximated and

\hat{y}_j = the value of the approximating function.

The approximating function, \hat{y} , should closely match the function, y , not only at the designs, $\{\mathbf{x}\}_j$, but over the entire region of interest.

1.1.1 Fit at the designs

The approximating function \hat{y} closely approximates the function y when s is small where

$$s = \sqrt{\frac{\delta^2}{N}} \quad (2)$$

and where δ^2 is the sum of the squares of the residuals thus

$$\delta^2 = \sum_{i=1}^N (y_i - \hat{y}_i)^2 \quad (3)$$

Let \bar{y} be the average value of the designs, y_i . Thus

$$\bar{y} = \frac{\sum_{i=1}^N y_i}{N} \quad (4)$$

In this study, one measure of the closeness of fit to be considered is the non-dimensional value v where

$$v = \frac{\sqrt{\frac{\sum_{i=1}^N (y_i - \hat{y}_i)^2}{N}}}{\bar{y}} * 100 \quad (5)$$

The coefficient v is the non-dimensional root mean square (RMS) error at the designs. Thus, $v = 0$ is a necessary and sufficient condition that the approximating function fit the actual function at the N design points.

1.1.2 Overall fit

Just because the approximating function exactly fits the function at N designs does not guarantee that it gives a good fit over the region of interest. It is therefore desirable over the region of interest to have a measure of the quality of overall fit. Several examples of this study considers a two dimensional region of interest. For these problems, the

rectangular region of interest is overlaid with a 31x31 evenly spaced grid of points. The value of the function and the approximating function is then compared at these NG=961 evenly spaced grid of points. Other examples consider a rectangular n dimensional region of interest. These regions of interest are also overlaid with a evenly spaced grid of points. The value of the function and the approximating function are then compared at these NG grid points. For these examples, a measure of the quality of overall fit is taken as

$$v_G = \frac{\sqrt{\frac{\sum_{i=1}^{NG} (y_i - \hat{y}_i)^2 / NG}{\bar{y}_G}}}{\bar{y}_G} * 100 \quad (6)$$

where \bar{y}_G is the average value of y at the grid points. A small value of v_G indicates that the approximating function did a good job of approximation over the region of interest.

1.2. Polynomial Approximations

With the polynomial response surface approach, the approximating function is taken as an $m=k+1$ term polynomial expression [1-3] thus

$$\hat{y} = b_0 + b_1 X_1 + \dots + b_k X_k \quad (7)$$

where X_j is some expression involving the design variables. For example, a second order polynomial approximation in two variables could be of the form

$$\hat{y} = b_0 + b_1 x_1 + b_2 x_2 + b_3 x_1^2 + b_4 x_1 x_2 + b_5 x_2^2 \quad (8)$$

The value of the function to be approximated at the N designs can be used to determine the $m = k + 1$ undetermined coefficients in the polynomial expression. For the N designs, Equation (7) yields

$$\begin{Bmatrix} y_1 \\ y_2 \\ \vdots \\ y_N \end{Bmatrix} = \begin{bmatrix} 1 & X_{1_1} & \dots & X_{k_1} \\ 1 & X_{1_2} & \dots & X_{k_2} \\ \dots & \dots & \dots & \dots \\ 1 & X_{1_N} & \dots & X_{k_N} \end{bmatrix} \begin{Bmatrix} b_0 \\ b_1 \\ \vdots \\ b_k \end{Bmatrix} \quad (9)$$

or

$$\{Y\} = [Z]\{b\} \quad (10)$$

where $\{Y\}$ is an $N \times 1$ matrix, $[Z]$ is an $N \times m$ matrix, and $\{b\}$ is an $m \times 1$ matrix.

1.2.1 Exactly-determined approximation

When $N = m$, the approximation is exactly-determined and the matrix $\{b\}$ can be determined from Equation (10).

1.2.2 Over-determined approximation

With $N > m$, Equation (10) can be solved in a least squares sense thus [1-3]

$$[Z]^t \{Y\} = [Z]^t [Z] \{b\} \quad (11)$$

or

$$\{b\} = ([Z]^t [Z])^{-1} [Z]^t \{Y\} \quad (12)$$

Equation (12) in effect, chooses the terms of $\{b\}$ so as to minimize the square of the residual as defined in Equation (2).

1.2.3 Under-determined approximation

When $N < m$, the approximation is under-determined. A solution can be obtained by choosing the terms of $\{b\}$ so as to minimize the square of the residual as defined in Equation (2). However, a direct solution can be obtained by using the concept of pseudo-inverse [4,5]. Assume that the rank of matrix $[Z]$ is N and define the pseudo-inverse of matrix Z , Z^* thus

$$[Z]^* = [Z]^t ([Z][Z]^t)^{-1} \quad (13)$$

where t denotes transpose. Solution of Equation (10) is then

$$\{b\} = [Z]^* \{Y\} + [Q] \{w\} \quad (14)$$

where $\{w\}$ is an $(m-N)$ column matrix of arbitrary coefficients and $[Q]$ is a $m \times (m-N)$ matrix formed from any $m-N$ independent columns of the matrix $[R]$ thus

$$[R] = [I] - [Z]^* [Z] \quad (15)$$

One solution to Equation (14) is to take all the arbitrary terms of $\{w\}$ as zero giving

$$\{b\} = [Z]^* \{Y\} \quad (16)$$

The basic solution to Equation (10) is Equation (16). Using that equation, at the designs, $\{x\}_j$, the value of \hat{y}_j matches the value of y_j . If w_i is the i th term in matrix $\{w\}$ and $\{q\}_i$ is the i th column of matrix $[Q]$, then at the designs, $\{x\}_j$, $\hat{y}_j = 0$ when

$$\{b\} = w_i \{q\}_i \quad (17)$$

Thus, the last term of the right hand side of Equation (14) gives \hat{y}_j values which match y_j at the designs, $\{x\}_j$, for any values of w_i .

1.3 Artificial Neural Nets

While the initial motivation for developing artificial neural nets was to develop computer models that could imitate certain brain functions, neural nets can be thought of as another way of developing a response surface. Different types of neural nets are available [6,7], but the type of neural nets considered in this paper are back propagation nets with one hidden layer as shown in Figure 1. This type of neural net has been used previously to develop

response surfaces [8-12] and is capable, with enough nodes on the hidden layer, of approximating any continuous function [13].

For the neural net of Figure 1, associated with each node on the hidden layer, node j , and each output node, node k , are coefficients or weights, θ_j and θ_k , respectively. These weights are referred to as the biases. Associated with each path, from an input node i to node j on the hidden layer, is an associated weight, w_{ij} and from node j on the hidden layer to output node k is an associated weight w_{jk} . Let q_i be inputs entered at node i . Node j on the hidden layer receives weighted inputs, $w_{ij}q_i$. It sums these inputs and uses an activation function to yield an output r_j . The activation function considered in this paper is the sigmoid function [6,7]

$$r_j = \frac{1}{1 + e^{-\sum w_{ij}q_i - \theta_j}} \quad (18)$$

Output node k then receives inputs $w_{jk}r_j$ which are summed and used with an activation function to yield an output s_k . Some variation of the delta-error back propagation algorithm [6,7] is then used to adjust the weights on each learning try so as to reduce the values between the predicted and desired outputs. In this investigation, studies were performed using the program NEWNET [14] which was developed especially for this investigation. NEWNET minimizes the sum of the squares of the residuals in Equation (2) with respect to the weights and biases of the net. Training of the net is thus formulated as an unconstrained minimization problem. Solution of this minimization problem is performed

using the method of Davidon, Fletcher, and Powell [15-16]. That algorithm performs a series of one dimensional searches along search directions. Search directions are determined by building an approximation to the inverse Hessian matrix using gradient information. Gradients required by that algorithm are obtained using back-propagation. One-dimensional searches are performed along the search directions using an interval shortening routine.

2. Levels of Designs

2.1 Taylor Series Approximation

The overriding factor which affects the accuracy of an approximation is the levels of the design parameters considered. It is instructive to consider a problem in two design variables. Suppose we wish to make a quadratic approximation of a function thus:

$$y = b_0 + b_1x_1 + b_2x_2 + b_3x_1^2 + b_4x_1x_2 + b_5x_2^2 \dots \quad (19)$$

Consider that the exact function is evaluated at 6 design points and the information thus generated will be used to determine the 6 undetermined coefficients in Equation (19).

Design variables at these design points are taken from the following sets:

$$\begin{aligned} x_1 & \text{ from the set } \{x_{1_1} \ x_{1_2} \dots x_{1_p}\} \\ x_2 & \text{ from the set } \{x_{2_1} \ x_{2_2} \dots x_{2_q}\} \end{aligned} \quad (20)$$

Here p discrete values are considered for x_1 and q discrete values are considered for x_2 . The variable x_1 is said to have p levels and x_2 is said to have q levels. The problem is to determine the minimum levels of the design variables, p and q , required to build the quadratic approximation. In this regard, it is instructive to consider a Taylor series approximation [17] of the function about the point $\{x_1=0, x_2=0\}$:

$$\tilde{y} = y(0,0) + \{\nabla y(0,0)\}'\{\Delta x\} + \{\Delta x\}'[H(0,0)]\{\Delta x\} + \dots \quad (21)$$

where

$$\{\Delta x\} = [(x_1 - 0) \ (x_2 - 0)]^t = [x_1 \ x_2]^t \quad (22)$$

$$\{\nabla y(0,0)\} = \left[\left(\frac{\partial y(0,0)}{\partial x_1} \quad \frac{\partial y(0,0)}{\partial x_2} \right) \right]^t \quad (23)$$

$$[H(0,0)] = \begin{bmatrix} \frac{\partial^2 y(0,0)}{\partial x_1^2} & \frac{\partial^2 y(0,0)}{\partial x_1 \partial x_2} \\ \frac{\partial^2 y(0,0)}{\partial x_1 \partial x_2} & \frac{\partial^2 y(0,0)}{\partial x_2^2} \end{bmatrix} \quad (24)$$

Entering Equations (22), (23), and (24) into Equation (21) gives

$$\begin{aligned} \tilde{y} = & y(0,0) + \frac{\partial y(0,0)}{\partial x_1} x_1 + \frac{\partial y(0,0)}{\partial x_2} x_2 + \frac{\partial^2 y(0,0)}{\partial x_1^2} x_1^2 + \\ & 2 \frac{\partial^2 y(0,0)}{\partial x_1 x_2} x_1 x_2 + \frac{\partial^2 y(0,0)}{\partial x_2^2} x_2^2 \end{aligned} \quad (25)$$

The derivatives in Equation (25) can be determined by finite difference equations [18]. The second derivative of y with respect to x_1 can be obtained using information at points indicated in Figure 2 by solid circles, the second derivative of y with respect to x_2 can be

obtained using information at points indicated by unfilled circles, and the mixed derivative can be obtained using information at points indicated by unfilled squares.

It can be seen in Figure 2 that at least three levels of both x_1 and x_2 must be used to obtain a quadratic approximation. If three levels are not provided, not information is available to calculate the higher derivatives in Equation (25). A complete 3 factorial design does not have to be used--only 6 selected points from the complete 3 factorial design. Information at those 6 points allow the undetermined coefficients to be exactly determined.

Consider now the design of Figure 3 which are also taken from the 3 factorial design. Even though 6 design points are used, this set of design points does not allow an approximation containing the x_2^2 term of Equation (25). However, with the design of Figure 3, an approximation of the form of Equation (26) could be obtained thus:

$$y = b_0 + b_1x_1 + b_2x_2 + b_3x_1^2 + b_4x_1x_2 \quad (26)$$

With the design of Figure 3, if a solution is attempted using Equations (19) and (12), a singular coefficient matrix will be encountered. A solution could be attempted using the pseudo-inverse concept of Equations (13) and (14). However, recent studies [19] have shown that non-unique solutions are obtained with this technique. Non-uniqueness makes these solutions undesirable. Using Equations (26) and (12), a slightly over-determined approximation is obtained.

Recent studies have found that the numerical performance of neural network approximations and polynomial approximations with the same number of associated undetermined parameters is comparable [19]. Thus, it is not expected that neural nets as approximators will perform better than polynomials when there are inadequacies in the training design, as in Figure 3. The next example investigates performance of both polynomial and neural net approximations.

2.2 Example

Consider the function

$$y=1+x_1+x_2+x_3+x_1^2+x_1x_2+x_1x_3+x_2^2+x_2x_3+x_3^2 \quad (27)$$

In the first phase of the investigation, approximations are to be made of this function using the design of Figure 4. The star pattern of design points in Figure 4 does not allow mixed derivatives of the function to be calculated using finite difference type formulae but does permit the other second derivatives to be calculated. Thus, information is available to make a polynomial approximation of the form

$$\tilde{y}=b_0+b_1x_1+b_2x_2+b_3x_3+b_4x_1^2+b_5x_2^2+b_6x_3^2 \quad (28)$$

The function y was evaluated at the design points shown in Figure 4 yielding 7 training pairs for calculating the 7 undetermined parameters in Equation (28). The value of the approximating function \hat{y} was then evaluated at a 5x5x5 grid of designs. These values of \hat{y}

were then used to evaluate v_G from Equation (6). The value of v_G obtained is shown in the first line of Table 2.1.

Table 2.1. Performance of Approximations for Various Designs

Number Designs Points	Description	Polynomial Approximation		Neural Net Approximation			
		No. Para.	v_G (%)	ih	No. Para.	No. Apx.	v_G (%)
7	Star--see Figure 4	7	34.6	2	11	10	25.5-97.3
12	Star--see Figure 5	7	34.6	2	11	10	32.9-93.5
10	Computer Generated	10	0.0	2	11	10	36.6-36.9
				3	16	10	21.9-36.7
27	3 factorial	10	0.0	3	16	2	16.6-16.7
				4	21	2	16.6-16.9
125	5 factorial	10	0.0	8	41	1	3.7

A neural net approximation was then considered. Previous studies [19] have indicated that it is desirable to have more training pairs than the number of undetermined parameters (weights and biases) associated with the net. If fewer training pairs than undetermined parameters are used, non-unique approximations should be expected. For a neural net with one hidden layer as shown in Figure 1, there are 6 parameters associated with a net with one node on the hidden layer and 11 parameters associated with a net with two nodes on the hidden layer. It was considered that one node on the hidden layer would yield an inadequate approximation. Thus 2 nodes on the hidden layer were considered. Thus, the

neural net approximation is under-determined. That is to say that there are fewer training pairs than there are undetermined parameters associated with the approximation. Non-unique approximations are to be expected. Indeed, this was the case. The 8 training pairs were used to make 10 different approximations by having training commence from a different randomly selected set of weights and biases. Once the nets were trained, the value of the approximating function, \hat{y} , was generated at the 5x5x5 set one grid points and the value of v_G was developed. The range of the values obtained is shown in Table 2.1. One can see that a large range of values is obtained. The best neural net approximation is only slightly better than the polynomial approximation while the worst neural net approximation is considerably worse. Just as with the polynomial approximation, the designs used to train the approximation can not yield information necessary to capture essential features of the function to be approximated.

The 12 designs of Figure 5 were next used in the training of a polynomial approximation and a 2 node neural net approximation. Even though more designs are used here than in Figure 4, the additional designs selected do not yield any more information about the nature of the function being approximated. Information is still not available for determining the mixed derivatives of the function to be approximated. Thus, the polynomial approximation of Equation (26) was considered. As there are now more training pairs than there are undetermined parameters, the approximation obtained is over-determined. As no new information is available with the 12 designs, the same polynomial approximation and thus

the same v_G as before are obtained. The value of v_G is shown in the second line of Table 2.1.

A neural net with 2 nodes on the hidden layer was then trained with the 12 training pairs. The net was trained 10 times starting from different randomly selected sets of weights and biases. Even though the number of training pairs, 12, is greater than the number of undetermined parameters associated with the net, 11, non-unique approximations were obtained as can be seen in Table 2.1. Thus, it can be concluded that for neural net approximations, having more training pairs than the number of associated undetermined parameters is only a necessary condition for obtaining a unique approximation but that it is not a sufficient condition. As the 12 designs offered no new information about the function being approximated over that offered by the 8 designs, then just as with the 8 design case, non-unique approximations were obtained.

The program DESIGNS [20], which was developed for this project, was used to generate 10 designs which contain the information necessary for calculating the 10 undetermined coefficients of the complete quadratic approximation of the form:

$$\tilde{y} = b_0 + b_1x_1 + b_2x_2 + b_3x_3 + b_4x_1^2 + b_5x_2^2 + b_6x_3^2 + b_7x_1x_2 + b_8x_1x_3 + b_9x_2x_3 \quad (29)$$

The location of these design points is shown in Figure 6. The polynomial approximation obtained by training the polynomial of Equation (29) with the computer generated designs exactly duplicated the test function of Equation (27). Thus, v_G for the 5x5x5 grid of points

was zero as seen in the third line of Table 2.1.

A neural net with 2 nodes on the hidden layer with 6 associated undetermined parameters and a neural net with 3 nodes on the hidden layer and 11 associated undetermined parameters were then trained 10 times with the computer generated training pairs. Each training started from a different randomly selected set of weights and biases. For the case of 2 nodes on the hidden layer, the approximation generated was over-determined and a unique approximation was obtained (the small range of v_G obtained most likely results from the exit criteria employed in the training algorithm). For the case of 3 nodes on the hidden layer, there are 11 associated undetermined parameters but only 10 training pairs. Thus the approximation is under-determined and a non unique approximation is obtained as can be seen in Table 2.1.

The performance of the neural net approximations was much poorer than that of the polynomial approximation on this problem. This poorer performance may be in part because the problem is biased towards the polynomial approximation as the function being approximated is 2 second order polynomial.

A complete 3^3 factorial design and a 5^3 factorial design were considered to see if good results could be obtained with the neural nets if more training pairs were employed. Indeed this was the case. However, many more training pairs were required to get a good approximation than were required with the polynomial approximation. The extra training

pairs were wasted on the polynomial approximation. Ten correctly selected training pairs is all that is required to get an exact second order approximation. The additional training pairs offered no new information to the polynomial approximation. The coefficient v_G was zero for training pairs using the 3 and 5 factorial designs and a second order polynomial approximation.

2.3 Conclusion

For a given order of approximation, a good design must use an adequate number of levels of the design variables or a poor approximation will be obtained. Likewise, design points must be located so that information is available for determining all of the undetermined coefficients of the approximating function. In many instances, especially when the region of interest is small, a second order polynomial approximation or neural net equivalent will be sufficient to build a response surface. A second order approximation requires a design containing 3 levels of the design variables. Program DESIGNS has been developed to generate a minimum point design which allows all of the coefficients of a second order polynomial approximating function to be obtained. This minimum point design can be augmented by randomly selected design points or by user selected points.

3. Standard Designs

3.1 Underlying Principle

When making a polynomial approximation of a function, the number of design levels required for each design variable depends upon the order of polynomial approximation being used. Consider for example the problem of approximating a function y , a function of one design variable. As previously discussed, two levels of the design variable would be required to make a linear approximation of the function, three levels of the design variable would be required to make a second order approximation, four levels of the design variable would be required to make a 3rd order approximation, etc. If y is a function of r design variables, a p th order polynomial approximation, \hat{y} , requires designs at $p+1$ levels in each design variable.

In response surface methodology, the term factor is used for design variable. A factorial design or factorial experiment is a design in which one uses each of the possible combinations of the levels of each factor. If m is the number of level of each factor and r is the number of factors, then the design would be referred to as a m^r factorial experiment. Table 3.1 gives the number of designs in various factorial experiments.

Table 3.1. Number of designs in a full factorial design

m = level r = factor	2	3	4
2	4	9	16
3	8	27	64
4	16	81	256
10	1024	59049	1.05E06

One can see that even for a small number of factors, complete factorial experiments become impractical if designs are computationally or experimentally expensive to obtain. One then is forced to use some sub-set of the factorial design or alternate designs containing requiring fewer design points. Concepts from statistics are normally used in selecting a sub-set of the factorial design or in developing alternate designs. Thus statistical concepts are reviewed.

3.2 Statistical Concepts

When making an approximation, \hat{y} , of a function, y , most approaches used to select design points for a design consider that

1. polynomial approximations are employed and
2. the value of the function, y_i , determined at the designs, $\{x\}_i$, contains some error, ϵ_i .

A measure of the error at point i is the variance of the error, $\text{var}(\epsilon_i) = \sigma^2$ where

$$\sigma^2 = \sum_{i=1}^n \frac{(y_i - \mu)^2}{n} \quad (30)$$

where

μ is the true mean of all possible observations of y_i and

n is the number of observations made.

In experimental investigations, ϵ_i is experimental error. When making approximations to analytical functions, ϵ_i is zero and the variance of the error at point i is zero. Often approximations are made to a function whose values must be obtained from some numerical algorithm such as the finite element method or finite difference method. Values of y_i obtained from such algorithms depend on control parameters which dictate the level of accuracy of the solution. For example, if y was a stress determined from a finite element analysis, then y could depend on a control parameter which specifies the coarseness of the finite element idealization. In this case, different values of y_i would be obtained for the i th design for different values of the control parameters and ϵ_i could be thought of as a numerical error.

It would be an interesting study to select designs such that approximations developed are insensitive to numerical errors such as finite element idealization error. However, the problem at hand is to find a good approximation to an analytical function or a good

approximation for output from a deterministic model. For the problem at hand, for a given design, x_i , one obtains the same functional value, y_i , no matter how many times the function is evaluated. Thus, the problems considered in this report contain no numerical error. However, as all known algorithms with one exception [21] consider that there is some experimental or numerical error, this section now further examines this case.

Errors in the value of y_i used to build an approximation affect the estimation of the undetermined coefficients, b_j , in the polynomial approximation and thus affect \hat{y}_i , the values of y_i predicted by the approximation. A measure of the error in b_j resulting from errors in y_i is the variance of b_j . For example, consider that y_i is obtained from a finite element analysis and that a p th order polynomial approximation is employed. The undetermined coefficients in that approximations, b_j , can be determined from Equation (12). If a number of approximations were now made with finite element results, obtained using different idealizations, the coefficient b_j for these approximations would be different. The variance of b_j is a measure of how much the b 's change for these different approximations. In like form, the different approximations yield different \hat{y}_i and the variance of \hat{y}_i is a measure of how much the \hat{y}_i values change from approximation to approximation.

From a numerical standpoint, it is desirable to have approximations that are not highly sensitive to the error ϵ_i . Approximations are insensitive to the error, ϵ_i , if the variance of b_j and the variance of \hat{y}_i is small. Most design selection algorithms currently in use attempt in some way to keep these variances small.

The variance of b_j is the j,j term of the variance-covariance matrix $\text{cov } b$ where (see Equation 3.11 of [3] or Equation 2.8 of [2])

$$[\text{cov } b] = \sigma^2([Z]'[Z])^{-1} \quad (31)$$

and the variance of \hat{y}_i is given by (see Equation 2.11 of [2])

$$\text{var } \hat{y}_i = \sigma^2 \{Z_i\}'([Z]'[Z])^{-1} \{Z_i\} \quad (32)$$

where $\{Z_i\}'$ is the $1 \times p$ vector whose elements correspond to the elements of a row of matrix $[Z]$.

Notice that these variance involve the matrix $[H]$ where

$$[H] = ([Z]'[Z])^{-1} \quad (33)$$

Design selection affects $[Z]$, which from Equation (33) affects $[H]$, which in turn affects the variances of b_j and \hat{y}_i . Many design point selection algorithms attempt to select designs which give an $[H]$ matrix which will keep the variances of b_j and \hat{y}_i small.

3.3 Orthogonal Designs

The associated undetermined coefficients of a polynomial approximation function can be found from Equation (12). The solution for these coefficients involve the matrix $[Z]$ (see Equations (9) and (10)). Let $\{Z_i\}$ be the i th column of matrix $[Z]$. A design is said to be

orthogonal if the columns of the $[Z]$ matrix are orthogonal, i.e. $\{Z_i\}^t\{Z_j\}=0, i \neq j$. There are interesting properties of orthogonal designs which have prompted their use. Thus orthogonal designs will now be presented in some detail.

3.3.1 Scaling

The discussion of orthogonality is simplified by working with scaled variables. Consider that the approximation in question involves k unscaled design variables \bar{x}_i and contains N design points. Instead of working with \bar{x}_i , the variables will be scaled. Let \bar{x}_{iu} be the u th level of unscaled variable i and x_{iu} be the scaled level. The desired scaling is

$$\sum_{u=1}^N x_{iu}^2 = N, \quad i=1, k \quad (34)$$

$$\sum_{u=1}^N x_{iu} = 0, \quad i=1, k \quad (35)$$

This scaling can be accomplished by having

$$x_{iu} = \frac{\bar{x}_{iu} - \bar{\bar{x}}_i}{S_i} \quad (36)$$

where

$$\bar{\bar{x}}_i = \text{the average of the levels of } \bar{x}_i \quad (37)$$

and

$$S_i^2 = \sum_{u=1}^N \frac{(\bar{x}_{iu} - \bar{x}_i)^2}{N} \quad (38)$$

With this scaling, N experimental design points of the orthogonal design give

$$[Z]'[Z] = N[I] \quad (39)$$

$$([Z]'[Z])^{-1} = \frac{1}{N}[I] \quad (40)$$

where [I] is the identity matrix.

3.3.1.1 Example of Scaled Designs:

Consider a 2 factorial design with levels of 4 and -4. For that design

$$\tilde{x}_1 = 0, \quad \tilde{x}_2 = 0 \quad (41)$$

and

$$S_1^2 = S_2^2 = \frac{(4-0)^2 + (-4-0)^2}{2}, \quad \text{or} \quad S_1 = S_2 = 4 \quad (42)$$

From Equation (3), the levels of the scaled variables are

$$x_{iu} = \frac{\bar{x}_{iu} - 0}{4} \quad (43)$$

or the levels of the scaled variables are 1 and -1.

3.3.2 Bias

Assume that the polynomial approximating function is inadequate. The coefficients of that polynomial can be determined from Equation (12). Let $\{\hat{b}_1\}$ be the coefficients thus obtained and let $[Z_1]$ be the corresponding $[Z]$ matrix. Then from Equation (12)

$$\{\hat{b}_1\} = ([Z_1]^t [Z_1])^{-1} [Z_1]^t \{Y\} \quad (44)$$

Assume that the function being approximated can be expressed as

$$\{Y\} = [Z] \{b\} \quad (45)$$

where

$$\{b\} = \begin{Bmatrix} \{b_1\} \\ \{b_2\} \end{Bmatrix}, \quad [Z] = [\begin{array}{cc} [Z_1] & [Z_2] \end{array}] \quad (46)$$

Entering Equations (40), (45), and (46) into Equation (44) gives

$$\{\hat{b}_1\} = \frac{1}{N} [I] [Z_1]^t ([Z_1] \quad [Z_2]) \begin{Bmatrix} \{b_1\} \\ \{b_2\} \end{Bmatrix} \quad (47)$$

Entering Equation (39) into Equation (47) gives

$$\{\hat{b}_1\} = \frac{1}{N} (N[I] \{b_1\} + [Z_1]^t [Z_2] \{b_2\}) \quad (48)$$

or

$$\{\hat{b}_1\} = \{b_1\} + \frac{1}{N}[Z_1]^t[Z_2]\{b_2\} = \{b_1\} + [A]\{b_2\} \quad (49)$$

where $[A]$ is called the alias matrix. One can see in Equation (49) that the coefficients $\{\hat{b}_1\}$ will only be correct estimates of $\{b_1\}$ if the columns of $[Z_1]$ are orthogonal to the columns of $[Z_2]$. Special situations where this orthogonality occurs are next discussed.

3.3.2.1 A bias example--linear approximating polynomial but the exact function contains linear terms and cross-product terms:

Consider a linear approximating polynomial

$$\hat{y} = \hat{b}_o + \sum_{i=1}^k \hat{b}_i x_i \quad (50)$$

where the exact function is

$$y = b_o + \sum_{i=1}^k b_i x_i + \sum_{i=1}^k \sum_{j=i}^k b_{ij} x_i x_j \quad (51)$$

where b_{ij} are the undetermined coefficients associated with the cross-product terms. For this problem, a full 2^k factorial design gives that the columns of $[Z_1]$ are orthogonal to the columns of $[Z_2]$ and thus

$$\{\hat{b}_1\} = \{b_1\} \quad (52)$$

3.3.2.2 A bias example--linear approximating function but the exact function is a complete quadratic polynomial:

Consider a linear approximating polynomial

$$\hat{y} = \hat{b}_o + \sum_{i=1}^k \hat{b}_i x_i \quad (53)$$

where the exact function is a complete second order polynomial thus

$$y = b_o + \sum_{i=1}^k b_{ii} x_i^2 + \sum_{i=1}^k \sum_{j=i}^k b_{ij} x_i x_j \quad (54)$$

Assume again that a full 2^k factorial design is used. For this problem the alias matrix is such that one obtains

$$\begin{aligned} \hat{b}_o &= b_o + \sum_{i=1}^k b_{ii} \\ \hat{b}_j &= b_j, \quad j=1, k \end{aligned} \quad (55)$$

Thus only \hat{b}_o is biased with the other coefficients unbiased or uncorrelated.

3.3.3 Orthogonal Designs for Linear Approximations

For a problem with r design variables, a full 2^r factorial design is an orthogonal design if the approximating function is a first order polynomial. There are several advantages in using such an orthogonal design when the approximating function is assumed to be linear. These advantages are:

1. The solution for the coefficients of the polynomial approximation require a matrix inverse (see Equation (12)). However, when the design is an orthogonal design, that inverse is very easily obtained using Equation (40). Thus there is a small computational advantage in using an orthogonal design.
2. Examples 3.3.2.1 and 3.3.2.2 indicate that under certain conditions, the coefficients obtained using an orthogonal design are unbiased. Obtaining unbiased coefficients is probably more important in developing response surface from experimental results than when developing response surfaces when results are from a deterministic model. With experimental studies, it may be important to ascertain the unbiased values of the linear coefficients. For the deterministic model however, one is looking for an approximating function which gives a good approximation throughout a region of interest. Whether the coefficients of the polynomial approximation are biased or unbiased is of little concern.
3. It can be proven that for linear polynomial approximations, an orthogonal design gives the minimum variance of the coefficients (see page 109 of [3]). It is important when modeling experimental results to obtain a model that is not overly sensitive to experimental error and thus there is an advantage in having a minimum variance of the coefficients.

However, for response surfaces of a deterministic model, variance of the coefficients is not relevant.

3.3.4 Orthogonal Designs for 2nd Order Polynomial Approximations

It is not possible to find an orthogonal design when using a second order polynomial approximating function of the form of Equation (8) (see page 107 of [2]). However, an orthogonal design can be found if one uses as the approximating function a second order orthogonal polynomial (page 130 of [3])

$$\hat{y} = b_o + \sum_{i=1}^k b_i x_i + \sum_{i=1}^k b_{ii} (x_i^2 - \bar{x}_i^2) + \sum_{i=1}^k \sum_{j=i}^k b_{ij} x_i x_j \quad (56)$$

where

$$\bar{x}_j^2 = \frac{\sum_{u=1}^N x_{ju}^2}{N} \quad (57)$$

and where

$$\begin{aligned} N &= \text{the number of design points and} \\ x_{ju} &= x_j \text{ for each of the design points.} \end{aligned} \quad (58)$$

The use of an orthogonal design still gives the small computational advantage that the inverse shown in Equation (12) is an inverse of a diagonal matrix. However, when using

second order approximations, it is not clear under what conditions one obtains unbiased coefficients. Also it can not be proven that orthogonal designs any longer give a minimum variance of the coefficients. Thus most of the reasons for using orthogonal designs found for linear approximations are not present when using second order approximations.

3.3.5 General Discussion of Orthogonal Designs

Orthogonal designs offer a small computational advantage that the matrix inverse required in solving for the coefficients of the polynomial approximating function is an inverse of a diagonal matrix. When approximating a deterministic model, properties of orthogonal designs which minimize the variance of the coefficients and which give unbiased coefficients are unimportant. For this case, the use of orthogonal designs can only be justified by how well they perform on test problems. Such test problems are presented later in this report.

3.4 Central Composite Designs--Designs for Fitting Second Order Models

It was shown in Section 2 that at least 3 levels of the design variables are required if one is to make a second order approximation. A workable alternative to using a 3^k factorial design is a class of designs called the central composite design. These types of designs are widely used by workers applying second order response surface techniques [3].

3.4.1 Format of the central composite design

The central composite design is a design composed of the 2^k factorial design augmented by additional points. The augmented design points are as follows:

$$\begin{array}{cccccc}
x_1 & x_2 & x_3 & \dots & x_k & \\
0 & 0 & 0 & \dots & 0 & \\
-\alpha & 0 & 0 & \dots & 0 & \\
\alpha & 0 & 0 & \dots & 0 & \\
0 & -\alpha & 0 & \dots & 0 & \\
0 & \alpha & 0 & \dots & 0 & \\
\dots & \dots & \dots & \dots & \dots & \\
0 & 0 & 0 & \dots & -\alpha & \\
0 & 0 & 0 & \dots & \alpha &
\end{array} \tag{59}$$

Figure 7 shows a central composite design for $k=3$. The value of α and the number of design points at the center of the design are varied to meet certain conditions. In the following, those conditions are chosen assuming that the approximating polynomial function is given by Equation (56).

3.4.1.1 Single center point rotatable second order experimental designs:

A design is said to be rotatable when the variance of the estimated response--that is, the variance of \hat{y} , which in general is a function of position in the design space, is instead only a function of the distance from the center of the design and not on the direction. In other words, a rotatable design is one for which the quality of the estimator \hat{y} is the same for two points that are the same distance from the center of the design [3]. It is possible to develop central composite designs which have a single center point. The value of α which will yield these rotatable second order designs are given in Table 3.2.

Table 3.2. Value of α for single center point rotatable central composite designs

k	α
2	1.414
3	1.682
4	2.000
5	2.378
5 (1/2 rep)	2.000
6	2.828
6 (1/2 rep)	2.378
7	3.364
7 (1/2 rep)	2.828
8	4.000
8 (1/2 rep)	3.364

Note in Table 3.2 that a rotatable second order experimental design can be obtained with a fractional factorial design augmented with additional design points as well as with a augmented full factorial design.

3.4.1.2 Multiple center point rotatable uniform precision designs:

In general, the variance of \hat{y} varies with distance from the center of the design. However, by varying the number of center points, N , the variance at a distance of unity from the center can be made approximately equal to the variance at the center of the design. Such designs are referred to as uniform precision designs. The uniform precision design is based on the philosophy that in the central region of the design space there should be uniform importance as far as the variance of response is concerned, as opposed to, for example, a

situation in which the variance is low in the center of the design but increases drastically as one moves away from the design center [3]. The number of center points, m , and the value of α can be varied so as to obtain a rotatable uniform precision designs. Table 3.3 gives those values.

Table 3.3. Values of m and α for multiple center point rotatable uniform precision designs

k	m	α
2	5	1.414
3	6	1.682
4	7	2.000
5	10	2.378
5 (1/2 rep)	6	2.000
6	15	2.828
6 (1/2 rep)	9	2.378
7 (1/2 rep)	14	2.828
8 (1/2 rep)	20	3.364

3.4.1.3 Single center point orthogonal central composite designs:

An orthogonal central composite design can be developed where $[Z]^t[Z]$ is diagonal. To obtain a design of this type a single center point can be used and the α value are taken from Table 3.4.

Table 3.4. Values of α for single center point orthogonal central composite designs

k	α
2	1.000
3	1.216
4	1.414
5	1.596
6	1.761
7	1.910
8	2.045

3.4.1.4 Rotatable orthogonal designs:

By varying the number of designs at the design center, m , and by selecting appropriate values for α , an orthogonal rotatable central composite design can be obtained. Values of m and α for such a design are given in Table 3.5.

Table 3.5. The value of m and α for multiple center point orthogonal rotatable central composite designs

k	m	α
2	8	1.414
3	9	1.682
4	12	2.000
5	17	2.378
5 (1/2 rep)	10	2.000
6	24	2.828
6 (1/2 rep)	15	2.378
7 (1/2 rep)	22	2.828
8 (1/2 rep)	33	3.364

3.4.2 Discussion of the central composite design

Orthogonal central composite designs have been shown to give a variance of response comparable to that obtained with a full 3^k factorial design. Thus, their use is justified when one has experimental error in the response function. Rotatable and uniform precision designs attempt to control the response variance. Thus their use is also justified when one has experimental error in the response function. However, when building a response surface for a deterministic model where there is no experimental error in the response function, their use is justified only by how well they perform on trial problems. Likewise, the designs were developed for the approximating function of Equation (56). If a different second order polynomial approximating function such as in Equation (8) were used or if a neural net was used to develop the response surface, then again the justification for the use of the various

central composite designs would have to be based on their performance on trial problems. Performance of various central composite designs on trial problems is next reported.

3.4.3 Example -- Fox's Banana Function

Fox investigated in Reference [16] a function

$$y=10x_1^4-20x_2x_1^2+10x_2^2+x_1^2-2x_1+5 \quad (60)$$

which has banana shaped contours as seen in Figure 8. The region of interest to be considered is $(-1.5 < x_1 < 1.5, -0.5 < x_2 < 2.0)$.

A second order polynomial approximation is to be made of this function using an orthogonal polynomial approximation as in Equation (56). A two variable orthogonal polynomial approximation is of the form

$$\hat{y}=b_0+b_1x_1+b_2x_2+b_{11}(x_1^2-\bar{x}_1^2)+b_{22}(x_2^2-\bar{x}_2^2)+b_{12}x_1x_2 \quad (61)$$

where

$$\bar{x}_j^2 = \frac{\sum_{u=1}^N x_{ju}^2}{N} \quad (62)$$

and where

$$N = \text{the number of design points and} \quad (63)$$

$$x_{j_n} = x_j \text{ at the design points}$$

In the first phase of this example, Fox's function was approximated using the second order orthogonal polynomial of Equation (61). The designs used in making the approximation were

1. a full 5^2 factorial design,
2. a full 3^2 factorial design,
3. single center point rotatable central composite design,
4. multiple center point rotatable uniform precision central composite design,
5. single center point orthogonal central composite design,
6. multiple center point rotatable orthogonal central composite design,
7. minimum point design from program DESIGNS,
- 8-10. minimum point design from program DESIGNS augmented by additional randomly selected design points, and
11. nine randomly selected design points.

Once an approximation was obtained, the approximate function was evaluated at a 31×31 grid of points over the region of interest. The approximate function values at these 961 points were used to develop the error parameter v_G from Equation (6). Because there are a differing number of functional evaluations required for each of the sundry designs tested, a comparison of the designs based on v_G is misleading. For example, the full 5^2 factorial

design has 25 design points each requiring a functional evaluation where as the multiple center point rotatable orthogonal central composite design has but 16 design points requiring 9 functional evaluations (in the following it is assumed that the function being approximated has no experimental or numerical error and thus the 8 design points at the design center require but one functional evaluation). Thus a comparison of performance based only on quality of fit is not a fair comparison. The 5^2 factorial might do a better job of approximating a function but the computational cost of the $25-9=16$ extra functional evaluations might make it a less desirable design.

For each design, design j , a measure of efficiency, E_j , was developed where

$$E_j = \frac{(v_G)_{design\ j} T_{design\ j}}{(v_G)_{design\ 1} T_{design\ 1}} \quad (64)$$

where T is the number of functional evaluations required for a given design.

The efficiency of all the designs was compared to design 1, the 5^2 factorial design. Table 3.6 gives, for each design tested, the number of design points, N ; for central composite designs, the number of design points at the center of the design, m ; the number of functional evaluations required, T ; the value of v ; the value of v_G ; and the value of E_j .

Table 3.6. Performance of various designs on Fox's Banana Function, orthogonal polynomial approximating function, $-1.5 < x_1 < 1.5$, $-.5 < x_2 < 2.0$

Design	N	m	T	v	v_G	E_j
5^2 factorial design	25	...	25	70.76	78.92	1.00
3^2 factorial design	9	...	9	64.07	102.46	.47
single center point rotatable central composite design	9	1	9	54.36	77.34	.35
multiple center point rotatable uniform precision central composite design	13	5	9	53.08	77.34	.35
single center point orthogonal central composite design	9	1	9	64.07	102.46	.47
multiple center point rotatable orthogonal central composite design	16	8	9	51.62	77.34	.35
minimum point design from program DESIGNS	6	...	6	0	162.62	.49
minimum point design from program DESIGNS augmented by 2 randomly selected design points	8	...	8	43.27	105.16	.43
minimum point design from program DESIGNS augmented by 3 randomly selected design points	9	...	9	53.53	88.63	.40
minimum point design from program DESIGNS augmented by 4 randomly selected design points	10	...	10	53.05	86.44	.44
random--9 points	9	...	9	21.05	460.96	2.10

Several items can be noted in Table 3.6:

1. The design composed of 9 randomly selected design points did poorly. Even though the design points were chosen randomly, it turned out that the design points were not well scattered in the design space but were heavily concentrated in one quadrant of the design space. The polynomial approximation fitted the function well at the design points but poorly over the region of interest.
2. The value of v_G was approximately the same for the single center point rotatable central composite design, the multiple center point rotatable uniform precision central composite design, and the multiple center point rotatable orthogonal central composite design. These three designs differ only in the number of design points at the center of the design space. These designs have 1, 5, and 8 designs at the center, respectively. The effect of putting more designs at the center is to translate the response surface toward the center response. For this problem, however, the actual and approximated response were very close at the design center point, even for only 1 design point at the center. Thus, adding more design points at the design center did little to translate the response surface and thus did not material effect the value of v_G .
3. The eleven designs of Table 3.5 were next used to build an approximation using the standard second order polynomial approximation of Equation (8) instead of the orthogonal polynomial approximation of Equation (61). Results identical to those of Table 3.5 were found. The type of approximating polynomial may effect variances but does not affect quality of fit at the design points or over the region of interest. For those problems where there is no experimental or numerical error associated with functional evaluations, one is

not interested in variance. Thus, there is little advantage in using the orthogonal polynomial approximating functions over a standard second order polynomial function.

4. Based on efficiency, the single center point rotatable central composite design, the rotatable uniform precision central composite design, and the rotatable orthogonal central composite design performed the best but none of the designs gave a good approximation over the region of interest. Over a small region of interest, one could expect that a second order polynomial approximation could well approximate the given function. Obviously, here the region of interest is too large for a second order approximation to be a good one. Thus a smaller region of interest was chosen, $-.5 < x_1, .5$, $-.5 < x_2 < .5$. Table 3.7 compares the eleven designs using this region of interest. Notice that over this smaller region of interest, all the designs gave a much better approximation to the function.

5. For the smaller region of interest, based on efficiency, the 3^2 factorial design, the single center point orthogonal central composite design, and the augmented minimum point designs performed the best. Obviously, the optimum choice of design is problem dependent. However, all designs except the randomly selected design performed much better than the 5^2 factorial design.

Table 3.7. Performance of various designs on Fox's Banana Function, orthogonal polynomial approximating function, $-.5 < x_1 < .5$, $-.5 < x_2 < .5$

Design	N	m	T	v	v_G	E_j
5 ² factorial design	25	...	25	11.16	8.57	1.00
3 ² factorial design	9	...	9	13.27	10.95	.46
single center point rotatable central composite design	9	1	9	6.58	14.74	.62
multiple center point rotatable uniform precision central composite design	13	5	9	5.88	14.74	.62
single center point orthogonal central composite design	9	1	9	13.27	10.95	.46
multiple center point rotatable orthogonal central composite design	16	8	9	5.47	14.74	.62
minimum point design from program DESIGNS	6	...	6	0	18.66	.52
minimum point design from program DESIGNS augmented by 2 randomly selected design points	8	...	8	5.74	11.82	.44
minimum point design from program DESIGNS augmented by 3 randomly selected design points	9	...	9	6.45	10.53	.44
minimum point design from program DESIGNS augmented by 4 randomly selected design points	10	...	10	6.33	10.29	.48
random--9 points	9	...	9	2.42	47.22	1.98

3.4.4 Conclusion

Second order polynomial approximations or neural net equivalents are often adequate for building response surfaces, especially if the region of interest is small. Central composite designs are convenient for building the second order approximations. They provide the necessary information for determining all of the coefficients of the approximating polynomial and give a good distribution of points in the design space. The approximating function can be made to closely fit the exact function at the design center by using multiple center points. When modeling deterministic systems, each functional evaluation at the design center yields the same function value. Thus, for deterministic models, only one functional evaluation need be performed at the center point even when multiple center points are used. Table 3.8 gives information relevant to central composite designs for various number of design variables, k . Central composite designs give over-determined second order polynomial approximations. In other words, there are more design points in the design than there are undetermined coefficients in a second order polynomial approximation. Table 3.8 also gives the percentage that the approximation is over-determined. Previous studies [19] have indicated that designs which give approximations that are around 20-50% over-determined tend to be efficient designs. One can see that the central composite designs are reasonable for $k < 6$. For larger k values, too many design points are being used by the central composite designs. For $k > 5$, an augmented minimum point design is a better choice.

Table 3.8. Information relevant to central composite designs for various number of design variables

Number of design variables, k	Number of coefficients in a 2nd order polynomial approximation	Number of functional evaluations required with a central composite design	% over-determined
1	3	4	33
2	6	8	33
3	10	14	40
4	15	24	60
5	21	42	50
6	28	76	171
7	36	142	294
8	45	272	504

4. Optimality Criteria

4.1 D, A, E, G, and V Optimality Criteria

It was pointed out in Section 3 that even for a small number of factors, a complete factorial experiment become impractical if functional evaluations are computationally or experimentally expensive to obtain and thus one is forced to use some sub-set of the factorial design or an alternate design requiring fewer experiments. Section 3 shows that the variances of the coefficients of a polynomial approximation and the variance of the predicted response involve the matrix $[H]$ given in Equation (33) and repeated here:

$$[H] = ([Z]^t [Z])^{-1} \quad (65)$$

Schoofs [22] lists five criteria for selecting a sub-set of the factorial designs. These criteria involve the matrix $[H]$. The criteria, referred to as optimality criteria, attempt to make $[H]$ minimal. However, "the minimum of a matrix is not a well defined concept and a number of operational criteria have been developed" [22]. The optimality criteria for selecting a subset of a full factorial design can be based on selecting the subset satisfying the following criteria:

1. D-optimality, which is achieved if the determinant of $[H]$ is minimal which in term gives that the product of the eigenvalues of $[H]$ is minimal.
2. A-optimality, which is achieved if the trace of $[H]$ is minimal which in term gives that the sum of the eigenvalues of $[H]$ is minimal.
3. E-optimality, which is achieved if the largest eigenvalue of $[H]$ is minimal.

4. G-optimality, which is achieved if the maximum over all candidate points of the estimated response variance is minimal.
5. V-optimality, which is achieved if the estimated response variance, averaged over all candidate points is minimal.

4.1.1 Criteria Applied to a One Dimensional Example

An example is considered here to compare the performance of the 5 optimality criteria.

The following test function of one variable was considered:

$$y=2+x+\sin\left[\frac{3\pi}{2}(x+1)\right], \quad -1 \leq x \leq 1 \quad (66)$$

This function was approximated with polynomials of order 1-4. The approximations shown in Figure 9 were developed using 13 designs, uniformly spaced in the region of interest. These approximations were then used to generate the functional values at 61 uniformly spaced points in the region of interest which were used to plot the curves of Figure 9.

Further approximations of Equation (66) were developed using various number of design points, n . The designs selected were

1. uniformly spaced design points, $n=5,7,9,11,13$;
2. randomly selected design points, $n=5,7,8,11,13$;
3. an n member subset of the 13 uniformly spaced design points, $n=5,7,9,11$.

Under item 3, the subset of design points was chosen using:

1. D-optimality,
2. A-optimality,
3. E-optimality,
4. G-optimality, and
5. V-optimality.

A FORTRAN program was written to perform the investigation under item 3. The demanding part of the programming was to identify all the possible subsets from the set of thirteen design points. After developing a procedure to identify all combinations, each subset was used to build the $[H]$ matrix. The "optimal" $[H]$ matrix was then determined using the five optimality criteria. The coefficient v_G was then computed for the optimal subset. Figures 10-13 show the value of v_G for the D, A, E, and G optimality criteria when a first, second, third, and fourth order approximation is being made, respectively, versus the number of design points specified in the subset. Also shown in those figures is the value of v_G for designs consisting of design points uniformly spaced in the region of interest.

It was found that for all subsets of size r from a design point set of size n that the estimated response variance, averaged over all candidate points, was invariant. This finding undoubtedly could also be proven theoretically but such a proof was not attempted. From this example, one can conclude that the V optimality criteria, which employs the estimated average response variance, is not a viable criteria for selecting a subset of design points from

a given set. From Figures 10-13, one can see that in most cases there is little difference in the performance of the various optimality criteria with criteria D and G performing slightly better than the other two criteria. As can be seen in Figure 12, on one occasion (when using a third order polynomial approximation and when selecting a subset of 5 design points from the 13 design point set) the G optimality criteria performed poorly while the D criteria did not. Thus, this example indicates that the D optimality criteria may be the criteria of choice. There is a further advantage in using the D optimality criteria. The requirement that the determinant of [H] is minimal is equivalent to a requirement that the determinant of [G] is maximal where

$$[G]=[Z]^T[Z] \quad (67)$$

Thus the D optimality criteria insures that the procedure for determining polynomial coefficients in Equation (12) will be well defined. In other words, Equation (12) uses the inverse of [G]. The D optimality criteria guarantees that [G] is not singular.

One can see in Figures 10-13 that, in most cases, all the optimality criteria performed worst than the uniformly spaced design case. This example indicates that a design picked using an optimality criteria may be no better than a design of the same size in which the design points are uniformly located in the design space.

4.2 S and Q Optimality Criteria

The previous optimality criteria involved only the matrix [H] and did not consider the

function to be approximated. Thus for a given number of design variables and level of approximation, the same designs would be selected no matter what the nature of the function to be approximated. Initially it was thought that a superior optimality criteria would have to consider the nature of the function. Thus two additional optimality criteria were examined:

1. S-optimality, which is achieved if the average error of approximation at the design points is minimal and
2. Q-optimality, which is achieved if the maximum error of approximation at the design points is minimal.

Here

$$\text{average error of approximation} = \frac{\sum_{i=1}^r (y_i - \hat{y}_i)^2}{r} \quad (68)$$

and

$$\text{maximum error of approximation} = \max (y_i - \hat{y}_i)^2, \quad i=1, \dots, r \quad (69)$$

where r is the size of the subset of design points to be selected. One can see that with the S and Q optimality criteria, the function to be approximated effects the design points selected.

4.2.2 Criteria Applied to a One Dimensional Example

The one dimensional example problem of Section 4.1.1 was then re-examined. Figures 14-17 show values of v_G using the S and Q optimality criteria and using a first, second, third, and fourth order polynomial approximation, respectively, versus size of the subset of design points. Also shown in these figures are results for uniformly spaced design points. One can see in these figures that terrible approximations were obtained with these criteria when only small subsets of design points were selected from the original set. Figures 18-20 indicate why such bad approximations are obtained with these two criteria.

Figure 18 depicts results obtained by having eleven design points selected, using the Q optimality criteria, from a set of 13 design points. The Q optimality criteria finds an approximation such that the maximum error of the approximation over eleven design points is minimal. One can see in Figure 18 that the approximating function did indeed well fit the exact function at the 11 design points selected. However, the approximating function did a poor job of approximation at the ends of the region of interest and thus would not yield a low value of v_G . Figure 19 is similar to Figure 18 except that this figure depicts results obtained by having 7 design points selected from the set of 13 design points. One can see that for the optimum design selected, there is an almost perfect approximation at the design points selected but over a much larger region the approximation is poor and thus a large value of v_G would be obtained. In Figure 20, only 5 design points are selected. Again at those design points, an almost perfect approximation is obtained but a terrible approximation is obtained over a large part of the region of interest and thus a large v_G

would be obtained. Thus we can conclude that the S and Q optimality criteria are not operative.

4.3 An Alternate Approach--Random Selection of Designs

The effect of randomly picking design points was next considered. Here designs are picked in the region of interest using a random number generator.

4.3.1 Random Selection of Designs Applied to a One Dimensional Example

For the one dimensional problem under consideration, first, second, third, and fourth order approximations were considered. Design point sets containing 5,7,9,11, and 13 design points were developed by randomly picking design points in the region of interest using a random number generator. Approximations were developed using the design sets. Results using these approximations are compared in Figures 21-24 to results using uniformly spaced design points. One can see in these figures that most of the time results from randomly picked design points are either as good as or not much worst than results from uniformly spaced design points. However, on two occasions, when the number of design points in the design set was small, a relatively poor approximation was obtained. Obviously where one is picking only a small number of points using a random number generator, there is a chance that a bad set of points can be generated and indeed on these two occasion a poor selection of points was made. In general however, when more design points are randomly selected, those points should be scattered throughout the design space and good approximations should be obtained. In conclusion, randomly selecting design points may be a viable method of design selection.

4.4 Larger Problems

Consider a problem in two variables and consider that the potential design points will be taken from a 6 x 6 grid of points. Let

r = total number of design points in the set of potential design points,

c = number of design points in the selected subset of design points,

nc = the number of different combinations of designs in the subset.

For the problem at hand, $r=36$. Subset sizes of $c=15, 20, 25$, and 30 are to be considered.

The number of possible combinations of design points in the subset, nc , is given by

$$nc = \frac{r!}{(r-c)! c!} \quad (70)$$

Table 4.1 summarizes the number of combinations for this study.

Table 4.1 Number of combinations of designs in a two variable study

r Total number of design points	c Number of point in subset	nc Number of combinations
36	15	5,567,902,560
36	20	7,307,872,110
36	25	600,805,296
36	30	1,947,792

One can see that for even small problems, it is infeasible to examine all possible combinations of subsets of size N from a given set of design points. Welch [23], instead of evaluating all possible N -point designs, developed a "branch and bound" algorithm which guarantees global D -optimal designs but which does not generate and evaluate all possible designs. However, even here the computing costs are high. Fedorov [24] developed another technique which neglects the integer character of the components of the design set and obtains a discrete design which is rounded off to an exact design. Reference [22] reports that these designs are considered only approximate. The most popular algorithm seems to be DETMAX by Mitchell [25]. Quoting reference [22], "The algorithm starts with an initial m -point ED (experimental design); the final goal is an optimal N -point ED. During each iteration step that candidate point, which results in the largest increase of $\det(M)$, is added to the design, and subsequently that point, which results in the smallest decrease of $\det(M)$, is removed from the design. The number m of points in the initial design may be larger or smaller than N . If necessary the algorithm first adds (if $m < N$) or rejects (if $m > N$) points until the number of points in the ED is equal to N . In order to avoid local optima the algorithm is able to perform 'excursions', in which several points are added at one go and subsequently the number of points is reduced to N . If the resulting N -point ED has not been improved, another excursion will be made from the same initial design. If the excursion is successful the resulting ED will be used as starting ED in a further attempt to maximize $\det(M)$. The algorithm terminates when, after several excursions, no better ED

is found. The algorithm generates high quality EDs against relatively low computing costs." An attempt is being made to obtain the algorithm DETMAX.

4.5 Optimality Criteria Based on Minimizing Uncertainty

Reference [21] considers problems where there is no experimental error. That reference uses an optimality criteria based on selecting a design which minimizes the uncertainty in the approximating function. That reference was given mixed reviews by a number of leading authorities in the field [21] (reviews follow the paper). The formulation is quite theoretical and difficult to follow. The formulation seems to have promise but requires additional theoretical development before it becomes operative.

4.6 Conclusion

There is little rational for using any of the investigated optimality criteria when building approximations of functions which contain no experimental error. However, the D-optimality criteria can conveniently be used as a heuristic in selecting design points.

Previous investigations have indicated that approximations should be over-determined. That is to say that more training pairs should be used to build an approximations than the number of associated undetermined parameters. It has been suggested that a 20-50% over-determined system might be reasonable. The program DESIGNS, described in Section 2, develops enough designs to exactly determine a quadratic approximation of a given function. The D-optimality criteria can be used as a heuristic for selecting design points to

supplement those generated by DESIGNS. The use of the D-optimality criteria to select the supplementary points would guarantee that no singular matrices would be encountered in determining the undetermined parameters associated with the polynomial approximation.

5. Significance Testing of Coefficients

5.1 Introduction

When the training pairs used to build a polynomial response surface contain experimental or numerical error, certain coefficients in the polynomial approximation may not be significant. In other words, even though one calculates a value for some coefficient, b_i , the experimental or numerical error may have such an effect on that coefficient that it could just as well be taken as zero as the value calculated. In situations like this, it may be advantageous to drop that term from the polynomial approximation and redevelop the response surface. Such a procedure is discussed in pages 34-38 of [3] and an automated procedure for performing such an operation was developed in [26]. Testing of significance involves the t-test which is next described.

5.2 t-test

Coefficients of the polynomial approximation are found from Equation (12). The determination of those coefficients involve the matrix $[H]$ where

$$[H] = ([Z]'[Z])^{-1} \quad (71)$$

A number of terms must now be defined:

$$\text{mean square error} = \text{MSE} = \frac{\sum_{i=1}^N (y_i - \hat{y}_i)^2}{N - m} \quad (72)$$

$$\text{standard error coefficient} = se_i = \sqrt{MSE H_{ii}} \quad (73)$$

$$t_i = \frac{|b_i|}{se_i} \quad (74)$$

where

N = the number of design points and

m = the number of coefficients in the polynomial approximation.

In making the test of significance, t_i from Equation (74) is compared to tabulated values of t_a . The value of t_a is taken from a table of "Percentage Points of the Student's t Distribution" [3]. The value taken depends on the level of significance desired. In lieu of using tabulated values, t_a is often taken as four [26]. If t_i is less than t_a ($t_i < t_a$), then that coefficient's importance in approximating the response is deemed to be insignificant and therefore may be eliminated from the response function.

The primary focus of this study was to examine methods of developing good response surfaces for deterministic models, i.e. for systems that contain no experimental or numerical error. Statistical testing of coefficients presupposes experimental or numerical error and thus is not relevant when approximating response which contains no error. However, the method was thought to perhaps offer a heuristic for improving the quality of a response surface even if experimental or numerical errors are not present. Thus, two examples were

examined. Results are next reported.

5.3 Example 1 -- Fox's Banana Function

Example 1 again examines Fox's Banana Function [16]. A complete second order polynomial approximation ($m=6$) and a complete third order polynomial approximation ($m=10$) were developed. These approximations were developed using a complete 6^2 factorial design ($N=36$). A **t-value**, t_i , was calculated for each parameter, b_i , and compared to $t_a = 4$. Parameter that lack significance ($t_i < t_a$) were eliminated. A new approximation was then developed using only the significant parameters. The values of v and v_G from Equations (5) and (6), respectively, were developed for the complete polynomial and for the polynomial containing only terms deemed significant. Results are shown in Figures 25 and 26. One can see in these figures that eliminating coefficients deemed insignificant had an adverse effect on the quality of the approximation over the region of interest.

5.4 Example 2

The effect of eliminating coefficients deemed insignificant was tested on the function

$$Y=(4+x_1)^3+\sin\left[\frac{3\pi}{2}(x_1+1)\right]+2+x_2^4+\sin\left(\frac{\pi}{2}\right)+7x_2x_1 \quad (75)$$

Again, a complete second order polynomial approximation ($m=6$) and a complete third order polynomial approximation ($m=10$) were developed. These approximations were developed using a complete 6^2 factorial design ($N=36$). A **t-value**, t_i , was calculated for

each parameter, b_i , and compared to $t_a = 4$. Parameter that lack significance ($t_i < t_a$) were eliminated. A new approximation was then developed using only the significant parameters. The values of v and v_G from Equations (5) and (6), respectively, were developed for the complete polynomial and for the polynomial containing only terms deemed significant. Results are shown in Figures 27 and 28. One can see in these figures that eliminating coefficients deemed insignificant offered no improvement in the quality of the response surface.

5.5 Conclusion

The applicability of significance testing of polynomial coefficients when modeling deterministic systems was considered. Two examples were examined to see if eliminating terms of polynomial approximations which were deemed to be insignificant by the t-test would improve the quality of the response surfaces developed. Based on these two examples, it was concluded that no improvement in the predictive capability of response surfaces over regions of interest would be obtained with such a procedure. The relevance of significance testing is when modeling systems containing numerical or experimental error.

6. Applicability of the Response Surface Technique

6.1 Introduction

The following study was performed to ascertain under what circumstances could the response surface technique be used to advantage in engineering optimization application. In this regard, assume that a quadratic polynomial approximations is to be made of functions of n variables. The number of undetermined coefficients in that approximation is:

$$\text{number of coefficients} = \frac{(n+1)(n+2)}{2} \quad (76)$$

Previous studies [19] have shown that the best approximations are obtained when the approximations are over-determined. Thus, the number of functional evaluations required to make the approximation is:

$$\text{number of functional evaluations} = \frac{\delta(n+1)(n+2)}{2} \quad (77)$$

where δ determines the degree that the approximation is over-determined.

The functional evaluations required to build the approximation are initially performed before the start of the optimization process. By using parallel processing, these functional evaluations may be less computationally expensive than evaluations made sequentially in a direct optimization procedure. The number of required evaluations of Equation (77) is then

equivalent to a reduced number of sequential evaluations thus:

$$\text{equivalent number functional evaluations} = \frac{\delta \beta (n+1)(n+2)}{2} \quad (78)$$

where β is a coefficient of efficiency associated with parallel processing.

An optimum solution can be attempted using the response surfaces developed instead of the original functions. However, because of the inexact nature of the approximations, a new set of response surfaces may have to be developed at the most recent approximate solution and another optimal solution attempted. This procedure may have to be repeated α times to reach the optimum solution for the original problem. The total number of equivalent functional evaluations performed in reaching this optimum is:

$$\text{total equivalent functional evaluations} = \frac{\alpha \beta \delta (n+1)(n+2)}{2} \quad (79)$$

If the solutions was attempted by direct optimization techniques instead of using response surfaces, Barthelemy [27] states that a solution can be obtained in most cases using no more than ψ first derivative evaluations. If the first derivatives are obtained by finite difference formulae, an estimate of the number of functional evaluations required by a direct solution procedure is:

$$\text{functional evaluations direct methods} = \psi(n+1) \quad (80)$$

If the response surface technique is to be competitive with the direct solution technique, then from Equations (4) and (5) one must have:

$$\frac{\alpha \beta \delta (n+1)(n+2)}{2} \leq \gamma \psi (n+1) \quad (81)$$

where γ is a convenience factor associated with using response surfaces. In other words, an investigator may tolerate more functional evaluations with the response surface technique than with the direct solution procedure just for the convenience of using response surfaces. Rearranging Equation (81) gives

$$[n+1] \left[\frac{\alpha \beta \delta (n+2)}{2\gamma} - \psi \right] \leq 0 \quad (82)$$

Since $(n+1)$ is positive, one obtains

$$\frac{\alpha \beta \delta (n+2)}{2\gamma} - \psi \leq 0 \quad (83)$$

or

$$n \leq \frac{2\psi\gamma}{\alpha\beta\delta} - 2 \quad (84)$$

In review

$$\begin{aligned}\alpha &= \text{number sequential optimizations} \\ \beta &= \text{parallel processing coefficient} \\ \delta &= \text{overdetermined coefficient} \\ \gamma &= \text{convenience coefficient} \\ \psi &= \text{direct solution coefficient}\end{aligned}\tag{85}$$

Reasonable ranges of the parameters involved are

$$\begin{aligned}\alpha &= 1.00 \rightarrow 4.00 \\ \beta &= 0.10 \rightarrow 1.00 \\ \delta &= 1.25 \rightarrow 1.75 \\ \gamma &= 1.00 \rightarrow 3.00 \\ \psi &= 6.00 \rightarrow 10.0\end{aligned}\tag{86}$$

For an approximate upper bound on the number of design variable that could be economical used with the response surface technique take:

$$\begin{aligned}\alpha &= 1.00 \\ \beta &= 0.10 \\ \delta &= 1.25 \\ \gamma &= 3.00 \\ \psi &= 10.0\end{aligned}\tag{87}$$

giving

$$n \leq 498\tag{88}$$

Under the most unfavorable set of circumstances, that is:

$$\begin{aligned}
 \alpha &= 4.00 \\
 \beta &= 1.00 \\
 \delta &= 1.75 \\
 \gamma &= 1.00 \\
 \psi &= 6.00
 \end{aligned}
 \tag{89}$$

one obtains

$$n \sim 0 \tag{90}$$

Thus depending upon the problem, one could use the response surface technique for $n=0$ to $n=500$ variables. Consider the following reasonable set of parameters

$$\begin{aligned}
 \alpha &= 3.00 \\
 \beta &= 0.50 \\
 \delta &= 1.25 \\
 \gamma &= 1.50 \\
 \psi &= 8.00
 \end{aligned}
 \tag{91}$$

giving

$$n \leq 13 \tag{92}$$

Thus, it is reasonable to assume that the response surface technique could be used for up to 10-15 design variables.

6.2 Conclusion

Under the most advantageous circumstances, the response surface technique applied to engineering optimization application could be used for up to 500 design variables. Under the worst set of circumstances, it is entirely inappropriate. Under normally expected circumstances, this technique might be used to advantage for 10-15 design variables.

7. Additional Examples

7.1 Introduction

The next several examples examine the effect of design selection on the quality of approximations. In each case, a second order polynomial approximation is made of a trial function. Different number of design variables are considered in each example. Thus, for each example different designs are appropriate. In the first example, there are 4 design variables. When there are fewer than 6 design variables, central composite designs are a possible appropriate choice. Other choices are the 3^k factorial design, the minimum point design, the augmented minimum point design, or randomly selected design. All of these designs are considered in that example. In the second and third examples, there are 15 and 20 design variables, respectively. Here, the 3^k factorial design and central composite designs contain too many design points to be practical. For these examples, the minimum point design, the augmented minimum point design, and the randomly selected design are appropriate and are considered.

7.2 The 35 Bar Truss with 4 Design Variables

In many response surface applications, the function to be approximated is a relatively smooth function of the design variables which can be approximated with a lower order polynomial or an artificial neural net with only a few nodes on the hidden layer. A problem of this type is shown in Figure 29. In this example, all loads shown in Figure 29 are in kips, all members of the lower chord of the truss are assumed to have area, A_1 , and all members

of the upper chord to have area, A_2 , all vertical and diagonal members to have area, A_3 . The depth of the truss is H . A response surface is to be constructed for the stress in member BC in terms of the design variables, x_i thus

$$\begin{aligned} x_i &= 1/A_i, \quad i=1,3 \\ x_4 &= .09375H - .4375 \end{aligned} \tag{93}$$

The region of interest is

$$\begin{aligned} 2 \text{ in}^2 &\leq A_i \leq 8 \text{ in}^2 \\ 6 \text{ ft} &\leq H \leq 10 \text{ ft} \end{aligned} \tag{94}$$

or in terms of the design variables

$$.125 \leq x_i \leq .5 \tag{95}$$

A number of designs were used to develop a second order polynomial approximation for the stress in member BC. Each approximation was then used to predict stress on a $5 \times 5 \times 5$ grid of points. The predicted stress and the actual stress on these $NG=625$ grid of points were then used to develop v_G from Equation (6). The parameter v_G is a measure of the quality of the approximation over the region of interest.

The different designs examined required different numbers of functional evaluation. So as to get a measure of the quality of fit of the approximation over the region of interest which

Table 7.1 The 35 bar truss with 5 design variables, 2nd order polynomial approximation

Description	m	α	T	F	v (%)	v _G (%)	E _i
3 ⁴ factorial design	81	81	3.34	2.41	1.00
single center point rotatable central composite design	1	2.000	25	25	0.66	2.67	0.34
multiple center point rotatable uniform precision central composite design	7	2.000	31	25	0.59	2.67	0.34
single center point orthogonal central composite design	1	1.414	25	25	1.47	2.37	0.30
multiple center point rotatable orthogonal central composite design	12	2.000	36	25	0.55	2.67	0.34
minimum point design from program DESIGNS	15	15	0.00	3.99	0.31
minimum point design from program DESIGNS augmented by 3 randomly selected design points	18	18	0.40	3.86	0.36
minimum point design from program DESIGNS augmented by 6 randomly selected design points	21	21	0.38	3.91	0.42
minimum point design from program DESIGNS augmented by 9 selected design points	24	24	0.41	3.77	0.46
randomly selected design	25	25	0.00	824.2	105

m=number of design points at the center of the design space

T = the total number of design points

F = the number of functional evaluations required

α = parameter which defines location of certain design points

takes into account the number of functional evaluations performed, the efficiency, E_j , from Equation (64) was developed for each design. Table 7.1 reports for each design considered, the efficiency, E_j , as well as other relevant information.

One can see in Table 7.1 that all the designs considered, except the randomly selected design, gave a good approximation over the region of interest. Randomly selected designs, which often work well, can sometimes suffer from the problem that the coefficient matrix used to solve for the approximation's associated parameters is poorly conditioned or that the design points selected are not well scattered throughout the design space. In either case, they can yield a poor approximation over the region of interest as in this example.

The 3^4 factorial design well approximated the trial function. However, because it uses so many design points its efficiency measure is poor and thus is not a design of choice. The single center point orthogonal central composite design and the minimum point design from program DESIGNS performed the best, based on their efficiency. However, excluding the randomly selected design and the 3^4 factorial design, all of the designs considered gave a low value of v_G and had approximately the same value of efficiency.

Under normal circumstances, information is not available to calculate v_G and one must use the parameter v as a measure of the quality of fit over the region of interest. However, the parameter v is only a measure of quality of fit over the region of interest if the approximation is over-determined. Thus, under normal circumstances one would not want

to use the minimum point design. This example indicates, that for problems of the size of this example, that any of the central composite designs or the augmented minimum point designs would be appropriate.

7.2 The 35 bar truss with 15 design variables

This example again considers the 35 bar truss of Figure 29. In this example, H is 10 ft., the areas of the 14 bars of the top and bottom chords are A_i , $i=1,14$, and the area of the vertical and diagonal members is A_{15} . The design variables of the problem are taken as

$$x_i = 1/A_i, \quad i=1,15 \quad (96)$$

The region of interest is

$$2 \text{ in}^2 \leq A_i \leq 8 \text{ in}^2 \quad (97)$$

or in terms of the design variables

$$.125 \leq x_i \leq .5 \quad (98)$$

Response surfaces were developed for the stress in member BC using a 2nd order polynomial approximation. The approximation were developed using various designs. To test the quality of the approximations over the region of interest, the function and the approximations were evaluated at $NG=500$ randomly selected test points over the region of interest. That information was then used to calculate v_G from Equation (6). The random

number generator used to develop design points uses, in generating its numbers, an initial seed parameter, IFLAG. A different value of IFLAG was used to generate the 500 test points than was used to generate random points in the randomly selected designs or in the augmented minimum point designs. Thus, the test set of points does not duplicate any of the design points in the designs considered. Results of this investigation are reported in Table 7.2.

One will notice in Table 7.2 that only minimum point designs, augmented minimum point designs, and randomly selected designs are considered. A 3^{15} factorial design contains over 14 million design points. Thus, the use of the 3^{15} factorial design is out of the question. For a problem in k design variables, the central composite design uses a 2^k factorial design augmented by $2k+1$ additional design points. Thus, such a single center point central composite design for this problem contains 32,799 design points. Here again, such a design is impractical. One can develop a central composite design by augmenting only a fraction of the 2^k factorial design. For this problem, a single center point central composite design using only a $1/4$ fraction of the 2^{15} factorial design would contain 8,223 design points which is still an impractical design. Thus, for problems of the size of this example, only the minimum point designs, augmented minimum point designs, and randomly selected designs are of reasonable size.

We can see in Table 7.2 that all of the designs with the exception of the "randomly selected-exactly determined design" did a good job of approximating truss behavior. A singular

matrix was encountered in Equation (10) for the randomly selected--exactly determined design. With completely randomly selected designs, there is always the possibility of having a poorly conditioned coefficient matrix $[Z]$ in Equation (10) and indeed this occurred in this problem. However, there was no problem with matrix conditioning using randomly selected over-determined designs.

Table 7.2 The 35 bar truss with 15 design variables, 2nd order polynomial approximation

Description	F	v %	v _G %	E _j
minimum point design from program DESIGN-exactly determined	136	0	1.263	1.0
augmented minimum point design--20% over-determined	163	0.083	0.294	0.28
augmented minimum point design--40% over-determined	190	0.087	0.060	0.07
random selection--exactly determined	136	*	*	*
random selection--20% over-determined	163	0.003	0.029	0.03
random selection--40% over-determined	190	0.003	0.010	0.01

* singular coefficient matrix

The efficiency parameter, E_j , is calculated in Table 7.2 but it is rather a meaningless parameter for this problem because all the designs so well fit the exact function. In real life

situations, one usually does not have available information for calculating v_G . Thus, the parameter v or like term must be used as a measure of the quality of the approximation. The parameter v is not a meaningful measure of the quality of fit over a region of interest unless the system is over-determined. Thus for this example, the design of choice would be either the 20% over-determined minimum point design or the 20% over-determined randomly selected design.

7.3 Analytical function--20 design variables

This example considers a problem with even more design variables. The function tested is:

$$y = 1. + \sum_{i=1}^{20} x_i + \sum_{i=1}^{20} \sum_{j=1}^{20} x_i x_j + \sum_{i=1}^{20} \sum_{j=1}^{20} x_i^2 * x_j \quad (99)$$

A second order polynomial function was used to build the response surface approximating this function. The polynomial approximating function had 231 undetermined coefficients. Because of the large size of this problem, factorial designs and central composite designs are not appropriate. A minimum point design, augmented minimum point designs, and randomly selected designs were considered. Values of the test function and approximate function were evaluated at $NG = 1000$ randomly selected points and the parameter v_G was developed using this information. The measure of efficiency of the designs examined along with other relevant information is given in Table 7.3.

Table 7.3 Analytical function with 20 design variables, 2nd order polynomial approximation

Description	F	v %	v _G %	E ₁
minimum point design from program DESIGN-exactly determined	231	0	88.93	1.0
augmented minimum point design--20% over-determined	277	5.83	49.82	0.67
augmented minimum point design--40% over-determined	323	9.58	18.03	0.28
random selection--exactly determined	231	*	*	*
random selection--20% over-determined	277	0.61	7.21	0.10
random selection--40% over-determined	323	0.46	1.20	0.02

* poorly conditioned coefficient matrix

Just as in Example 7.2, a exactly determined randomly selected design gave a poorly conditioned coefficient matrix. These examples indicate that randomly selected exactly determined designs should be avoided. The 40% over-determined randomly selected design did an excellent job of modeling the test function and was the most efficient design considered. It seems that on problems with a large number of design variables that randomly selected over-determined designs should be expected to work well.

7.4 Conclusion

The examples of this section have shown that design selection depends on the number of design variables. If the number of design variables is less than 6, appropriate designs are:

1. augmented minimum point designs
2. central composite designs
3. over-determined randomly selected designs.

When there are more than 6 design variables, the central composite designs contain too many design point for consideration. For more than 6 design variables, appropriate designs are then

1. augmented minimum point designs
2. over-determined randomly selected designs.

The example examined indicate that in all cases, over-determined designs should be used. They the most efficient designs. Also, when a design is over-determined the coefficient v can be used as a measure of the quality of the approximation over a region of interest. Being able to use v as a measure of the quality of fit over the region of interest is very important because, in general, information is not available to determined the parameter v_G .

8. Augmented Minimum Point Designs

8.1 Introduction

Design selection in the literature concentrates on linear or quadratic response surfaces. This study has also concentrated on quadratic approximations for several reasons:

1. linear approximations, in most instances, will be inadequate to model functions of interest,
2. for many problems, a 2nd order approximation will be adequate to model response especially if the region of interest is limited,
3. there is a scarcity of literature which address design selection for cubic or higher order polynomial approximations, and
4. in optimization process using response surfaces, for moderate size problems, it is more computationally efficient to perform a sequence of quadratic approximations than one cubic or higher order approximation. This fact is next discussed.

The number of terms in a second order polynomial in n design variables is

$$\text{number terms quadratic} = (n+1) + \frac{n(n+1)}{2} \quad (100)$$

The number of terms in a 3rd order polynomial in n design variables is

$$\text{number terms cubic} = 1 + \frac{3}{2}n(n+1) + \frac{n!}{6(n-3)!} \quad (101)$$

Table 8.1 gives, for various number of design variables, the number of terms in a 2nd order and 3rd order polynomial and their ratio.

Table 8.1 Number of terms in a 2nd and 3rd order polynomial and their ratio

number of design variables, n	number of terms in quadratic	number of terms in cubic	cubic/quadratic
3	10	20	2
6	28	84	3
9	55	220	4
12	91	455	5
15	136	816	6

One can see that for problems with more than 6 design variables, it will probably be more computationally efficient in an optimization algorithm to utilize a sequence of quadratic response surfaces than one 3rd or higher order response surface. When there are 6 or fewer design variables, 3rd or 4th order response surfaces may be used to advantage.

In this report, the term "minimum point design" refers to a design that has just enough design points to allow the determination of coefficients of an approximating polynomial. The term "augmented minimum point design" is a minimum point design which contains

additional design points. Thus, augmented minimum point designs are over-determined designs. The studies that have been performed in this report indicate that augmented minimum point designs are competitive with, if not better than, central composite designs for developing a 2nd order response surface. A program DESIGNS [20] was developed for generating augmented minimum point designs for developing a 2nd order response surface. That program is described in Section 8.2.

When there are 6 or fewer design variables, it may be computationally beneficial to use a 3rd order or 4th order response surface. Thus, the program DESIGN4 [28] was developed to generate augmented minimum point designs for a 4th order response surface. The program DESIGN4 is discussed in Section 8.3. The program can also be used to develop a 3rd order response surface. The 3rd order minimum point design is a subset of the 4th order minimum point design. Thus the 4th order minimum point design will give an over-determined 3rd order approximation. Additional randomly selected design points can be added to the 4th order minimum point design to give the desired degree that the 3rd order approximation is to be over-determined.

8.2 Augmented Minimum Point Designs for 2nd Order Approximations

The basic building block for program DESIGNS is the star pattern of design points. Figure 4 shows the star pattern for 3 design variables. This pattern of design points allows one to determine those coefficients of a 2nd order polynomial approximation associated with the

terms

$$1, x_i, x_i^2, \quad i=1,n \quad (102)$$

To be able to determine the coefficients associated with the terms

$$x_i x_j, \quad i \neq j \quad (103)$$

one must supplement the star pattern with one additional design point in the x_i, x_j planes. Figure 30 shows the additional design point in the x_i, x_j plane. Figure 6 shows the total minimum point design for 3 design variables.

Studies of this report indicate that designs should be over-determined. Having a design that is 20-50% over-determined is a good compromise between keeping down the number of design points while still getting a good approximation. The program DESIGNS augments the minimum point design with a user selected number of random design points.

8.2.1 Specifics of program DESIGNS

A listing of the FORTRAN program DESIGNS is found in Appendix 1 and a copy of that program is found in file "designs.f" on the floppy disk accompanying this report. The program should be compiled with a F77 compiler with the compiled program called "design". To run the program just enter "design" from the keyboard. The program prompts the user for

1. the number of design variables,
2. the number of designs points to augment the minimum point design, and
3. a seed parameter, IFLAG, which is used to generate the random numbers (IFLAG can be entered as any positive integer).

The program then generates a design in local coordinates with the maximum range on each design variable of -1 to +1. The program then

4. asks the user to enter an integer which specifies whether design point coordinates are to be also generated in global coordinates. If they are to be calculated in global coordinates, the program then
5. prompts the user to enter the range of design variables in global coordinates.

Results with commentary are written to file "design.res". Design points without commentary are written to file "design.run".

8.3 Augmented Minimum Point Design for 3rd and 4th Order Approximation

A 3^k factorial design is used as the building block of this minimum point design. The 3^k factorial design provides information for calculating the coefficients associated with the terms

$$1, x_i, x_i x_j, x_i^2, x_i^2 x_j, x_i^2 x_j^2, \quad j \neq i \quad (104)$$

Additional points are then added at -1 and 1 (in local coordinates) along the x_i axis. These

points together with the 3^k factorial design point give the star pattern which can be seen in Figure 31. With this arrangement of points, there are 5 design points along the x_i axis which provides information for calculating the coefficient associated with the terms

$$x_i^4 \quad (105)$$

Additional design points are then placed in each x_i, x_j plane which provides information for calculating the coefficient associated with the terms

$$x_i^3 x_j \quad (106)$$

These points are also shown in Figure 31.

8.3.1 Specifics of program DESIGN4

A listing of the FORTRAN program DESIGN4 is found in Appendix 2 and a copy of that program is found in file "design4.f" on the floppy disk accompanying this report. The program should be compiled with a F77 compiler with the compiled program called "design4". To run the program just enter "design4" from the keyboard. The program

prompts the user for needed information. Prompts and response are similar to those for the program DESIGNS.

8.4 Conclusion

A minimum point design is a design that has just enough design points to allow the determination of the coefficients of an approximating polynomial. An augmented minimum point design is a minimum point design which contains additional design points. Augmented minimum point designs are competitive with, if not better than, central composite designs for developing a 2nd order response surface. Minimum point designs should be augmented with enough points that the approximation is 20-50% over-determined. A program DESIGNS was developed for generating augmented minimum point designs for developing a 2nd order response surface.

When there are more than 6 design variables, 3rd or higher order approximations require so many design points that it is computationally better to perform a sequence of 2nd order approximations in an optimization process than one higher order approximation. When there are 6 or fewer design variables, a 2nd order approximation may often be satisfactory. However, for those cases where it is desirable to use a higher order approximation, program DESIGN4 was developed. That program generates designs which can be used to develop 3rd or 4th order approximations.

9. Solution Algorithm

9.1 Introduction

In this investigation, the program NEWPSI was used to perform the studies involving polynomial approximations. That program can investigate under-determined, exactly-determined, or over-determined approximations of various orders. The version submitted with this report can handle up to 15 design variables as programmed. The order of polynomial it can handle is as follows:

1. one design variable, up to a 20th order polynomial
2. two design variables, up to a 5th order polynomial
3. for 2-15 design variables, a second order polynomials.

One can use up to 250 designs to train the approximation. In calculating v_G , it can handle up to 2000 grid points.

The program solves for the undetermined parameters associated with the approximation. It then evaluates the approximate function at the design points and calculates the error parameter, v . It then reads in the design points and function value on the test grid. The approximate function is evaluated at the grid points and the error parameter, v_G , is then evaluated.

9.2 Program Specifics

A listing of the FORTRAN program NEWPSI is found in Appendix 3 and a copy of that

program is found in file "newpsi.f" on the floppy disk accompanying this report. The program should be compiled with a F77 compiler and the compiled program called "newpsi". To run the program just enter "newpsi" from the keyboard. Data is read from the file "newpsi.dat". Data can be in free format. The program asks for the following data:

1. a value of the print code, ip; (If ip=4, great quantities of output are generated for program checkout. Normally the program is run with ip=0 for normal output).
2. the number of design variable, nd;
3. the order of the polynomial being considered, np;
4. the number of design points in the design, m;
5. the design and function value at the design points, x(i,j), y(i);
6. the number of design points on the grid, ng; and
7. the design and function value at the grid points, xx(i,j), yy(j).

Output is written to the screen and to file "newpsi.res".

10. Conclusion

For a given order of approximation of a function, f , the quality of the approximation is affected by

- a. the number of levels of the design variables,
- b. the location of the design points, and
- c. the degree which the approximation is over-determined.

For an n th order approximation,

1. there must be $n+1$ levels of the design variables;
2. the design points must be located so that information is available for calculating all of the n th derivatives of f ;
3. the approximation should be, at least, 20-50% over-determined.

For example, for a 2nd order approximation in 3 design variables, there must be at least 3 levels of the design variables, design points must be located so that information is available for calculating

$$\frac{\partial f}{\partial x_i}, \frac{\partial^2 f}{\partial x_i \partial x_j}, \quad i=1,3; j=1,3 \quad (107)$$

A complete 2nd order polynomial approximation contains 10 undetermined coefficients.

Thus, at least 10 design points are required to provide information for calculating these

coefficients. To have the approximation 30% over-determined, one would want to use 13 design points.

For second order approximations, when there are fewer than 6 design variables, central composite designs meet requirements 1-3. However, for 6 or more design variables, these designs contain too many design points. A minimum point design is one which contains just enough design points, meeting the derivative requirements of item 1 and 2 above, to exactly-determine the approximation. An augmented minimum point design is a minimum point design supplemented with additional design points. The program DESIGNS was developed to yield augmented minimum point designs for 2nd order approximations. The quality of approximations developed using designs from program DESIGNS was comparable to, if not better than, other standard designs such as the central composite designs.

For more than 6 design variables, 3rd and 4th order approximations require so many design points to determine the coefficients in those approximations that it is more computationally efficient to develop a number of 2nd order approximations than one approximation of 3rd or higher order. For 6 or fewer design points, 2nd order approximations may be quite adequate. However, for those cases where one wishes to use a 3rd or 4th order approximation, the program DESIGN4 was developed. That program generates an augmented minimum point design for developing a 4th order approximation.

Previous studies have shown that the quality of approximations using neural networks is

comparable to those using polynomial approximations when the number of undetermined parameters associated with the approximations is the same. Thus, neural networks trained with designs from DESIGNS or DESIGN4 should offer approximations of comparable quality to those obtained using polynomial approximations with the same number of associated undetermined parameters.

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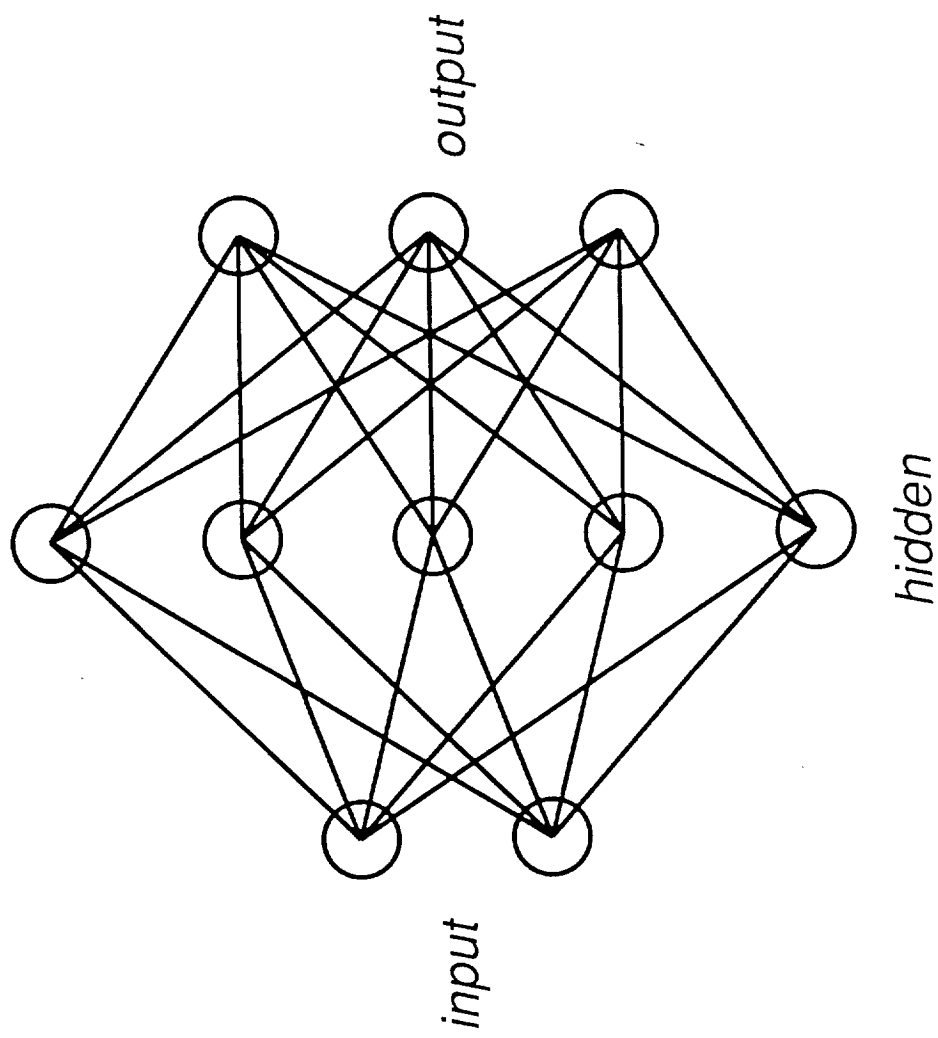


Figure 1. Artificial neural net

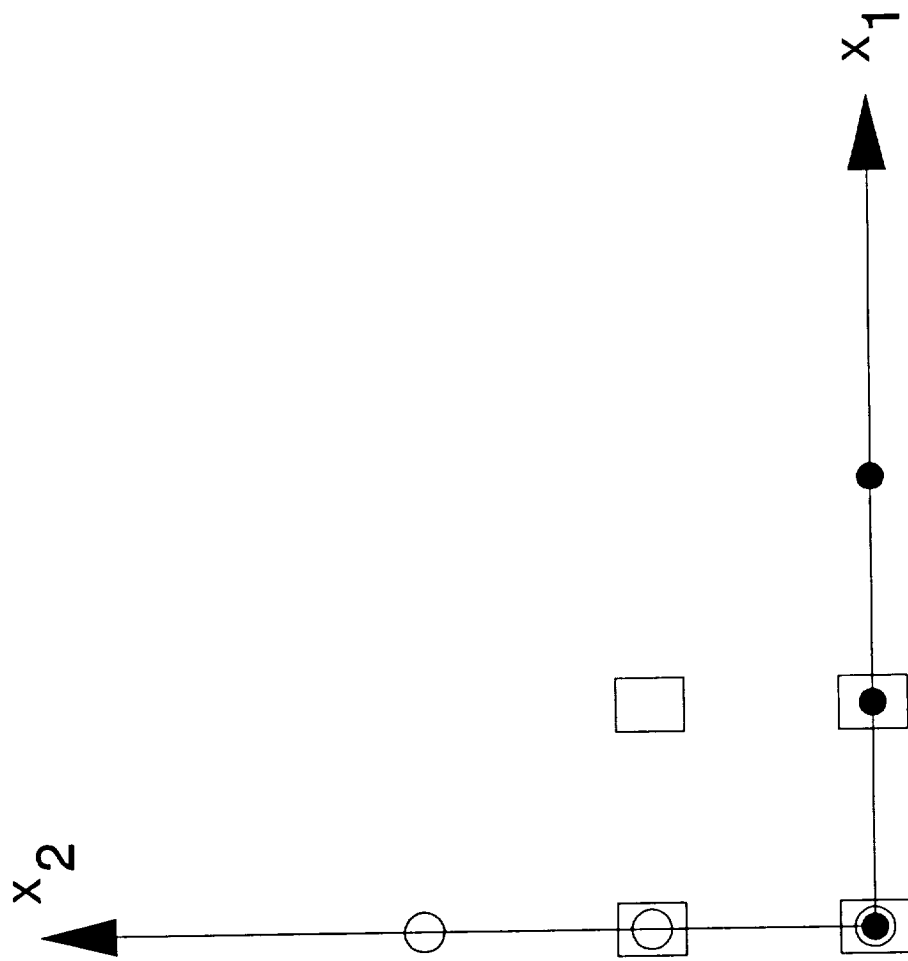


Figure 2. Complete design

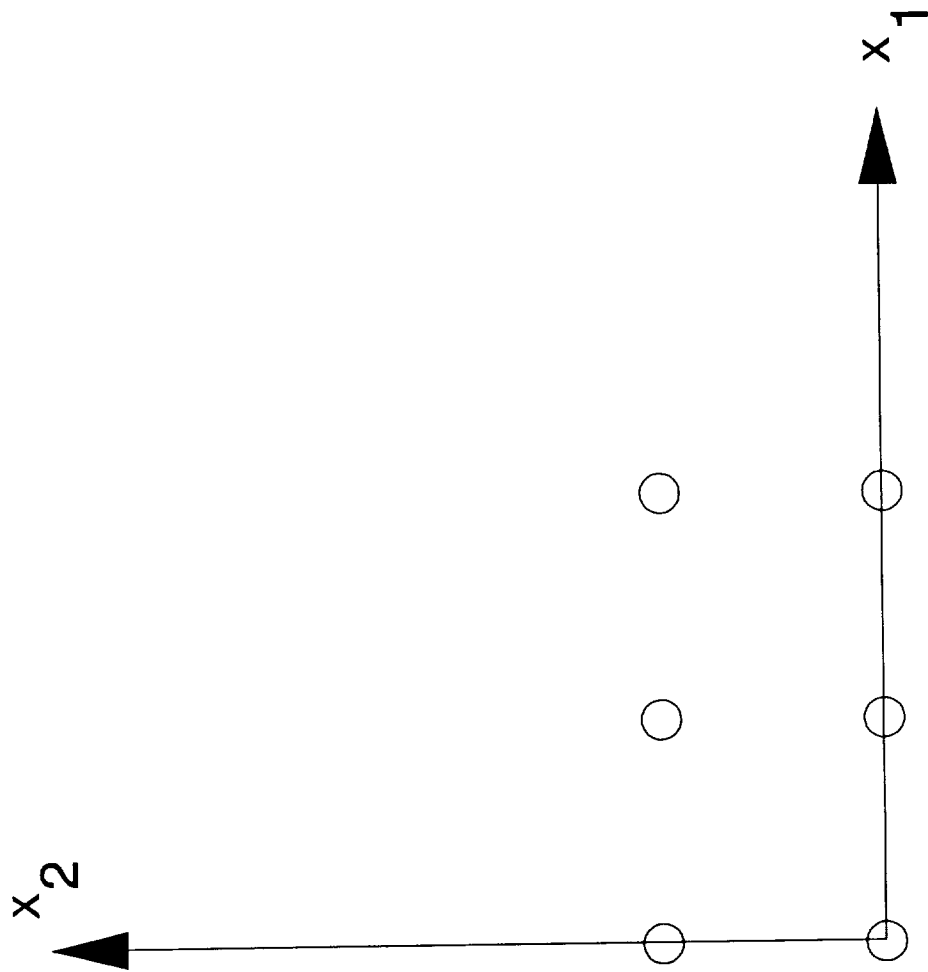


Figure 3. Deficient design

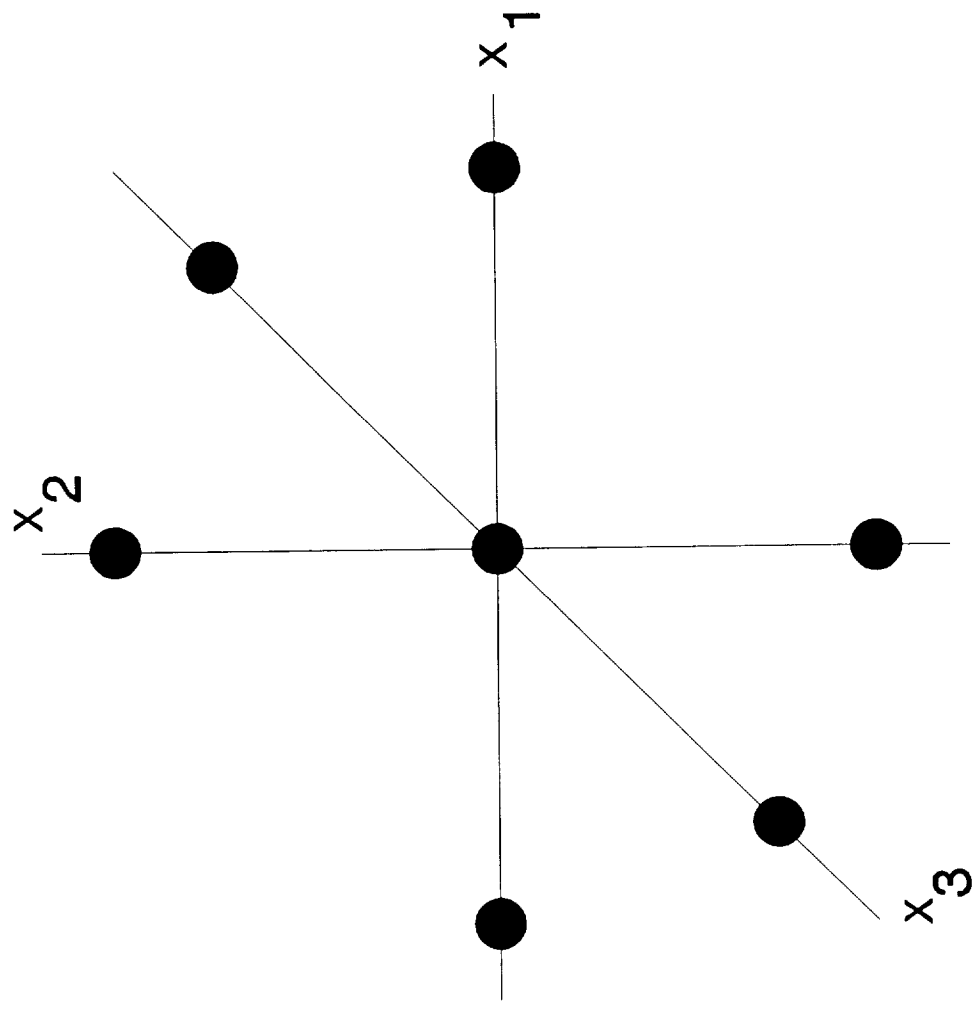


Figure 4. Star pattern of design points--7 design points

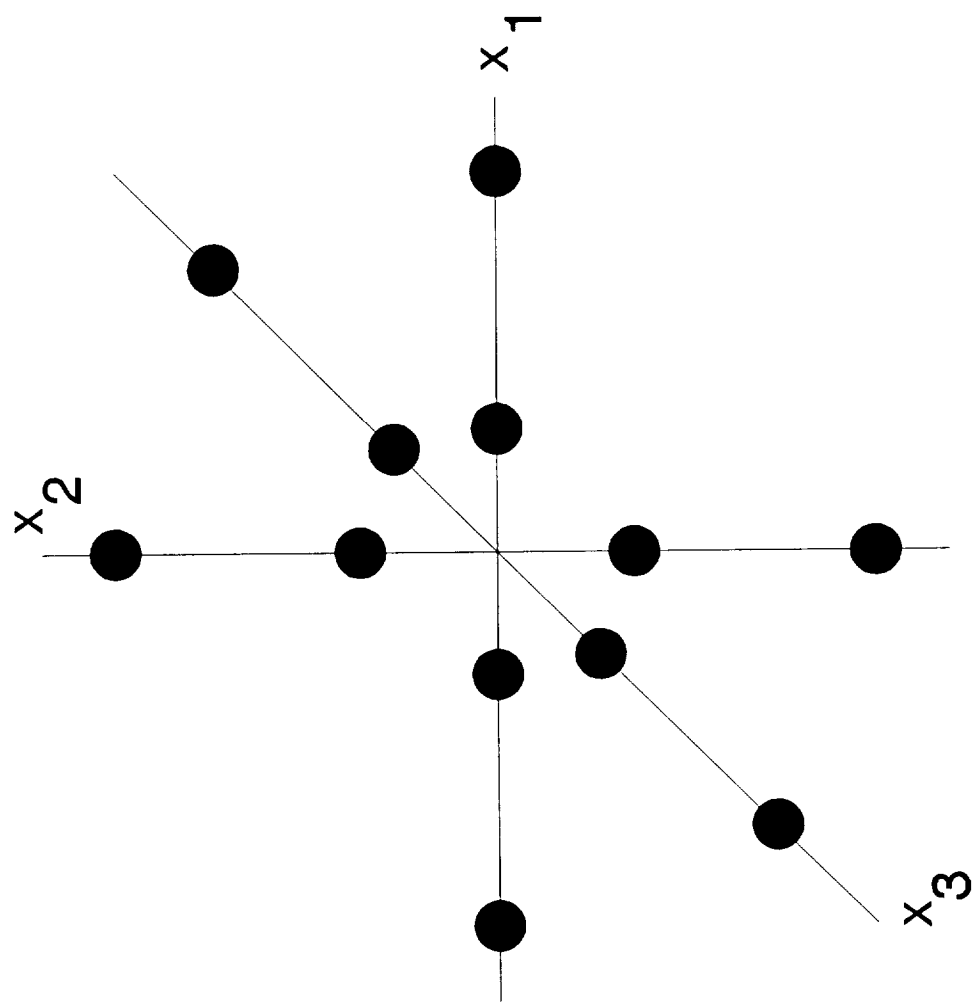


Figure 5. Star pattern of design points--12 design points

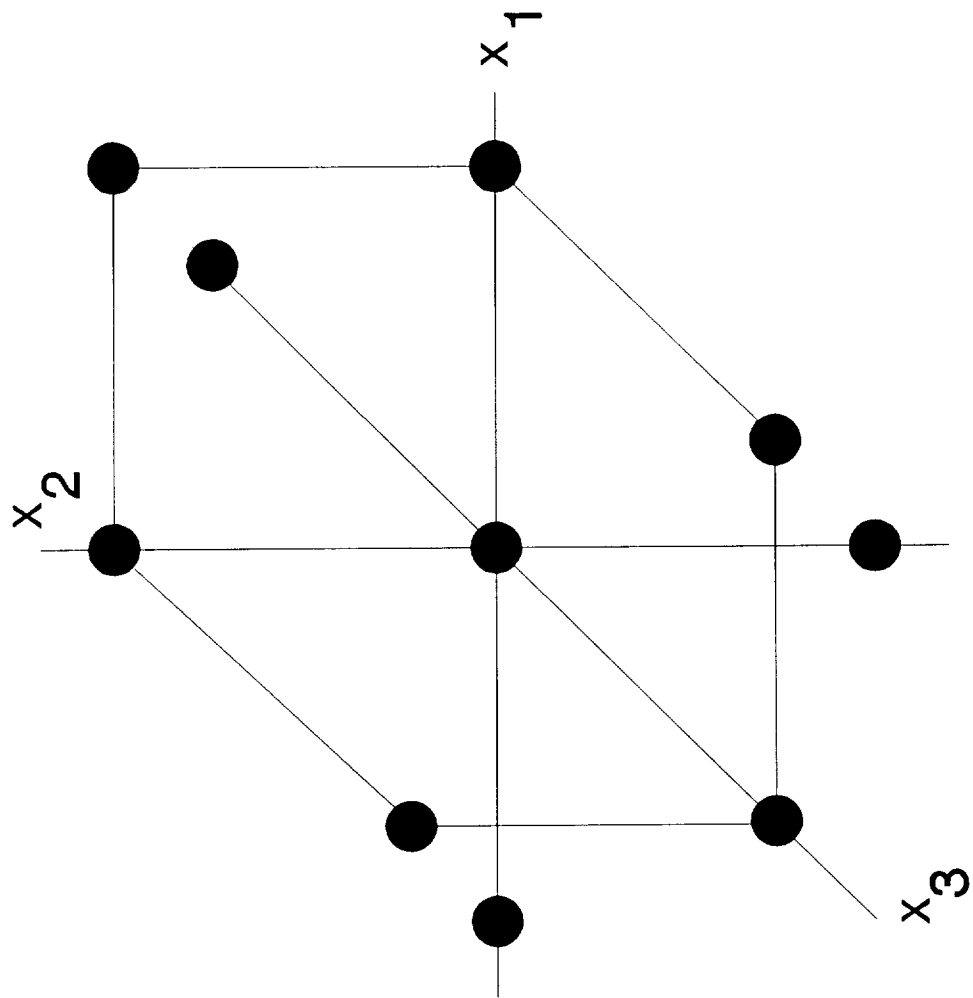


Figure 6. Design from program DESIGNS

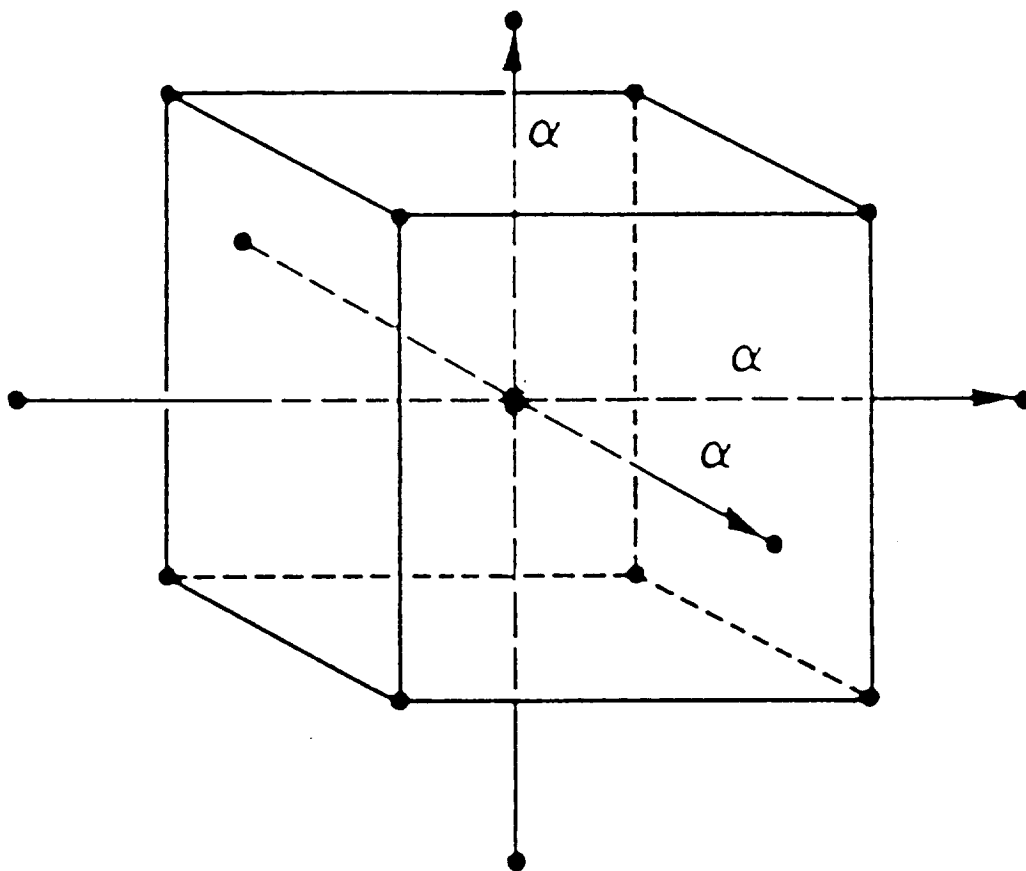


Figure 7. Central composite design for $k=3$

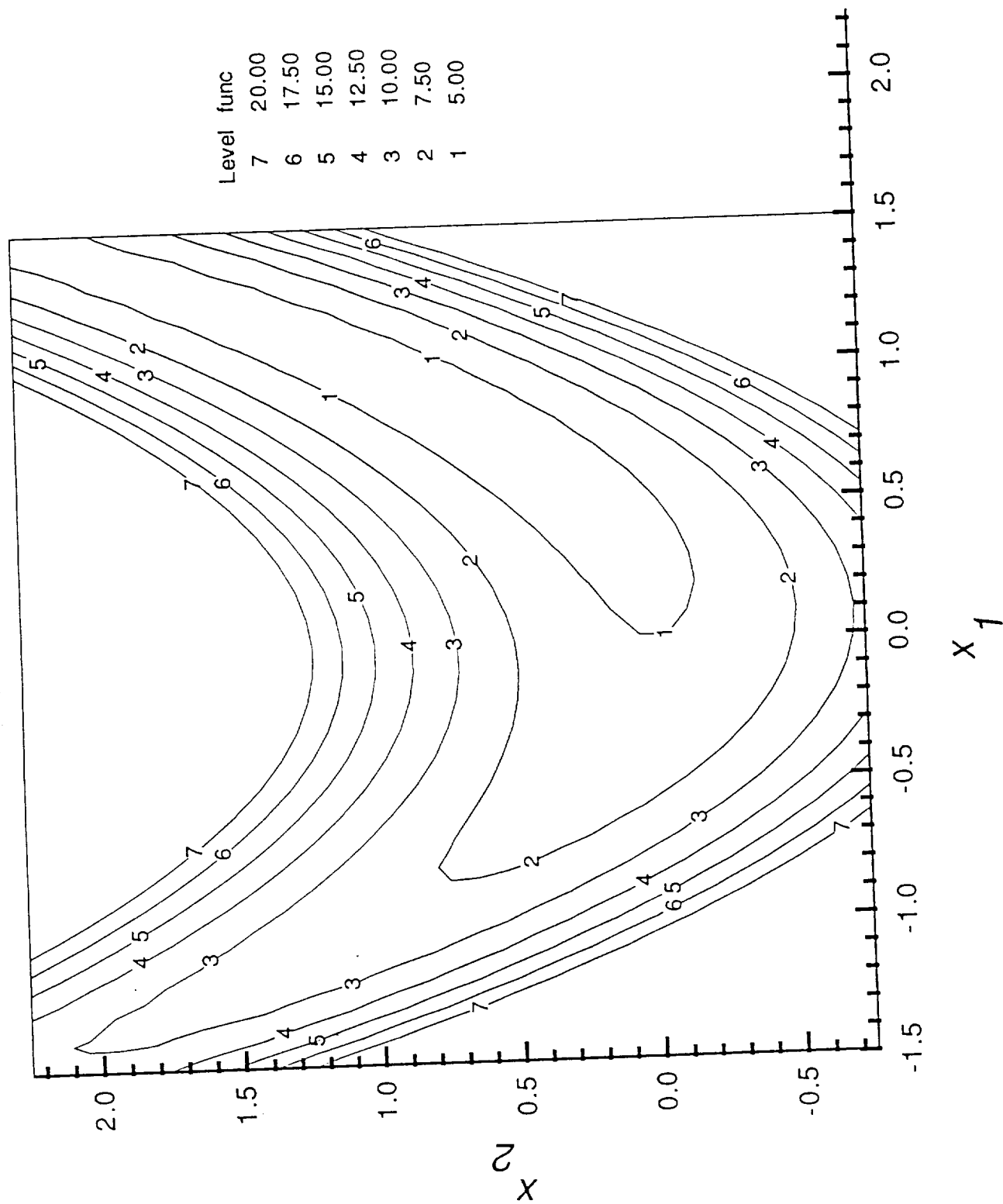


Figure 8. Fox's banana function

One Dimensional Example

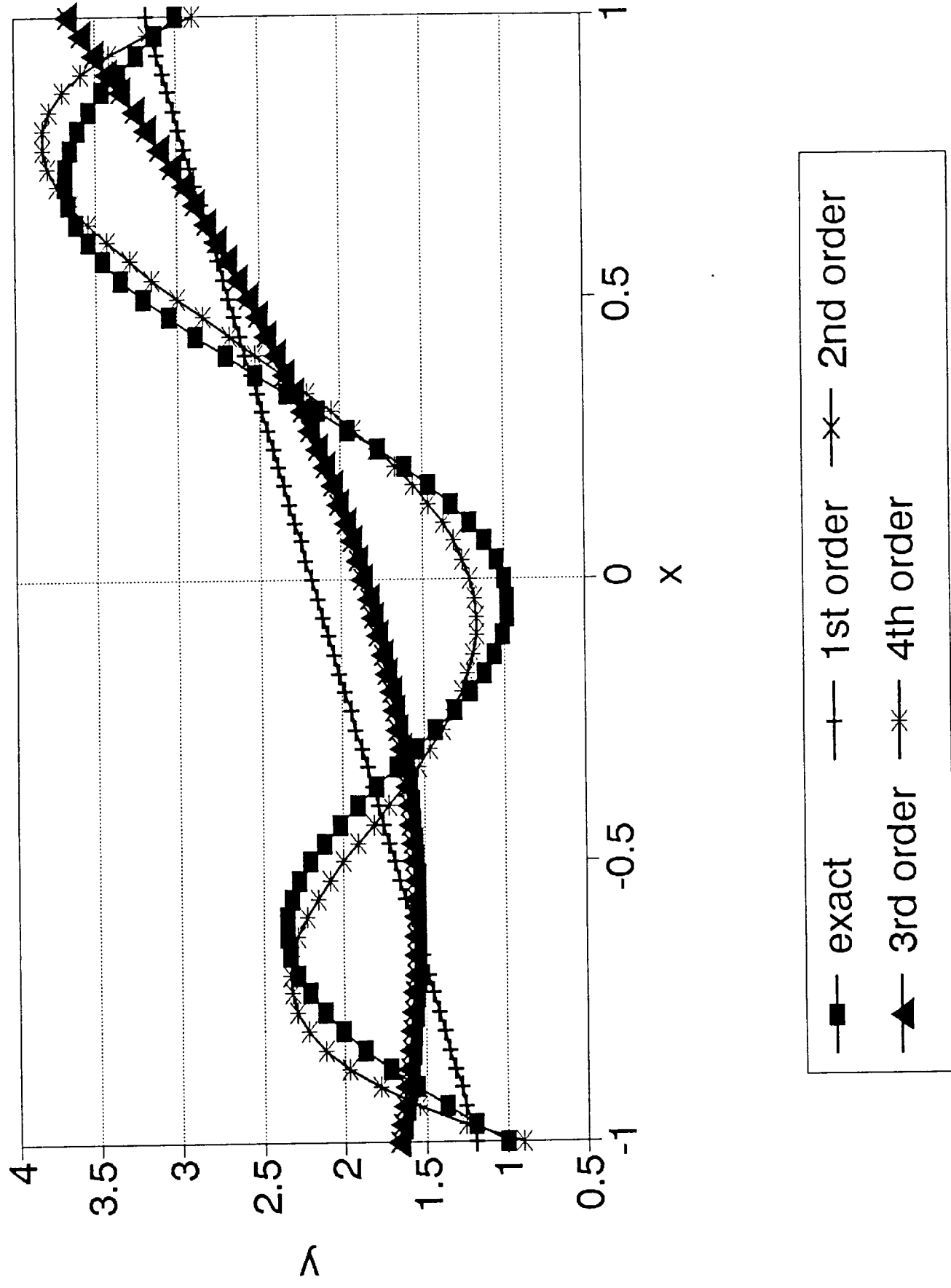


Figure 9. One dimensional example

Parameter vG

First Order Polynomial Approximation

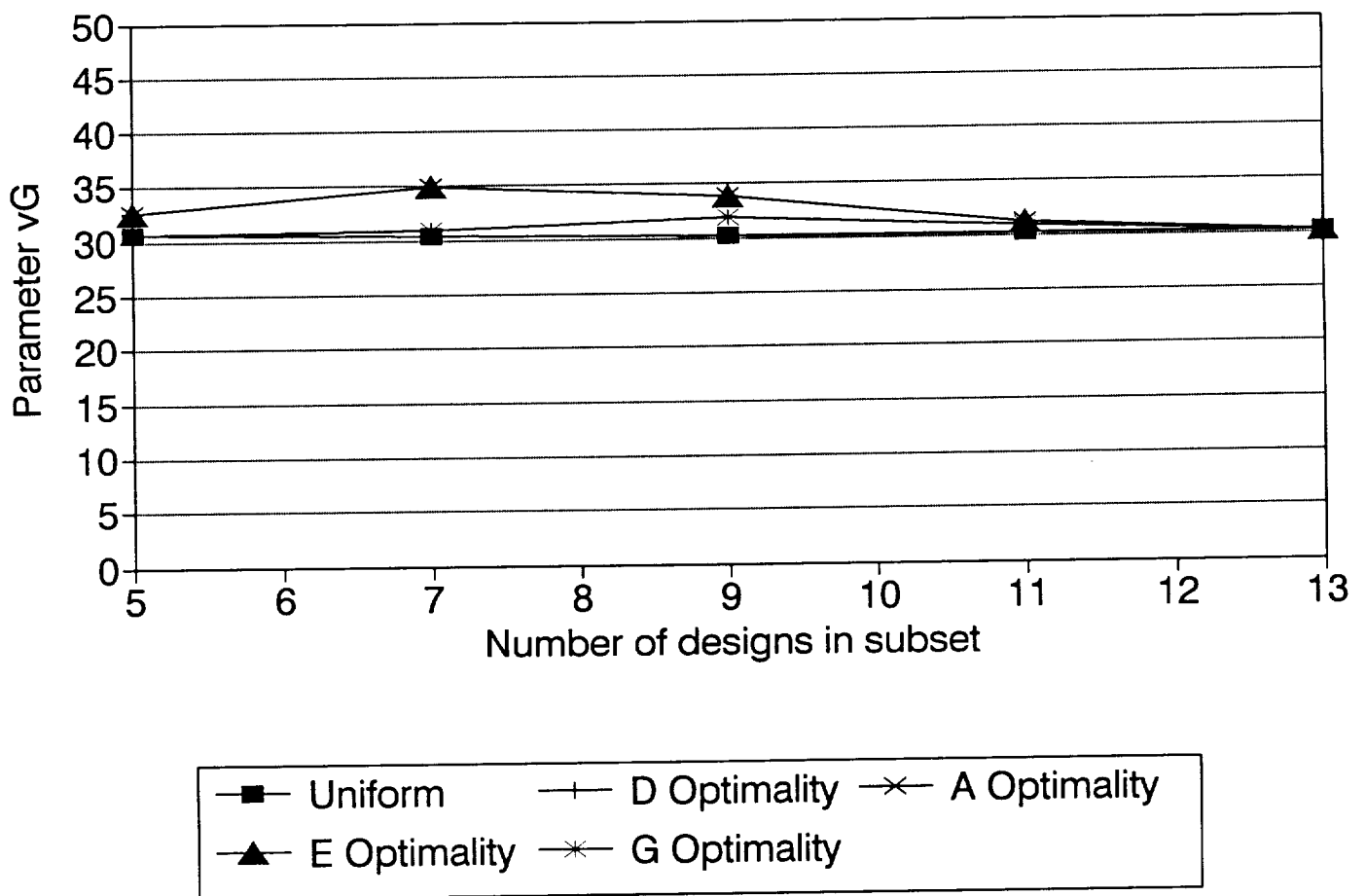


Figure 10. D, A, E, and G optimality, first order approximation

Parameter vG

Second Order Polynomial Approximation

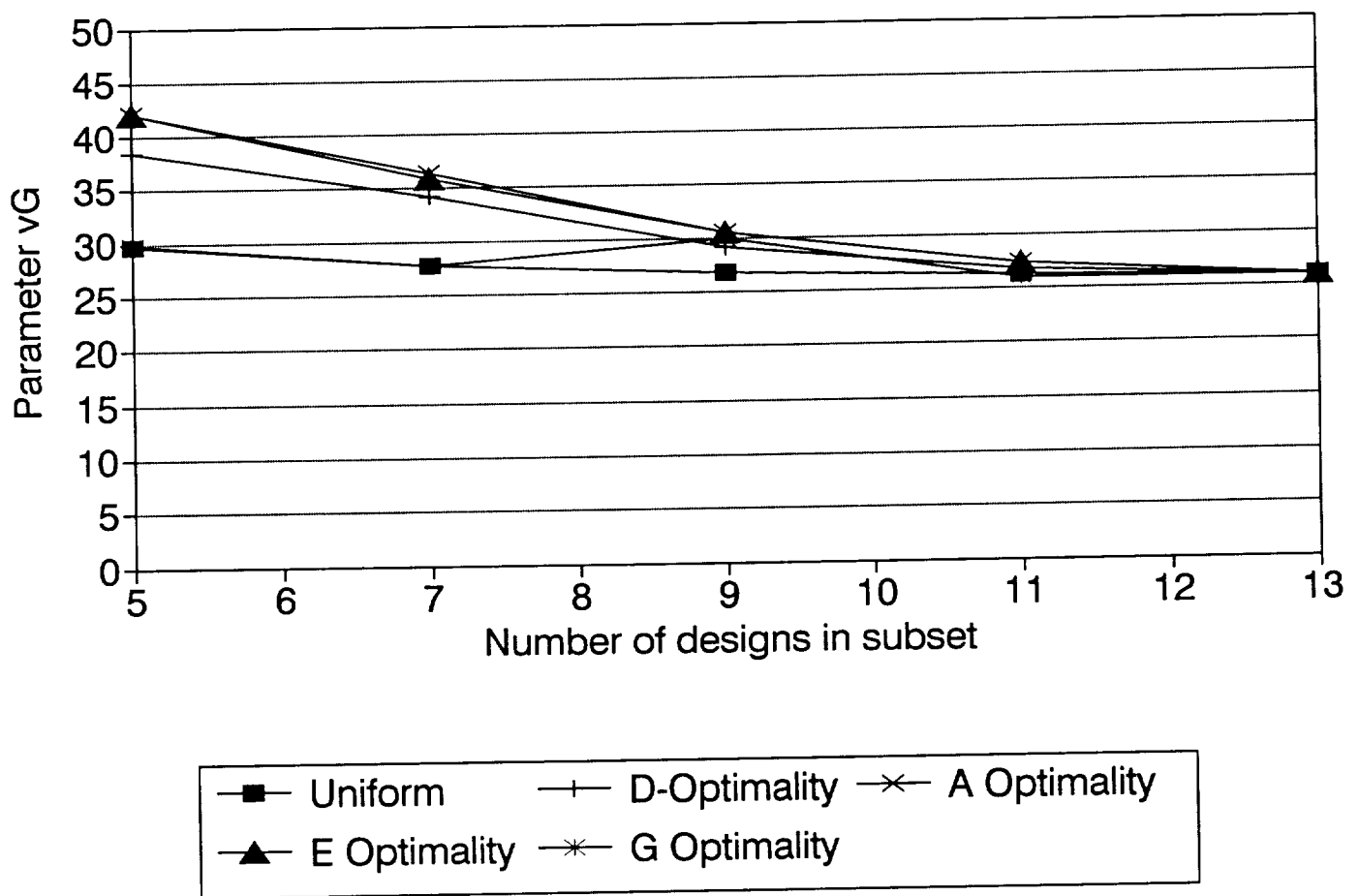


Figure 11. D, A, E, and G optimality, second order approximation

Parameter vG

Third Order Polynomial Approximation

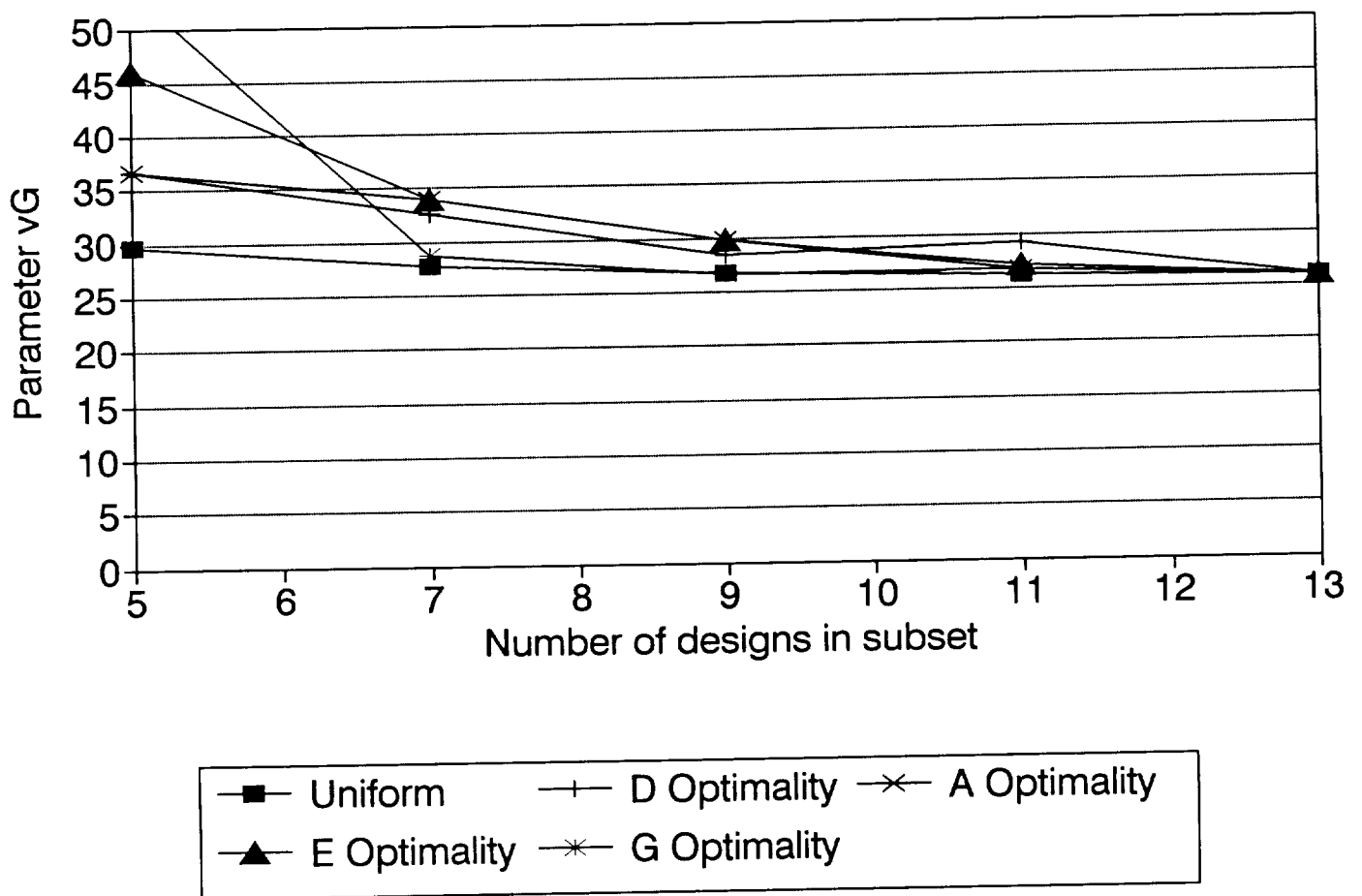


Figure 12. D, A, E, and G optimality, third order approximation

Parameter vG

Fourth Order Polynomial Approximation

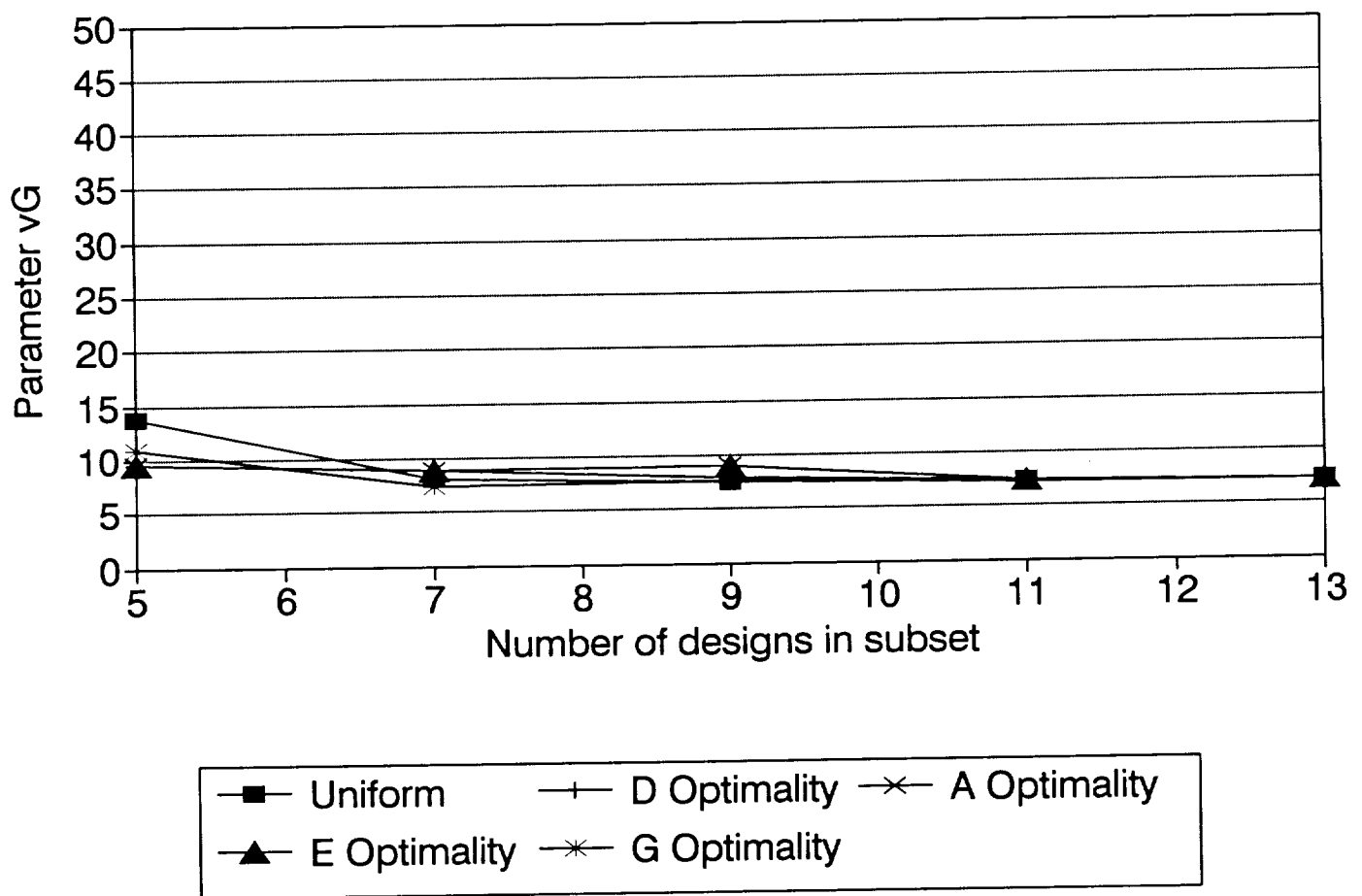


Figure 13. D, A, E, and G optimality, fourth order approximation

Parameter vG

First Order Polynomial Approximation

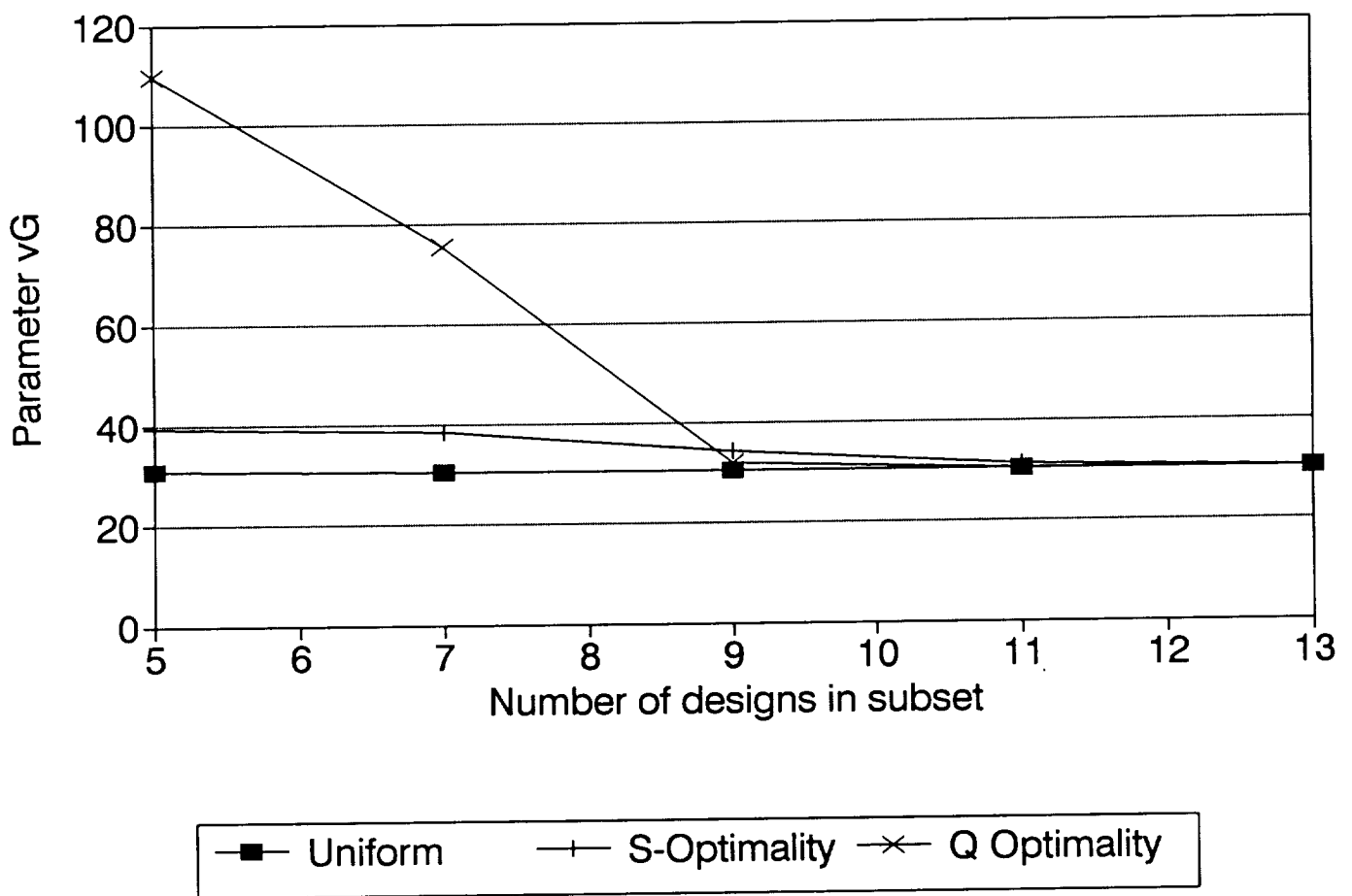


Figure 14. S and Q optimality, first order approximation

Parameter vG

Second Order Polynomial Approximation

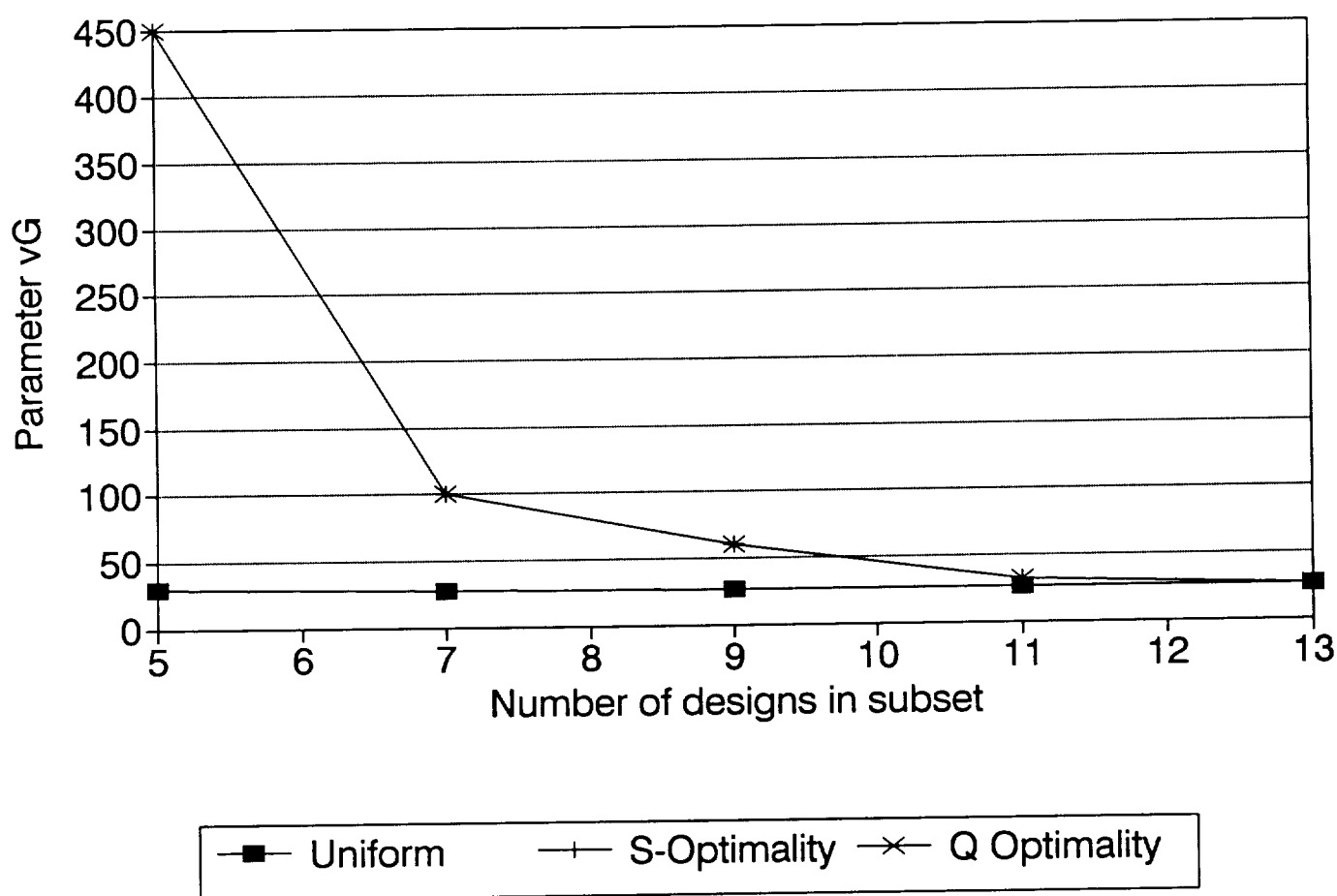


Figure 15. S and Q optimality, second order approximation

Parameter vG

Third Order Polynomial Approximation

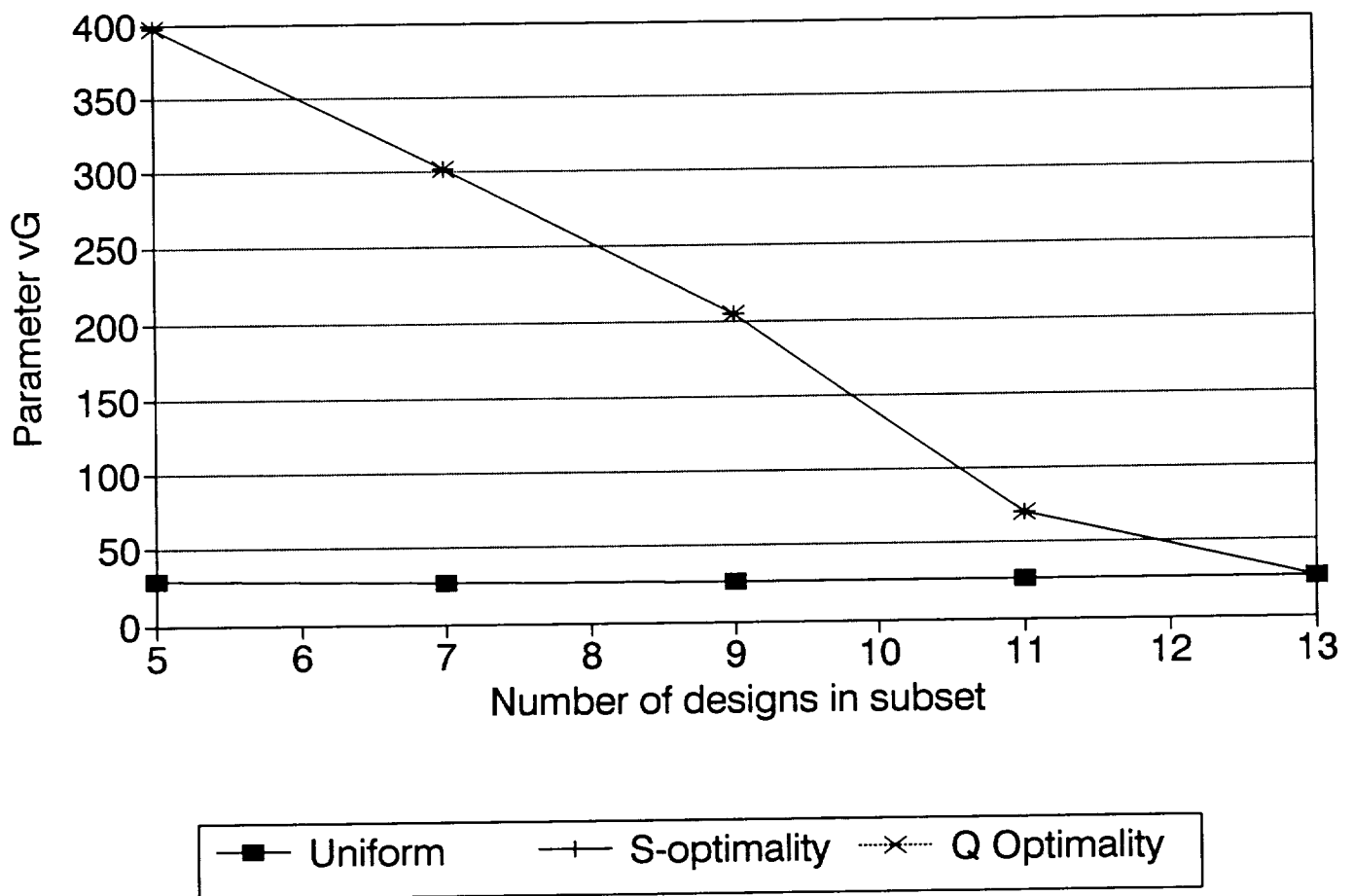


Figure 16. S and Q optimality, third order approximation

Parameter vG

Fourth Order Polynomial Approximation

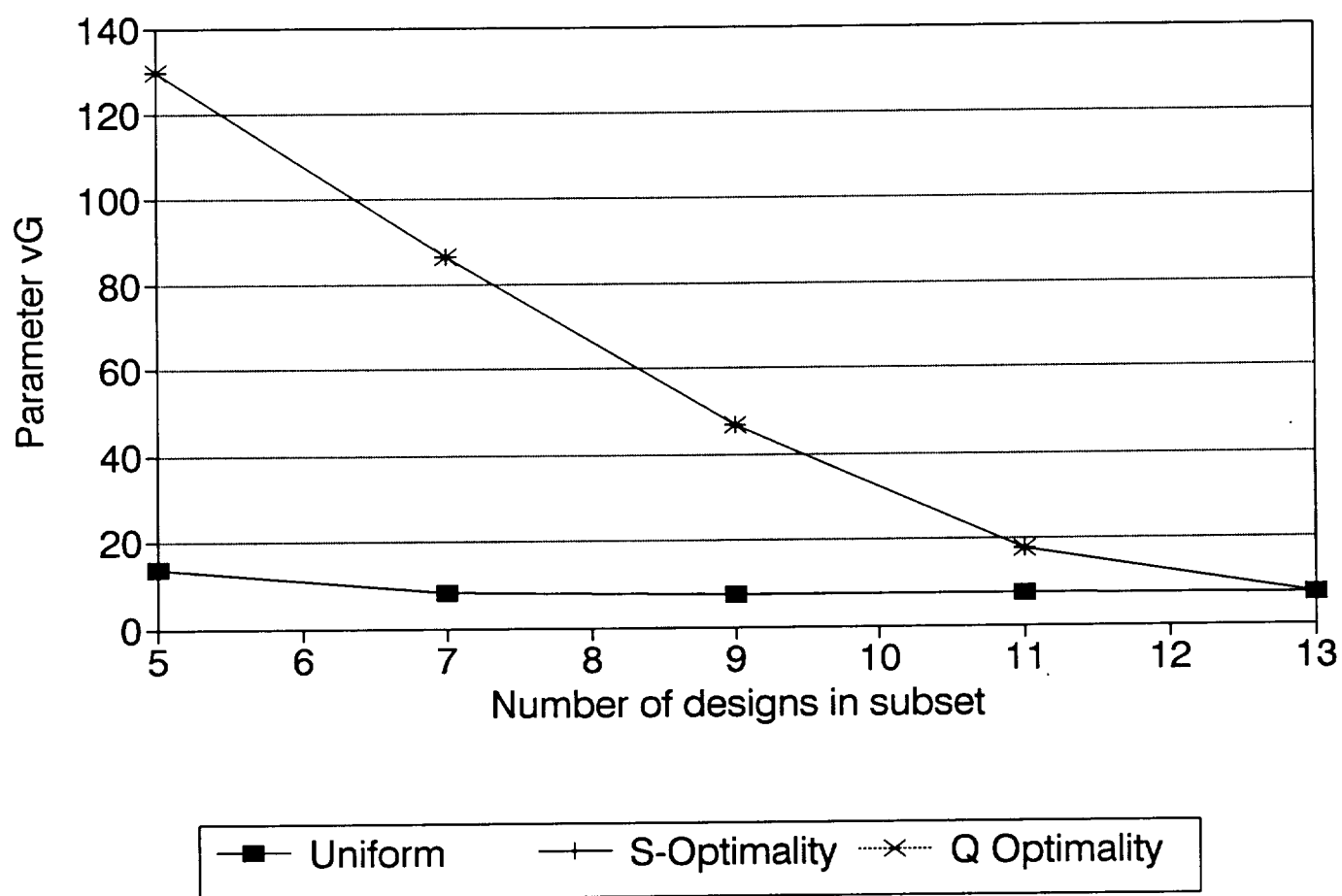


Figure 17. S and Q optimality, fourth order approximation

Y and its Approximation

Q optimality, 11 points out of 13

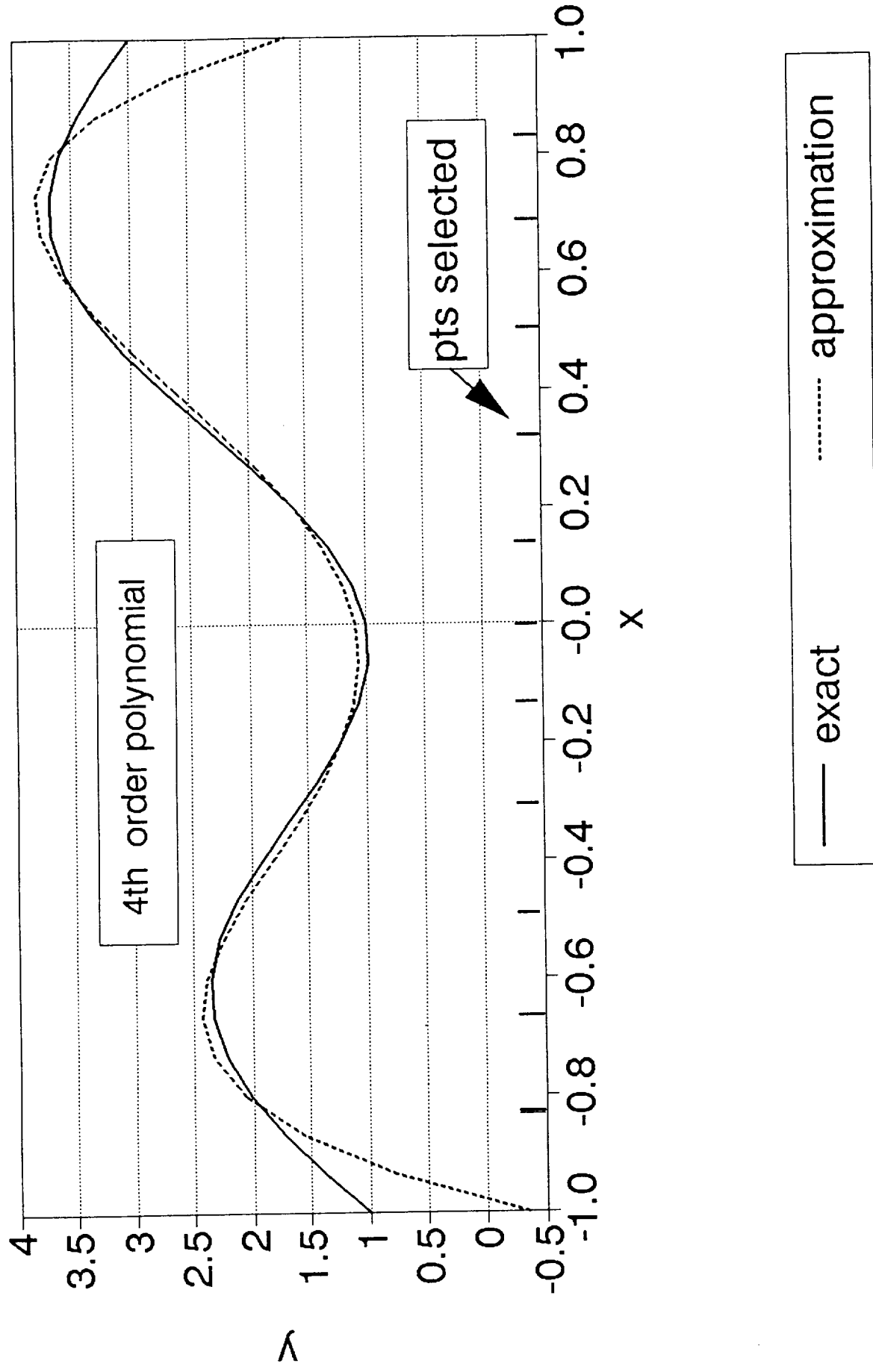


Figure 18. Q optimality, 11 out of 13 points selected

Y and its Approximation

Q optimality, 7 out of 13 points

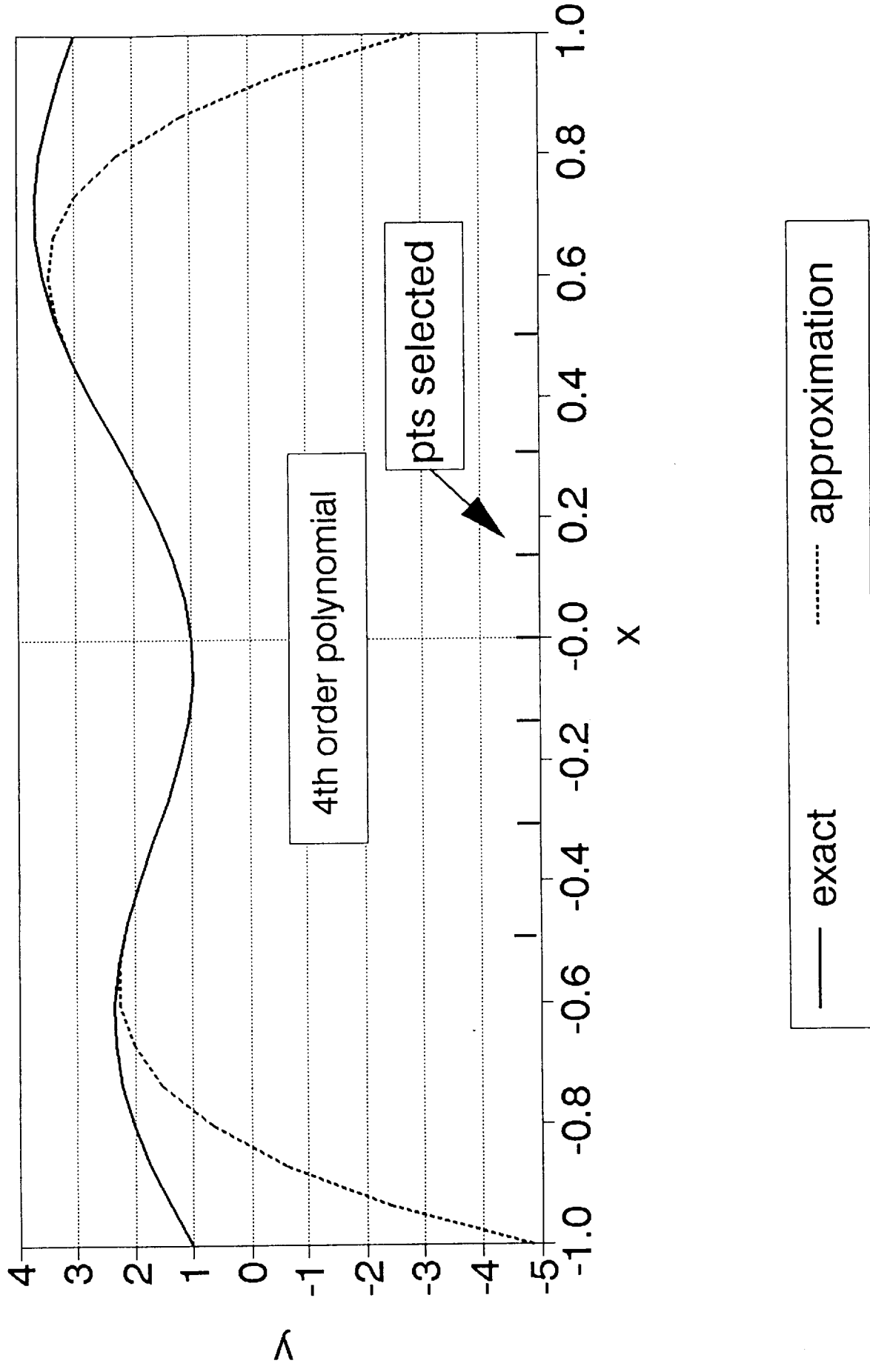


Figure 19. Q optimality, 7 out of 13 points selected

Y and its Approximation

Q optimality, 5 out of 13 points

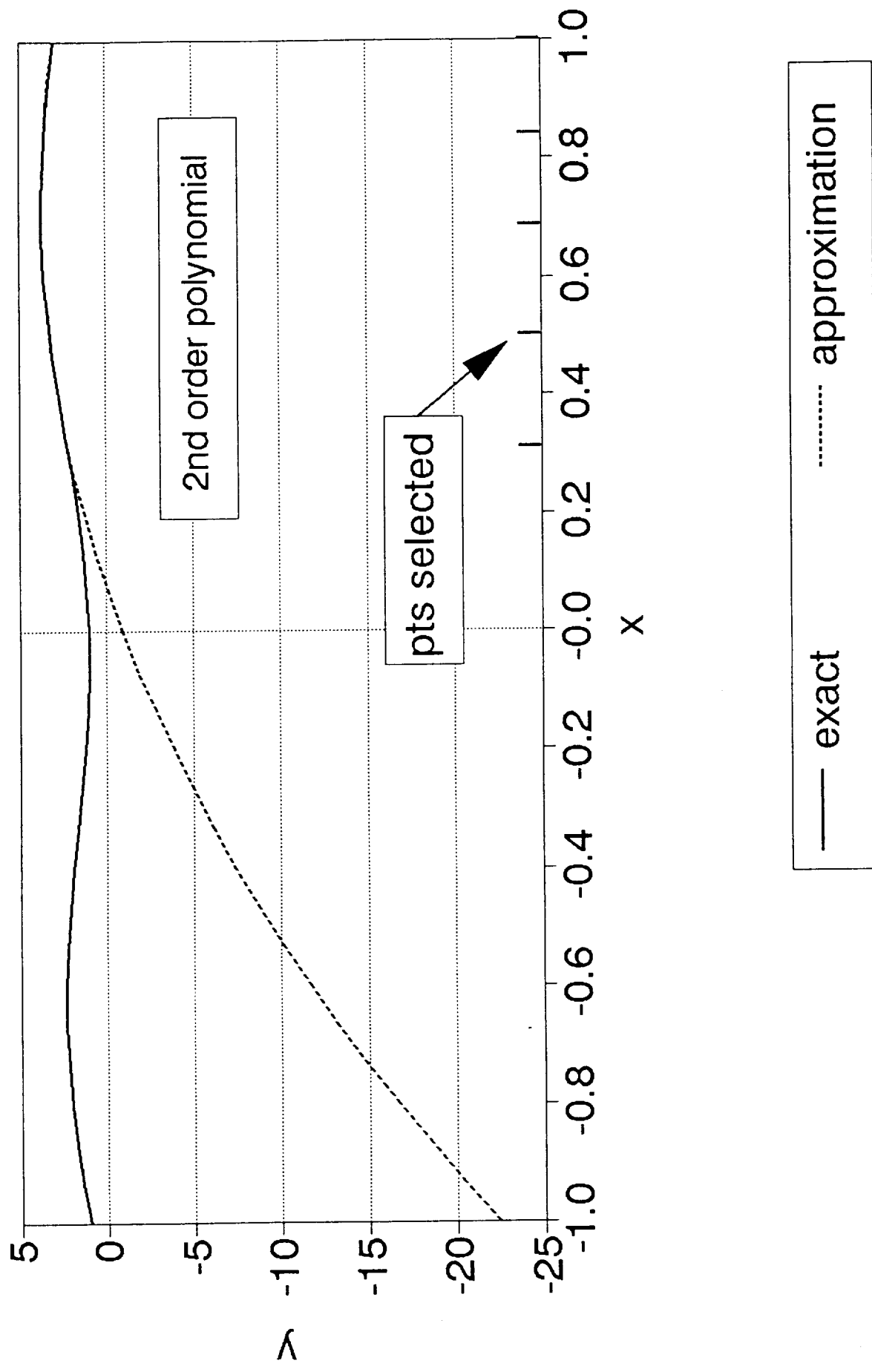


Figure 20. Q optimality, 5 out of 13 points selected

Parameter vG

First Order Polynomial Approximation

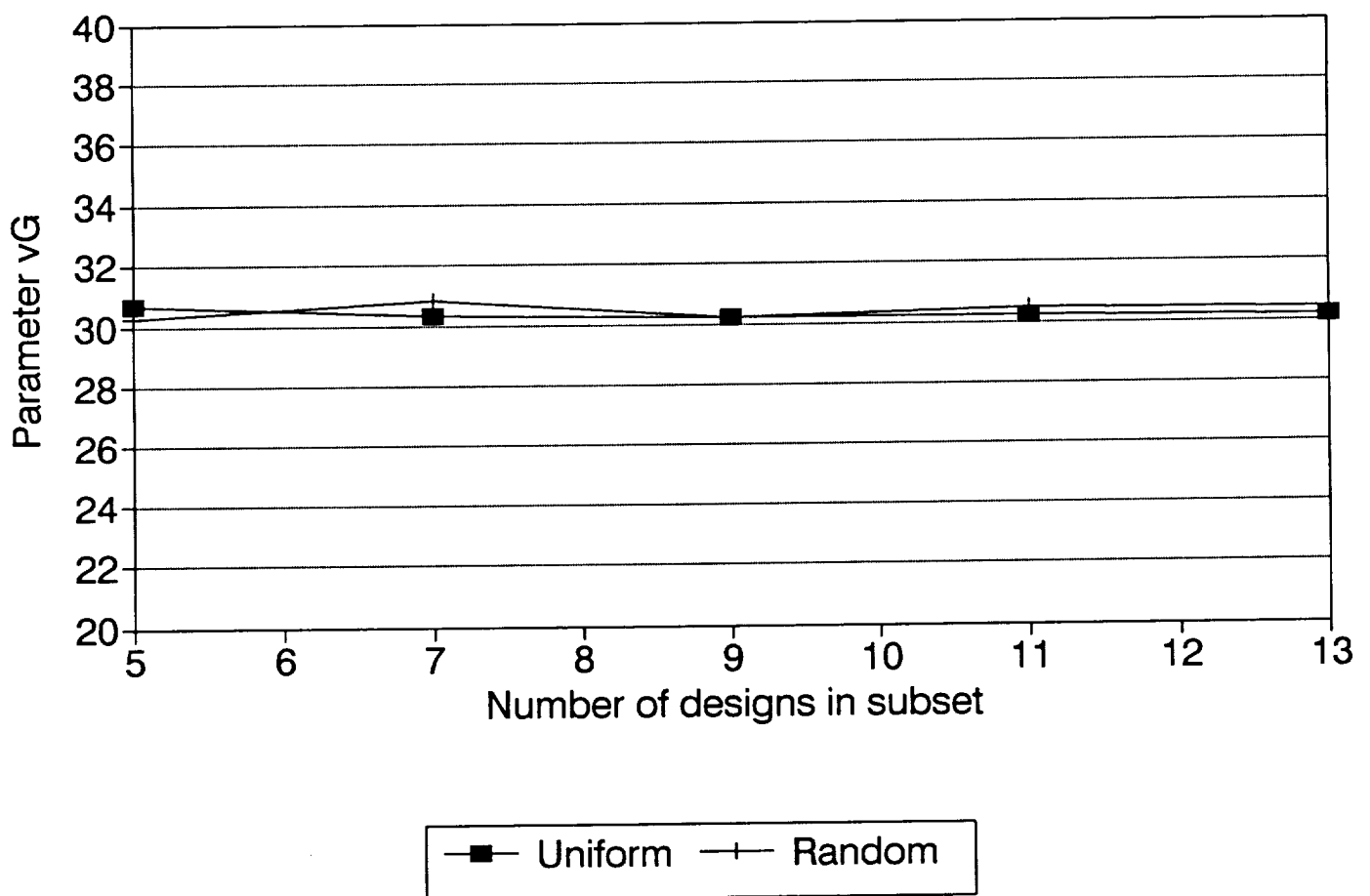


Figure 21. Random points, first order approximation

Parameter vG

Second Order Polynomial Approximation

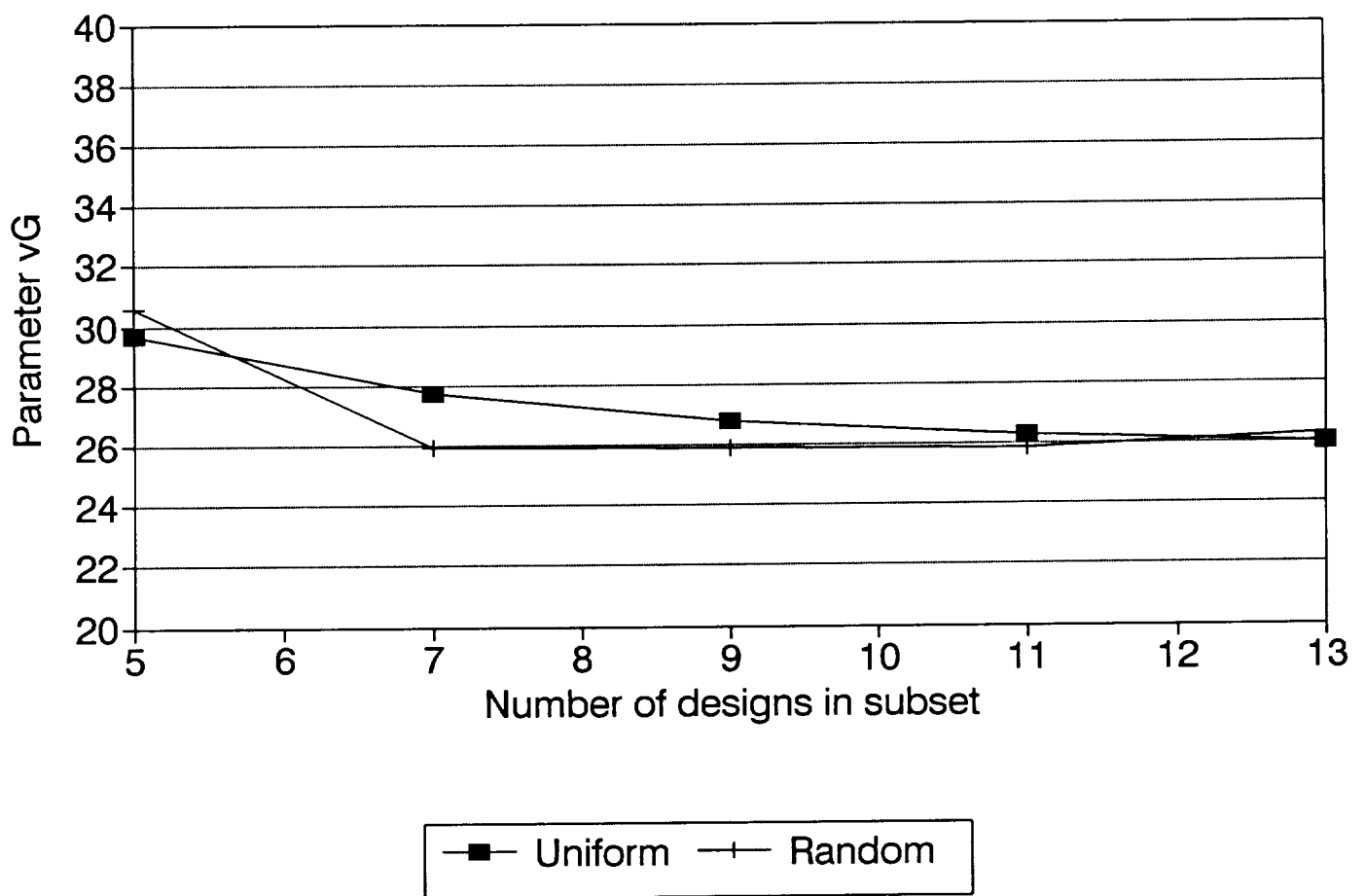


Figure 22. Random points, second order approximation

Parameter vG

Third Order Polynomial Approximation

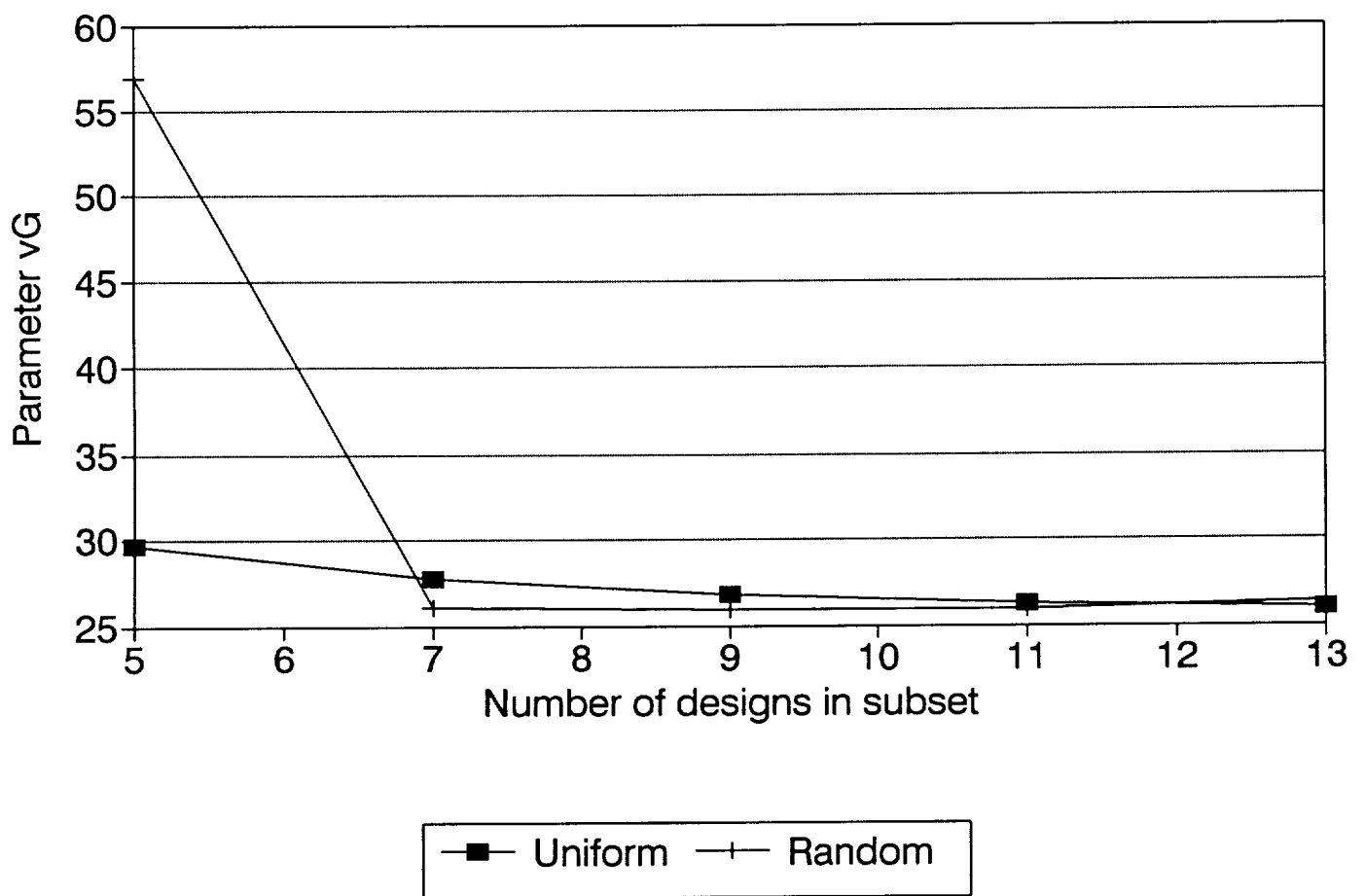


Figure 23. Random points, third order approximation

Parameter vG

Fourth Order Polynomial Approximation

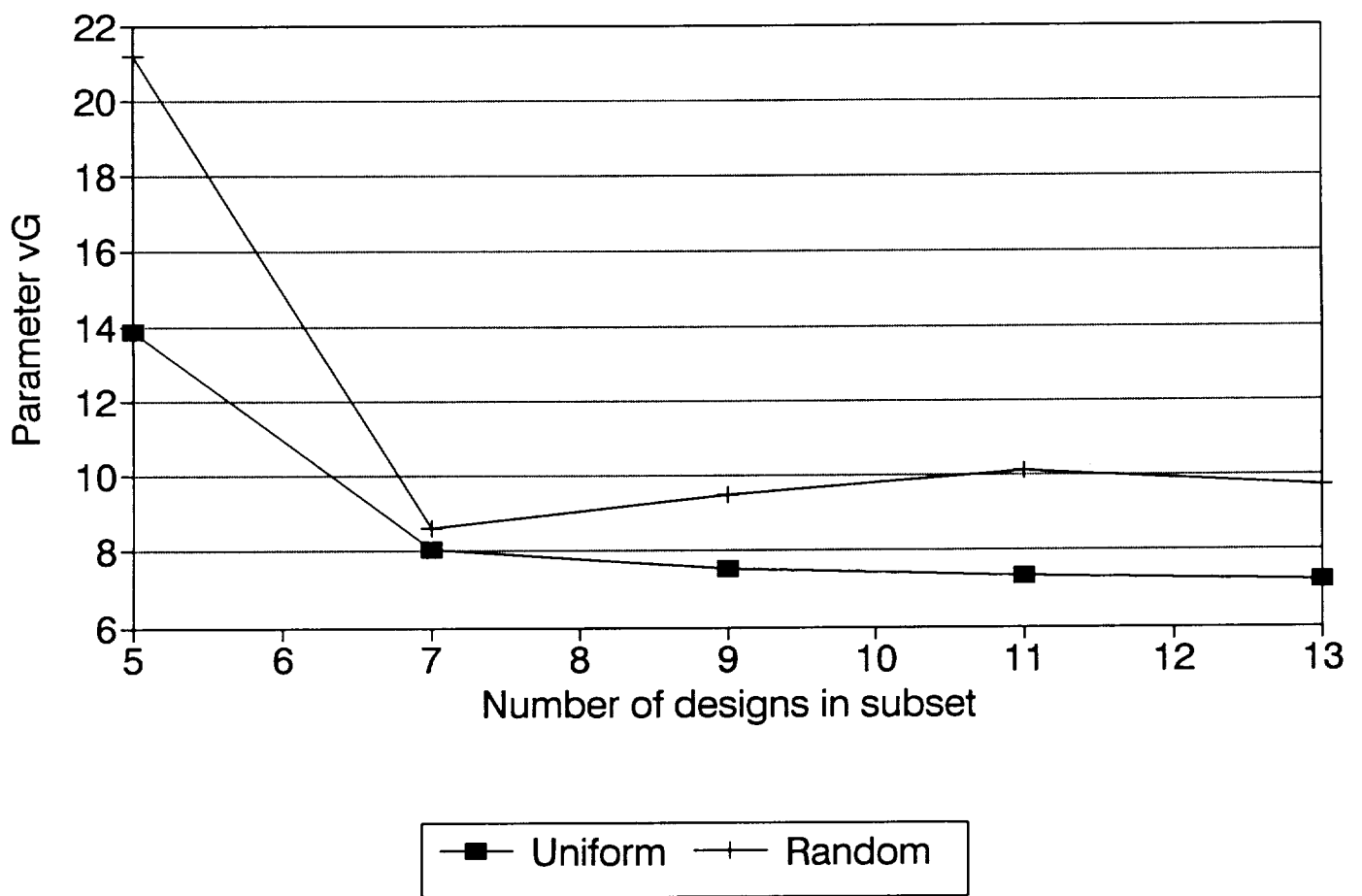


Figure 24. Random points, fourth order approximation

MEASUREMENTS OF QUALITY OF FIT
BEFORE AND AFTER t-test
PERFORMED ON

$$y = 10x_1^4 - 20x_2x_1^2 + 10x_2^2x_1^2 + x_1^2 - 2x_1 + 5$$

"Fox's Banana Function"

SECOND ORDER APPROXIMATION

$$Y = b_0 + b_1x_1 + b_2x_2 + b_3x_1^2 + b_4x_1x_2 + b_5x_2^2$$

Before t-test

After t-test

ν : 26.8

ν : 41.8

V_G : 102.11

V_G : 175.82

Solution of Coefficients

Solution of Coefficients

$$b = \begin{Bmatrix} 121.2 \\ -836.3 \\ 66.7 \\ 393.9 \\ -100 \\ 10 \end{Bmatrix},$$

$$b = \begin{Bmatrix} 0 \\ -814.0 \\ 0 \\ 352.6 \\ 0 \\ 0 \end{Bmatrix}$$

Figure 25. Significance testing, Example 1, 2nd order approximation

MEASUREMENTS OF QUALITY OF FIT
BEFORE AND AFTER t-test
PERFORMED ON

$$y = 10x_1^4 - 20x_2x_1^2 + 10x_2^2x_1^2 + x_1^2 - 2x_1 + 5$$

"Fox's Banana Function"

THIRD ORDER APPROXIMATION

$$Y = b_0 + b_1x_1 + b_2x_2 + b_3x_1^2 + b_4x_1x_2 + b_5x_2^2 + b_6x_1^3 + b_7x_1^2x_2 + b_8x_1x_2^2 + b_9x_2^3$$

Before t-test

After t-test

ν : 2.9

ν : 6.4

V_G : 53.71

V_G : 112.38

Solution of Coefficients

Solution of Coefficients

$$b = \begin{Bmatrix} -12.1 \\ 283.7 \\ 0 \\ -306.1 \\ 0 \\ 10 \\ 100 \\ -20 \\ 0 \\ 0 \end{Bmatrix}$$

$$b = \begin{Bmatrix} 0 \\ 385.0 \\ 0 \\ -349.3 \\ 0 \\ 0 \\ 103.8 \\ -17.2 \\ 0 \\ 0 \end{Bmatrix}$$

Figure 26. Significance testing, Example 1, 3rd order approximation

MEASUREMENTS OF QUALITY OF FIT
BEFORE AND AFTER t-test
PERFORMED ON

$$Y = (4 + x_1)^3 + \sin\left[\frac{3\pi}{2} * (x_1 + 1)\right] + 2 + x_2^4 + \sin\left(\frac{\pi}{2}\right) + 7x_2x_1$$

SECOND ORDER APPROXIMATION

$$Y = b_0 + b_1x_1 + b_2x_2 + b_3x_1^2 + b_4x_1x_2 + b_5x_2^2$$

Before t-test

After t-test

ν : 6.2

ν : 8.6

V_G : 90.02

V_G : 123.67

Solution of Coefficients

Solution of Coefficients

$$b = \begin{Bmatrix} 97.6 \\ 35.0 \\ -108.4 \\ 19.4 \\ 7 \\ 44.3 \end{Bmatrix} ,$$

$$b = \begin{Bmatrix} 96.4 \\ 0 \\ -90.9 \\ 29.0 \\ 0 \\ 44.3 \end{Bmatrix}$$

Figure 27. Significance testing, Example 2, 2nd order approximation

MEASUREMENTS OF QUALITY OF FIT
BEFORE AND AFTER t-test
PERFORMED ON

$$Y = (4 + x_1)^3 + \sin\left[\frac{3\pi}{2} * (x_1 + 1)\right] + 2 + x_2^4 + \sin\left(\frac{\pi}{2}\right) + 7x_2x_1$$

THIRD ORDER APPROXIMATION

$$Y = b_0 + b_1x_1 + b_2x_2 + b_3x_1^2 + b_4x_1x_2 + b_5x_2^2 + b_6x_1^3 + b_7x_1^2x_2 + b_8x_1x_2^2 + b_9x_2^3$$

Before t-test

After t-test

ν : 0.7

ν : 0.7

V_G : 27.87

V_G : 29.92

Solution of Coefficients

Solution of Coefficients

$$b = \begin{Bmatrix} 64.1 \\ 50.7 \\ 28.6 \\ 10.8 \\ 7 \\ -30.7 \\ 1.2 \\ 0 \\ 0 \\ 10 \end{Bmatrix}$$

$$b = \begin{Bmatrix} 64.1 \\ 50.8 \\ 28.6 \\ 10.8 \\ 7 \\ -30.7 \\ 0 \\ 0 \\ 0 \\ 10 \end{Bmatrix}$$

Figure 28. Significance testing, Example 2, 3rd order approximation

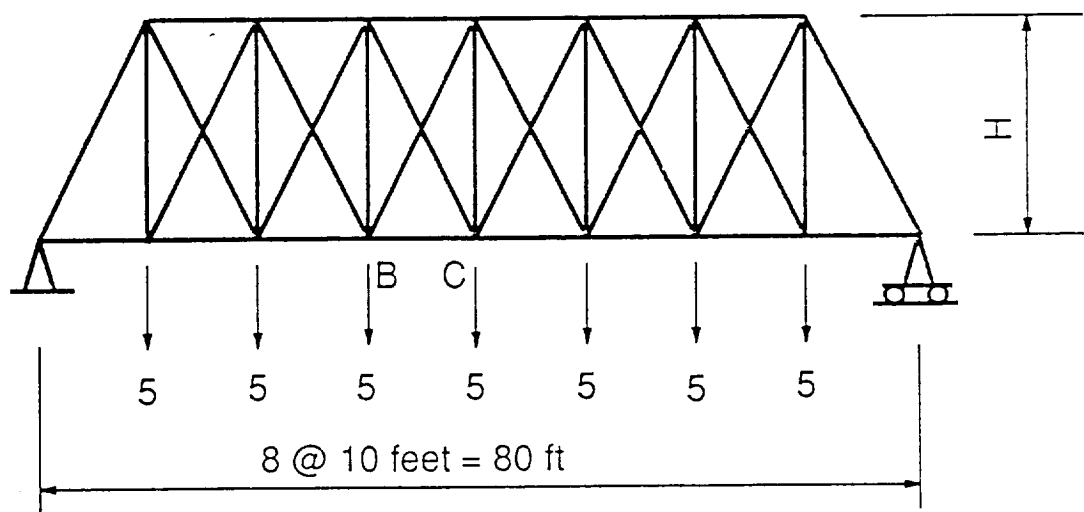


Figure 29. The 35 bar truss

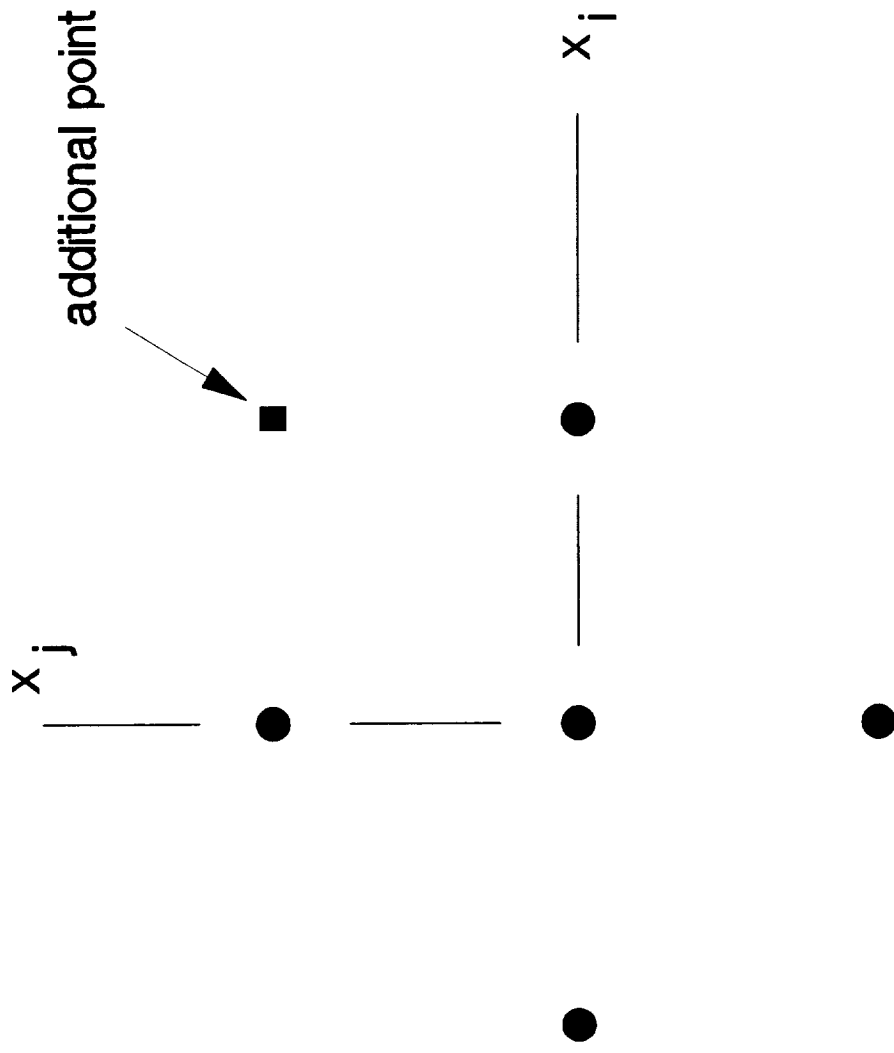


Figure 30. Additional points to complete a second order design

added to find coefficients of terms of Eq. (106)

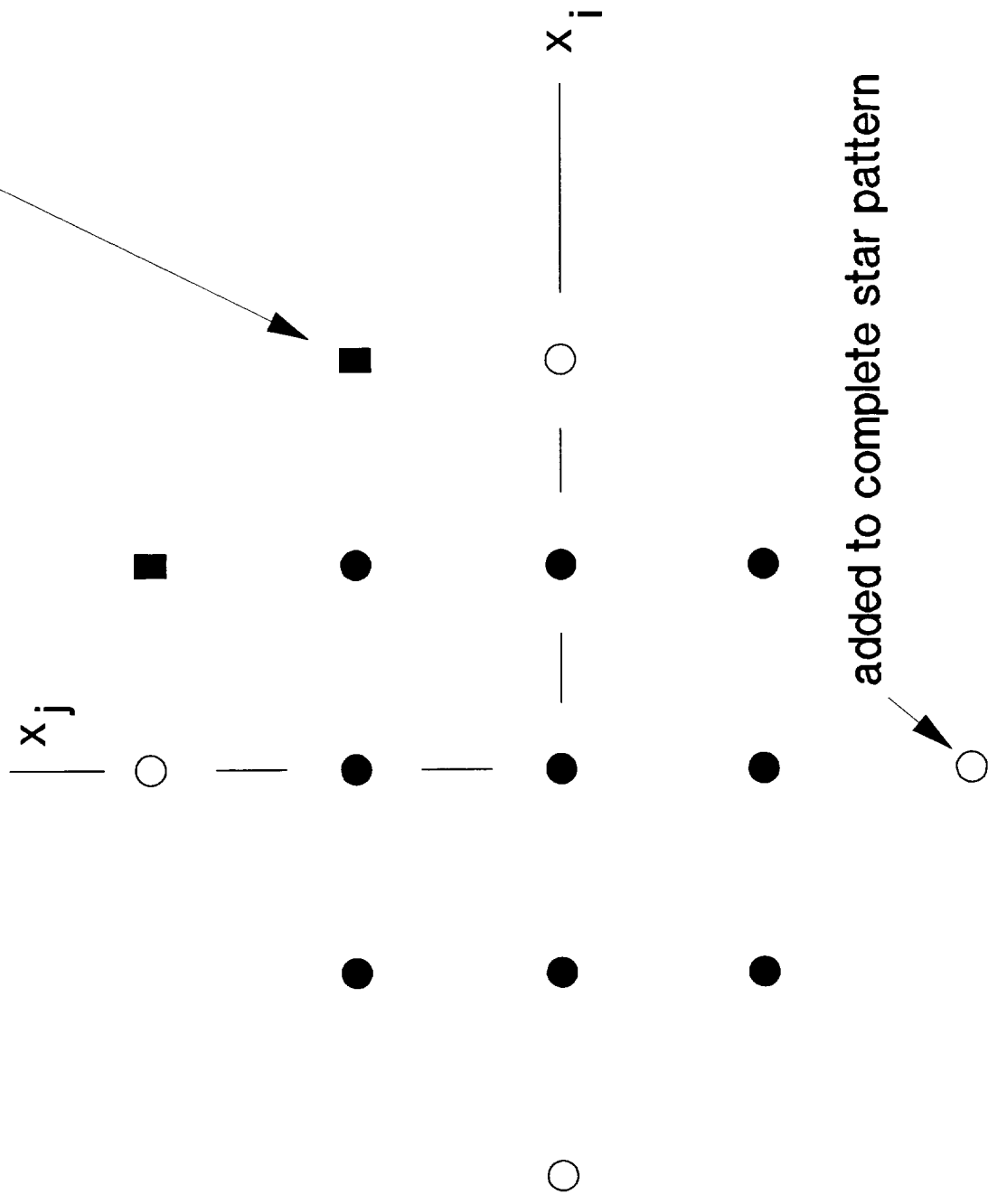


Figure 31. Additional points to complete a fourth order design

Appendix 1 Program DESIGNS

PROGRAM DESIGNS

PROGRAM TO GENERATE DESIGNS FOR 2ND ORDER POLYNOMIAL
PROGRAM DIMENSIONED FOR UP TO 20 VARIABLES
RESULTS TO SCREEN AND TO FILE designs.res
DESIGN IN GLOBAL COORDINATES TO FILE designs.run

DEFINITIONS

N = NUMBER OF DESIGN VARIABLES
M = NUMBER OF RANDOM DESIGNS POINTS

DIMENSION X(2000,20)
DIMENSION XBB(20),XBE(20),A(20),B(20)

1 FORMAT(I5,6F10.6)
2 FORMAT(' PROGRAM GENERATES DESIGNS FOR FITTING 2ND ORDER',
X' POLYNOMIAL')
3 FORMAT(' ENTER NUMBER OF DESIGN VARIABLES')
4 FORMAT(' NUMBER OF DESIGN VARIABLES = N =', I3)
11 FORMAT(6F10.6)

OPEN(UNIT=7,FILE='designs.res')
OPEN(UNIT=8,FILE='designs.run')

WRITE(6,2)
WRITE(6,3)
READ(5,*)N
WRITE(6,4)N
SET UP TERMS

NP1=N+1
NM1=N-1
M=(N*N+3*N+2)/2
MP1=M+1

ZERO DESIGN MATRIX

DO100I=1,M
DO100J=1,N
100 X(I,J)=0.
II=0

.....
GENERATE THE FIRST N+1 POINTS FOR FITTING A LINEAR FUNCTION
THE FIRST POINT IS WHEN ALL X'S ZERO, ALREADY DONE
GENERATE NEXT N POINTS

DO101I=1,N
II=I+1
101 X(II,I)=1.

.....
GENERATE NEXT N POINTS
THE 2N+1 POINTS THUS GENERATED WILL ALLOW ADDING SQUARED TERMS

DO102I=1,N
II=I+N+1
102 X(II,I)=-1.

.....
GENERATE NEXT N(N-1)/2 POINTS
THE (N*N+3*N+2)/2 POINTS THUS GENERATED WILL ALLOW ADDING CROSS
PRODUCT TERMS. WE WILL THEN HAVE COMPLETE 2ND ORDER POLYNOMIAL
APPROXIMATION

```

-      ILAST=2*N+1
      IDO=N-1
      J=1
-      JJ=2
103  CONTINUE
      DO104 I=1,IDO
-      II=I+ILAST
      X(II,J)=1.
      X(II,JJ)=1.
      JJ=JJ+1
-      104  CONTINUE
      ILAST=ILAST+IDO
      IDO=IDO-1
-      J=J+1
      JJ=J+1
      IF(J.LE.NM1)GOTO103
-
C
C      IF WE GOT HERE WE HAVE DEVELOPED THE MINIMUM POINT DESIGN
C
      WRITE(6,*)' WE HAVE GENERATED ',II,' POINTS IN THE MIN PT DESIGN'
-      WRITE(7,*)' WE HAVE GENERATED ',II,' POINTS IN THE MIN PT DESIGN'
      WRITE(6,*)' DESIGN POINTS WRITTEN TO FILE designs.res'
C
-
C      DEVELOP DESIGN POINTS TO AUGMENT THE MINIMUM POINT DESIGN
C      READ IN THE NUMBER OF RANDOM DESIGN POINTS TO BE DEVELOPED
      WRITE(6,*)' ENTER THE NUMBER OF RANDOM GENERATED DESIGN PTS',
-      X' DESIRED=M'
      READ(5,*)M
      WRITE(6,*)' NUMBER OF RANDOM DESIGN POINTS M=',M
      WRITE(7,*)' NUMBER OF RANDOM DESIGN POINTS M=',M
-      WRITE(6,*)' IFLAG IS ANY POSITIVE INTEGER USED TO START RANDOM',
      X' PROCESS'
      WRITE(6,*)' ENTER IFLAG'
-      READ(5,*)IFLAG
      WRITE(6,*)' IFLAG=',IFLAG
      WRITE(7,*)' IFLAG=',IFLAG
-      DO850 I=1,M
      II=II+1
      DO851 J=1,N
      IFLAG=IFLAG+1
-      XDUM=RAND(IFLAG)
      X(II,J)=2.*XDUM-1.
851  CONTINUE
-      850  CONTINUE
C
C      IF WE GOT HERE WE HAVE FINISHED GENERATING THE RANDOM DESIGN PTS
C      WRITE(6,*)' RANDOM DESIGN POINTS WRITTEN TO FILE designs.res'
-
C
C      PRINT OUT THE MINIMUM POINT MATRIX IN LOCAL COORDINATES
C
-      WRITE(7,*)' DESIGN MATRIX IN LOCAL COORDINATES'
      ITOTAL=II
      DO700 I=1,ITOTAL
-      WRITE(7,1) I, (X(I,J),J=1,N)
700  CONTINUE
C
C      SEE IF WE ARE TO GENERATE DESIGNS IN GLOBAL COORDINATES
-
      WRITE(6,*)' ITEST=1 IF DESIGN POINTS ARE TO BE IN GLOBAL',
      X' COORDINATES'
      WRITE(6,*)' OTHERWISE, ITEST=0'
-

```



```

-      WRITE(6,*)' ENTER ITEST'
      READ(5,*)ITEST
      IF(ITEST.NE.1)GOTO860
-C     IF WE GOT HERE WE ARE TO GENERATE DESIGNS IN GLOBAL COORDINATES
      WRITE(6,*)' ENTER LOWER AND UPPER RANGE ON EACH DESIGN VARIABLE'
      WRITE(6,*)' i.e. ENTER XBB(I) TO XBE(I)'
      DO861I=1,N
      READ(5,*)XBB(I),XBE(I)
      WRITE(6,*)' I,XBB(I),XBE(I)=' ,I,XBB(I),XBE(I)
      WRITE(7,*)' I,XBB(I),XBE(I)=' ,I,XBB(I),XBE(I)
- 861  CONTINUE
      GOTO862
860  CONTINUE
-C
C     IF WE GOT HERE LOWER BOUND VARIABLE IN GLOBAL COORDINATES IS -1
C     IF WE GOT HERE UPPER BOUND VARIABLE IN GLOBAL COORDINATES IS 1
      DO863I=1,N
      XBB(I)=-1.
      XBE(I)=1.
- 863  CONTINUE
- 862  CONTINUE
      WRITE(7,*)' I,XBB(I),XBE(I),A(I),B(I)'
      DO1301I=1,N
      A(I)=(XBE(I)-XBB(I))/2.
      B(I)=(XBE(I)+XBB(I))/2.
      WRITE(7,*)I,XBB(I),XBE(I),A(I),B(I)
1301 CONTINUE
      DO1202I=1,ITOTAL
      DO1202J=1,N
1202 X(I,J)=A(J)*X(I,J)+B(J)
-      WRITE(6,*)' DESIGN IN GLOBAL COORDINATES WRITEN TO designs.res'
      WRITE(6,*)' DESIGN IN GLOBAL COORDINATES WRITEN TO designs.run'
      WRITE(7,*)' DESIGN IN GLOBAL COORDINATES'
      WRITE(8,*)ITOTAL
      DO970I=1,ITOTAL
      WRITE(7,1)I,(X(I,J),J=1,N)
      WRITE(8,11)(X(I,J),J=1,N)
- 970  CONTINUE
      STOP
      END

```

Appendix 2
Program DESIGN4

PROGRAM DESIGN4

PROGRAM TO GENERATE DESIGNS FOR 4TH ORDER POLYNOMIAL
PROGRAM DIMENSIONED FOR UP TO 6 VARIABLES
RESULTS TO SCREEN AND TO FILE design4.res
DESIGN IN GLOBAL COORDINATES TO FILE design4.run

DEFINITIONS

N = NUMBER OF DESIGN VARIABLES
M = NUMBER OF RANDOM DESIGNS POINTS

DIMENSION X(2000,6)
DIMENSION XBB(10),XBE(10),A(10),B(10)

1 FORMAT(I5,6F10.6)
2 FORMAT(' PROGRAM GENERATES DESIGNS FOR FITTING 4TH ORDER',
X' POLYNOMIAL')
3 FORMAT(' ENTER NUMBER OF DESIGN VARIABLES')
4 FORMAT(' NUMBER OF DESIGN VARIABLES = N =', I3)
11 FORMAT(6F10.6)
OPEN(UNIT=7,FILE='design4.res')
OPEN(UNIT=8,FILE='design4.run')
WRITE(6,2)
WRITE(6,3)
READ(5,*)N
WRITE(6,4)N
IF(N.EQ.6)GOTO601
IF(N.EQ.5)GOTO501
IF(N.EQ.4)GOTO401
IF(N.EQ.3)GOTO301
IF(N.EQ.2)GOTO201
IF(N.EQ.1)GOTO101
WRITE(6,*)' PROGRAM CAN NOT DO MORE THAN 6 DESIGN VARIABLES'
WRITE(7,*)' PROGRAM CAN NOT DO MORE THAN 6 DESIGN VARIABLES'
STOP

DEVELOP 3 FACTORIAL DESIGN TO GET 4 DESIGN VARIABLE PRODUCT TERMS

101 CONTINUE
II=0
DO100I1=1,101,50
II=II+1
X(II,1)=FLOAT(I1-51)/100.
100 CONTINUE
GOTO701
201 CONTINUE
II=0
DO200I1=1,101,50
DO200I2=1,101,50
II=II+1
X(II,1)=FLOAT(I1-51)/100.
X(II,2)=FLOAT(I2-51)/100.
200 CONTINUE
GOTO701
301 CONTINUE
II=0
DO300I1=1,101,50
DO300I2=1,101,50
DO300I3=1,101,50
II=II+1
X(II,1)=FLOAT(I1-51)/100.
X(II,2)=FLOAT(I2-51)/100.

```

-      X(II,3)=FLOAT(I3-51)/100.
300  CONTINUE
      GOTO701
- 401  CONTINUE
      II=0
      DO400I1=1,101,50
-      DO400I2=1,101,50
      DO400I3=1,101,50
      DO400I4=1,101,50
      II=II+1
-      X(II,1)=FLOAT(I1-51)/100.
      X(II,2)=FLOAT(I2-51)/100.
      X(II,3)=FLOAT(I3-51)/100.
-      X(II,4)=FLOAT(I4-51)/100.
400  CONTINUE
      GOTO701
- 501  CONTINUE
      II=0
      DO500I1=1,101,50
-      DO500I2=1,101,50
      DO500I3=1,101,50
      DO500I4=1,101,50
      DO500I5=1,101,50
-      II=II+1
      X(II,1)=FLOAT(I1-51)/100.
      X(II,2)=FLOAT(I2-51)/100.
-      X(II,3)=FLOAT(I3-51)/100.
      X(II,4)=FLOAT(I4-51)/100.
      X(II,5)=FLOAT(I5-51)/100.
500  CONTINUE
-      GOTO701
C
601  CONTINUE
-      II=0
      DO600I1=1,101,50
-      DO600I2=1,101,50
      DO600I3=1,101,50
-      DO600I4=1,101,50
      DO600I5=1,101,50
      DO600I6=1,101,50
-      II=II+1
      X(II,1)=FLOAT(I1-51)/100.
      X(II,2)=FLOAT(I2-51)/100.
-      X(II,3)=FLOAT(I3-51)/100.
      X(II,4)=FLOAT(I4-51)/100.
      X(II,5)=FLOAT(I5-51)/100.
-      X(II,6)=FLOAT(I6-51)/100.
600  CONTINUE
      GOTO701
701  CONTINUE
-C
C      ENTER REST OF POINTS IN THE STAR FORMATION
C
-      DO702I=1,N
      II=II+1
      DO703J=1,N
- 703  X(II,J)=0.
      X(II,I)=1.
- 702  CONTINUE
      DO704I=1,N
-

```

```

-      II=II+1
      DO705J=1,N
705  X(II,J)=0.
-      X(II,I)=-1.
704  CONTINUE
C
C      ENTER TERMS TO CALCULATE COEFFICIENT ASSOCIATED WITH THE TERM
C      X(I)**3*X(J)
C
      NM1=N-1
-      IDO=N-1
      J=1
      JJ=2
- 803  CONTINUE
      DO804I=1,IDO
      II=II+1
      X(II,J)=1.
-      X(II,JJ)=.5
      II=II+1
-      X(II,J)=.5
      X(II,JJ)=1.
-      JJ=JJ+1
804  CONTINUE
-      IDO=IDO-1
      J=J+1
      JJ=J+1
      IF(J.LE.NM1)GOTO803
-
C      IF WE GOT HERE WE HAVE DEVELOPED THE MINIMUM POINT DESIGN
C
-      WRITE(6,*)' WE HAVE GENERATED ',II,' POINTS IN THE MIN PT DESIGN'
      WRITE(7,*)' WE HAVE GENERATED ',II,' POINTS IN THE MIN PT DESIGN'
      WRITE(6,*)' DESIGN POINTS WRITTEN TO FILE design4.res'
-
C      DEVELOP DESIGN POINTS TO AUGMENT THE MINIMUM POINT DESIGN
C      READ IN THE NUMBER OF RANDOM DESIGN POINTS TO BE DEVELOPED
C      WRITE(6,*)' ENTER THE NUMBER OF RANDOM GENERATED DESIGN PTS',
-  X' DESIRED=M'
      READ(5,*)M
      WRITE(6,*)' NUMBER OF RANDOM DESIGN POINTS M=',M
-      WRITE(7,*)' NUMBER OF RANDOM DESIGN POINTS M=',M
      WRITE(6,*)' IFLAG IS ANY POSITIVE INTEGER USED TO START RANDOM',
-  X' PROCESS'
      WRITE(6,*)' ENTER IFLAG'
      READ(5,*)IFLAG
      WRITE(6,*)' IFLAG=',IFLAG
-      WRITE(7,*)' IFLAG=',IFLAG
      DO850I=1,M
      II=II+1
-      DO851J=1,N
      IFLAG=IFLAG+1
      XDUM=RAND(IFLAG)
-      X(II,J)=2.*XDUM-1.
851  CONTINUE
- 850  CONTINUE
C
C      IF WE GOT HERE WE HAVE FINISHED GENERATING THE RANDOM DESIGN PTS
-      WRITE(6,*)' RANDOM DESIGN POINTS WRITTEN TO FILE design4.res'
C
C      PRINT OUT THE MINIMUM POINT MATRIX IN LOCAL COORDINATES
-

```

```

- C      WRITE(7,*)' DESIGN MATRIX IN LOCAL COORDINATES'
      ITOTAL=II
-      DO700I=1,ITOTAL
      WRITE(7,1)I,(X(I,J),J=1,N)
700 CONTINUE
- C      SEE IF WE ARE TO GENERATE DESIGNS IN GLOBAL COORDINATES
- C
      WRITE(6,*)' ITEST=1 IF DESIGN POINTS ARE TO BE IN GLOBAL',
X' COORDINATES'
-      WRITE(6,*)' OTHERWISE, ITEST=0'
      WRITE(6,*)' ENTER ITEST'
      READ(5,*)ITEST
-      IF(ITEST.NE.1)GOTO860
- C      IF WE GOT HERE WE ARE TO GENERATE DESIGNS IN GLOBAL COORDINATES
      WRITE(6,*)' ENTER LOWER AND UPPER RANGE ON EACH DESIGN VARIABLE'
      WRITE(6,*)' i.e. ENTER XBB(I) TO XBE(I)'
-      DO861I=1,N
      READ(5,*)XBB(I),XBE(I)
      WRITE(6,*)' I,XBB(I),XBE(I)=' ,I,XBB(I),XBE(I)
-      WRITE(7,*)' I,XBB(I),XBE(I)=' ,I,XBB(I),XBE(I)
861 CONTINUE
      GOTO862
- 860 CONTINUE
- C
- C      IF WE GOT HERE LOWER BOUND VARIABLE IN GLOBAL COORDINATES IS -1
- C      IF WE GOT HERE UPPER BOUND VARIABLE IN GLOBAL COORDINATES IS 1
-      DO863I=1,N
      XBB(I)=-1.
      XBE(I)=1.
- 863 CONTINUE
862 CONTINUE
      WRITE(7,*)' I,XBB(I),XBE(I),A(I),B(I)'
      DO1301I=1,N
      A(I)=(XBE(I)-XBB(I))/2.
      B(I)=(XBE(I)+XBB(I))/2.
      WRITE(7,*)I,XBB(I),XBE(I),A(I),B(I)
- 1301 CONTINUE
      DO1202I=1,ITOTAL
      DO1202J=1,N
- 1202 X(I,J)=A(J)*X(I,J)+B(J)
      WRITE(6,*)' DESIGN IN GLOBAL COORDINATES WRITEN TO design4.res'
      WRITE(6,*)' DESIGN IN GLOBAL COORDINATES WRITEN TO design4.run'
      WRITE(7,*)' DESIGN IN GLOBAL COORDINATES'
      WRITE(8,*)ITOTAL
      DO970I=1,ITOTAL
      WRITE(7,1)I,(X(I,J),J=1,N)
      WRITE(8,11)(X(I,J),J=1,N)
970 CONTINUE
      STOP
      END

```

Appendix 3
Program NEWPSI

PROGRAM newpsi

the program develops a polynomial approximation which
may be either under, exactly, or over determined
it can handle up to 15 design variables as programmed.
The order of polynomial it can handle is as follows:
1. one one design variable, up to a 20th order polynomial
2. two design variables, up to 5th order polynomial
3. for 2-15 design variables, a 2nd order polynomial
One can use up to 250 designs to train the approximation.
It can handle up to 2000 grid points

IMPLICIT REAL*8 (A-H,O-Z)
dimension x(250,15),y(250),a(250,136)
dimension yhat(250)
dimension b(136)
dimension xx(2000,15),yy(2000),abig(2000,136)
dimension yyhat(2000)

1 FORMAT(9F8.4)
2 FORMAT(3F12.6)
3 FORMAT(F10.6,1H,,F10.6,1H,,F10.6,1H,,F10.6,1H,,F10.6,1H,,F10.6,
X1H,,F10.6)

OPEN(UNIT=5,FILE='newpsi.dat')
OPEN(UNIT=7,FILE='newpsi.res')
OPEN(UNIT=8,FILE='newpsi.plot')

read in data

read in the print code
read(5,*)ip

enter number of design variables, nd
read(5,*)nd

enter THE DEGREE OF POLYNOMIAL TO BE CONSIDERED, np
READ(5,*)np

ENTER NUMBER OF DESIGNS FOR PROBLEM,M
READ(5,*)M

write(6,*)' print code ip=',ip
write(6,*)' number of design variables, nd=',nd

```

-- write(6,*)' degree of polynomial being considered=np=',np
-- write(6,*)' number of designs m=',m
-- write(7,*)' print code ip=',ip
-- write(7,*)' number of design variables, nd=',nd
-- write(7,*)' degree of polynomial being considered=np=',np
-- write(7,*)' number of designs m=',m

```

```

_C
C read in designs and set up matrix a
C

```

```

-- write(7,*)' x(i,j),y(i)'
-- DO101I=1,M
-- read(5,*)(x(i,j),j=1,nd),y(i)
-- write(7,*)(x(i,j),j=1,nd),y(i)

```

```

- 101 continue

```

```

C
C set up the coefficient matrix, a, in the matrix equation
C y=a x
C

```

```

call geta(ip,m,nd,np,n,x,a)

```

```

-C SEE WHETHER SYSTEM IS UNDER, EXACTLY, OR OVER DETERMINED
C

```

```

IF(M.GE.N)GOTO400
IF WE GOT HERE WE ARE UNDER-DETERMINED
WRITE(6,*)' SYSTEM IS UNDER-DETERMINED'
WRITE(7,*)' SYSTEM IS UNDER-DETERMINED'
CALL PSI(ip,M,N,A,Y,B)
GOTO402

```

```

400 CONTINUE

```

```

IF(M.GT.N)GOTO401
IF WE GOT HERE WE ARE EXACTLY DETERMINED
WRITE(6,*)' SYSTEM IS EXACTLY DETERMINED'
WRITE(7,*)' SYSTEM IS EXACTLY DETERMINED'
CALL EXACT(ip,M,A,Y,B)
GOTO402

```

```

401 CONTINUE

```

```

IF WE GOT HERE WE ARE OVER-DETERMINED
WRITE(6,*)' SYSTEM IS OVER-DETERMINED'
WRITE(7,*)' SYSTEM IS OVER-DETERMINED'
CALL OVER(ip,M,N,A,Y,B)

```

```

- 402 CONTINUE

```

```

C
C EVALUATE APPROXIMATION AT DESIGNS
C

```

```

-- WRITE(6,*)' MATRIX OF COEFFICIENTS, B(I)'
-- WRITE(7,*)' MATRIX OF COEFFICIENTS, B(I)'
-- WRITE(6,*)(B(I),I=1,N)
-- WRITE(7,*)(B(I),I=1,N)
-- WRITE(7,*)' MATRICES Y(I) AND YHAT(I)'

```

```

_C
C recalculate matrix a
C call geta(ip,m,nd,np,n,x,a)

```

```

_C
C calculate approximation at designs and print results
C

```

```

-- write(7,*)' y(i),yhat(i)'
-- DO102I=1,M
-- YHAT(I)=0.
-- DO103J=1,N

```

```

-      yhat(i)=yhat(i)+a(i,j)*b(j)
103  CONTINUE
-      WRITE(7,*)Y(I),YHAT(I)
102  CONTINUE
-
-  C
-  C      evaluate function at grid
-  C      read(5,*)ng
-  C      write(6,*)' number of designs on grid = ngn',ng
-  C      write(7,*)' number of designs on grid = ngn',ng
-  C      write(7,*)' xx(i,j),yy(i)'
-  C      DO601I=1,ng
-  C      read(5,*)(xx(i,j),j=1,nd),yy(i)
-  C      write(7,*)(xx(i,j),j=1,nd),yy(i)
- 601  continue
-  C      call getabg(ip,ng,nd,np,n,xx,abig)
-  C      write(7,*)' yy(i),yyhat(i) at grid'
-  C      DO602I=1,ng
-  C      YYHAT(I)=0.
-  C      DO603J=1,N
-  C      yyhat(i)=yyhat(i)+abig(i,j)*b(j)
- 603  CONTINUE
-  C      WRITE(7,*)YY(I),YYHAT(I)
-
-  C
-  C      write the plot file
-  C      write(8,*)(xx(i,j),j=1,nd),yyhat(i)
-
-  C
- 602  CONTINUE
-
-  C
-  C      calculate statistical terms
-  C
-  C      call statit(m,y,yhat,ng,yy,yyhat)
-  C
-  C      STOP
-  C      END
-  C      subroutine geta(ip,m,nd,np,n,x,a)
-
-  C
-  C      *****
-  C      *****
-  C
-  C      This subroutine generates the matrix a where the matrix
-  C      equation is  $y = a \cdot b$ . Here y are the training functions,
-  C      b are undetermined coefficients. The algorithm is programmed
-  C      to handle
-  C      1. any level of approximation for one design variable
-  C      2. up to 5th order polynomial in two design variables
-  C      3. quadratic approximation in more than two design variabaaales
-  C
-  C      *****
-  C      *****
-  C
-  C      IMPLICIT REAL*8 (A-H,O-Z)
-  C      dimension x(250,15),a(250,136)
-
-  C
-  C      do for each design
-  C
-  C      do300i=1,m
-
-  C
-  C      *****
-
-  C
-  C      if nd is not equal to 1 go to 400
-
-

```



```

-      if(nd.ne.1)goto400
c
c      *****
-c      *****
c
c      here we have nd=1, i.e. one design variable
-c      we will develop a's for all np's
c
c      a(i,1)=1.
-      j=1
-      do201k=1,np
-      j=j+1
-      a(i,j)=x(i,1)**k
- 201 continue
-      n=np+1
-      goto301
c
c      *****
c 400 continue
c
c      if nd is not equal to 2 go to 500
c      if(nd.ne.2)goto500
-c
c      *****
c      *****
c
c      if we got here we have 2 design variables
c
c      x1=x(i,1)
-      x2=x(i,2)
c
c      *****
-c
c      add the constant and linear terms
c
c      a(i,1)=1.
-      a(i,2)=x1
-      a(i,3)=x2
-      n=3
-      if(np.lt.2)goto301
c
c      *****
-c
c      add the 2nd order terms
c
c      a(i,4)=x1**2
-      a(i,5)=x1*x2
-      a(i,6)=x2**2
-      n=6
-      if(np.lt.3)goto301
c
c      *****
-c
c      add the cubic terms
c
c      a(i,7)=x1**3
-      a(i,8)=x1**2*x2
-      a(i,9)=x1*x2**2
-      a(i,10)=x2**3
-

```

```

-      n=10
      if(np.lt.4)goto301
- c
- c *****
- c
- c add the 4th order terms
- c
      a(i,11)=x1**4
      a(i,12)=x1**3*x2
      a(i,13)=x1**2*x2**2
-      a(i,14)=x1*x2**3
      a(i,15)=x2**4
      n=15
-      if(np.lt.5)goto301
- c
- c *****
- c
- c add the 5th order terms
- c
      a(i,16)=x1**5
-      a(i,17)=x1**4*x2
      a(i,18)=x1**3*x2**2
-      a(i,19)=x1**2*x2**3
-      a(i,20)=x1*x2**4
      a(i,21)=x2**5
      n=21
-      if(np.lt.6)goto301
- c
- c *****
- c
- c algorithm not programed for polynomials of order larger than 5
- c
      write(6,*)' for two design variables, algorithm not programed for'
-      write(6,*)' polynomials of order larger than 5'
      write(7,*)' for two design variables, algorithm not programed for'
-      write(7,*)' polynomials of order larger than 5'
      stop
- c
- c *****
- c *****
- c
500 continue
- c
- c if we got here number of design variables >2
- c
- c *****
- c
- c enter constant and linear terms
- c
      a(i,1)=1.
-      j=1
      do501k=1,nd
      j=j+1
-      a(i,j)=x(i,k)
501 continue
      n=j
      if(np.lt.2)goto301
- c
- c *****
- c

```

```

c      enter the quadratic terms
c
c      do502k=1,nd
c      do502L=k,nd
c      j=j+1
c      a(i,j)=x(i,k)*x(i,L)
502  continue
c      n=j
c      if(np.lt.3)goto301
c
c      *****
c
c      algorithm not programmed for more than quadratic approximation
c      when number of design variables >2
c
c      write(6,*)' algorithm not programmed for more than quadratic'
c      write(6,*)' approximation when number of design variables >2'
c      write(7,*)' algorithm not programmed for more than quadratic'
c      write(7,*)' approximation when number of design variables >2'
c      stop
c
c      *****
c      *****
c
c      print out some results
c
301  continue
c      if(ip.lt.4)goto302
c      write(6,*)' a(i,j)',(a(i,j),j=1,n)
c      write(6,*)' '
c      write(7,*)' a(i,j)',(a(i,j),j=1,n)
c      write(7,*)' '
302  continue
c
c      *****
c
300  continue
c      write(6,*)' number of undetermined coef=n=',n
c      write(7,*)' number of undetermined coef=n=',n
c
c      return
c      end
c      subroutine getabg(ip,m,nd,np,n,x,a)
c
c      *****
c      *****
c
c      This subroutine generates the matrix a where the matrix
c      equation is  $y = a \cdot b$ . Here y are the training functions,
c      b are undetermined coefficients. The algorithm is programmed
c      to handle
c      1. any level of approximation for one design variable
c      2. up to 5th order polynomial in two design variables
c      3. quadratic approximation in more than two design variables
c
c      *****
c      *****
c
c      IMPLICIT REAL*8 (A-H,O-Z)
c      DIMENSION A(2000,136)

```

DIMENSION X(2000,15)

do for each design

do300i=1,m

if nd is not equal to 1 go to 400

if(nd.ne.1)goto400

here we have nd=1, i.e. one design variable

we will develop a's for all np's

a(i,1)=1.

j=1

do201k=1,np

j=j+1

a(i,j)=x(i,1)**k

201 continue

n=np+1

goto301

400 continue

if nd is not equal to 2 go to 500

if(nd.ne.2)goto500

if we got here we have 2 design variables

x1=x(i,1)

x2=x(i,2)

add the constant and linear terms

a(i,1)=1.

a(i,2)=x1

a(i,3)=x2

n=3

if(np.lt.2)goto301

add the 2nd order terms

a(i,4)=x1**2

a(i,5)=x1*x2

a(i,6)=x2**2

n=6

if(np.lt.3)goto301

```
*****
```

```
add the cubic terms
```

```
a(i,7)=x1**3
```

```
a(i,8)=x1**2*x2
```

```
a(i,9)=x1*x2**2
```

```
a(i,10)=x2**3
```

```
n=10
```

```
if(np.lt.4)goto301
```

```
*****
```

```
add the 4th order terms
```

```
a(i,11)=x1**4
```

```
a(i,12)=x1**3*x2
```

```
a(i,13)=x1**2*x2**2
```

```
a(i,14)=x1*x2**3
```

```
a(i,15)=x2**4
```

```
n=15
```

```
if(np.lt.5)goto301
```

```
*****
```

```
add the 5th order terms
```

```
a(i,16)=x1**5
```

```
a(i,17)=x1**4*x2
```

```
a(i,18)=x1**3*x2**2
```

```
a(i,19)=x1**2*x2**3
```

```
a(i,20)=x1*x2**4
```

```
a(i,21)=x2**5
```

```
n=21
```

```
if(np.lt.6)goto301
```

```
*****
```

```
algorithm not programed for polynomials of order larger than 5
```

```
write(6,*)' for two design variables, algorithm not programed for'
```

```
write(6,*)' polynomials of order larger than 5'
```

```
write(7,*)' for two design variables, algorithm not programed for'
```

```
write(7,*)' polynomials of order larger than 5'
```

```
stop
```

```
*****
```

```
*****
```

```
500 continue
```

```
if we got here number of design variables >2
```

```
*****
```

```
enter constant and linear terms
```

```
a(i,1)=1.
```

```
j=1
```

```

-      do501k=1,nd
-      j=j+1
-      a(i,j)=x(i,k)
- 501 continue
-      n=j
-      if(np.lt.2)goto301
-
- c      *****
- c
- c      enter the quadratic terms
- c
-      do502k=1,nd
-      do502L=k,nd
-      j=j+1
-      a(i,j)=x(i,k)*x(i,L)
- 502 continue
-      n=j
-      if(np.lt.3)goto301
-
- c      *****
- c
- c      algorithm not programmed for more than quadratic approximation
- c      when number of design variables >2
- c
-      write(6,*)' algorithm not programmed for more than quadratic'
-      write(6,*)' approximation when number of design variables >2'
-      write(7,*)' algorithm not programmed for more than quadratic'
-      write(7,*)' approximation when number of design variables >2'
-      stop
-
- c      *****
- c      *****
- c
- c      print out some results
- c
- 301 continue
-      if(ip.lt.4)goto302
-      write(6,*)' a(i,j)',(a(i,j),j=1,n)
-      write(6,*)' '
-      write(7,*)' a(i,j)',(a(i,j),j=1,n)
-      write(7,*)' '
- 302 continue
-
- c      *****
- c      *****
- c
- 300 continue
-      write(6,*)' number of undetermined coef=n=',n
-      write(7,*)' number of undetermined coef=n=',n
-
- c      return
-      end
-      SUBROUTINE PSI(IP,M,N,DUMA,Y,XX)
-      IMPLICIT REAL*8 (A-H,O-Z)
-
- c      DIMENSION DUMA(250,136)
-      DIMENSION A(21,21),B(21,21),D(21,21),DI(21,21),BPI(21,21)
-      DIMENSION C(21,21),FI(21,21),CPI(21,21),H(21,21),HI(21,21)
-      DIMENSION API(21,21)
-      DIMENSION F(21,21)
-      DIMENSION IPIVOT(21),IWK(21,2)
-

```

```
DIMENSION Y(250)
DIMENSION XX(21)
```

```
THIS SUBROUTINE CALCULATES PSEUDO INVERSE OF MATRIX A
M      = ROW DIMENSION OF A LESS THAN N
N      = COLUMN DIMENSION OF A
```

```
COPY DUMA TO A
```

```
DO90I=1,M
DO90J=1,N
90 A(I,J)=DUMA(I,J)
```

```
PRINT MATRIX A
if(ip.lt.4)goto50
WRITE(6,*)' MATRIX A'
WRITE(7,*)' MATRIX A'
CALL WRITIT(M,N,A)
50 continue
```

```
SET UP MATRIX B
```

```
DO100I=1,M
DO100J=1,M
100 B(I,J)=A(I,J)
if(ip.lt.4)goto51
WRITE(6,*)' MATRIX B'
WRITE(7,*)' MATRIX B'
CALL WRITIT(M,M,B)
51 continue
```

```
GET D= B TRAN * B
```

```
DO101I=1,M
DO101J=1,M
D(I,J)=0.
DO101K=1,M
101 D(I,J)=D(I,J)+B(K,I)*B(K,J)
if(ip.lt.4)goto52
WRITE(6,*)' MATRIX D'
WRITE(7,*)' MATRIX D'
CALL WRITIT(M,M,D)
52 continue
```

```
GET INVERSE OF D=DI
```

```
MAX=21
MDUM=0
IOP=0
CALL MATINV(MAX,M,D,MDUM,DI,IOP,DETERM,ISCALE,IPIVOT,IWK)
WRITE(6,*)' DETERM=',DETERM,' ISCALE=',ISCALE
WRITE(7,*)' DETERM=',DETERM,' ISCALE=',ISCALE
DO300I=1,M
DO300J=1,M
300 DI(I,J)=D(I,J)
if(ip.lt.4)goto53
WRITE(6,*)' MATRIX DI'
WRITE(7,*)' MATRIX DI'
CALL WRITIT(M,M,DI)
53 continue
```

GET PSEUDO INVERSE OF $B = BPI = DI * B \text{ TRANS}$

```
DO102I=1,M
DO102 JQ=1,M
BPI(I,JQ)=0.
DO102J=1,M
102 BPI(I,JQ)=BPI(I,JQ)+DI(I,J)*B(JQ,J)
   if(ip.lt.4)goto54
   WRITE(6,*)' MATRIX BPI'
   WRITE(7,*)' MATRIX BPI'
   CALL WRITIT(M,M,BPI)
54 continue
```

SET UP MATRIX $C = A$

```
DO103I=1,M
DO103J=1,N
103 C(I,J)=A(I,J)
   if(ip.lt.4)goto55
   WRITE(6,*)' MATRIX C'
   WRITE(7,*)' MATRIX C'
   CALL WRITIT(M,N,C)
55 continue
```

SET UP MATRIX $F = C * C \text{ TRANS}$

```
DO104I=1,M
DO104J=1,M
F(I,J)=0.
DO104K=1,N
104 F(I,J)=F(I,J)+C(I,K)*C(J,K)
   if(ip.lt.4)goto56
   WRITE(6,*)' MATRIX F'
   WRITE(7,*)' MATRIX F'
   CALL WRITIT(M,M,F)
56 continue
```

GET THE INVERSE OF $F = FI$

```
CALL MATINV(MAX,M,F,MDUM,FI,IOP,DETERM,ISCALE,IPIVOT,IWK)
WRITE(6,*)' DETERM=',DETERM,' ISCALE=',ISCALE
WRITE(7,*)' DETERM=',DETERM,' ISCALE=',ISCALE
DO301I=1,M
DO301J=1,M
301 FI(I,J)=F(I,J)
   if(ip.lt.4)goto57
   WRITE(6,*)' MATRIX FI'
   WRITE(7,*)' MATRIX FI'
   CALL WRITIT(M,M,FI)
57 continue
```

GET THE PSEUDO INVERSE OF $C = CPI = C \text{ TRANS} * FI$

```
DO105IQ=1,N
DO105J=1,M
CPI(IQ,J)=0.
DO105I=1,M
105 CPI(IQ,J)=CPI(IQ,J)+C(I,IQ)*FI(I,J)
   if(ip.lt.4)goto58
```



```

-      WRITE(6,*) ' MATRIX CPI'
-      WRITE(7,*) ' MATRIX CPI'
-      CALL WRITIT(N,M,CPI)
-- 58 continue
C
C      SET UP MATRIX H = PSEUDO INVERSE OF B = BPI
C
-      DO106I=1,M
-      DO106J=1,M
106 H(I,J)=BPI(I,J)
-      if(ip.lt.4)goto59
-      WRITE(6,*) ' MATRIX H'
-      WRITE(7,*) ' MATRIX H'
-      CALL WRITIT(M,M,H)
- 59 continue
C
C      GET INVERSE OF H = HI
-      CALL MATINV(MAX,M,H,MDUM,HI,IOP,DETERM,ISCALE,IPIVOT,IWK)
-      WRITE(6,*) ' DETERM=',DETERM,' ISCALE=',ISCALE
-      WRITE(7,*) ' DETERM=',DETERM,' ISCALE=',ISCALE
-      DO302I=1,M
-      DO302J=1,M
302 HI(I,J)=H(I,J)
-      if(ip.lt.4)goto60
-      WRITE(6,*) ' MATRIX HI'
-      WRITE(7,*) ' MATRIX HI'
-      CALL WRITIT(M,M,HI)
- 60 continue
C
C      GET PSEUDO INVERSE OF A = API = CPI * HI * BPI
-C
-      DO107IQ=1,N
-      DO107J=1,M
-      API(IQ,J)=0.
-      DO107I=1,M
-      DO107K=1,M
107 API(IQ,J)=API(IQ,J)+CPI(IQ,I)*HI(I,K)*BPI(K,J)
-      if(ip.lt.4)goto61
-      WRITE(6,*) ' MATRIX API'
-      WRITE(7,*) ' MATRIX API'
-      CALL WRITIT(N,M,API)
- 61 continue
C
C      GET XX = API * Y
-C
-      DO108IQ=1,N
-      XX(IQ)=0.
-      DO108J=1,M
108 XX(IQ)=XX(IQ)+API(IQ,J)*Y(J)
-      JDUM=1
-      if(ip.lt.4)goto62
-      WRITE(6,*) ' MATRIX XX'
-      WRITE(7,*) ' MATRIX XX'
-      CALL WRITIT(N,JDUM,XX)
- 62 continue
C
-      RETURN
-      END
-      SUBROUTINE WRITIT(MM,NN,XX)
-      IMPLICIT REAL*8 (A-H,O-Z)
-

```

```

-      DIMENSION XX(21,1)
1  FORMAT(1X)
2  FORMAT(10F7.2)
-      WRITE(6,1)
      DO100I=1,MM
      WRITE(6,2) (XX(I,J),J=1,NN)
-      WRITE(7,2) (XX(I,J),J=1,NN)
100  CONTINUE
      RETURN
      END
-      SUBROUTINE EXACT(IP,M,A,Y,B)
      IMPLICIT REAL*8 (A-H,O-Z)

C
-      DIMENSION a(250,136),b(136),y(250)
      DIMENSION IPIVOT(250),IWK(250,2)
      DIMENSION C(136,1)
      DO100I=1,M
-      100  C(I,1)=Y(I)
      MAX=250
      MDUM=1
      IOP=0
      CALL MATINV(MAX,M,A,MDUM,C,IOP,DETERM,ISCALE,IPIVOT,IWK)
      WRITE(6,*)' DETERM=',DETERM,' ISCALE=',ISCALE
-      WRITE(7,*)' DETERM=',DETERM,' ISCALE=',ISCALE
      DO101I=1,M
      B(I)=C(I,1)
-      101  CONTINUE
      if(ip.lt.4)goto50
      WRITE(6,*)' MATRIX B',(B(I),I=1,M)
      WRITE(7,*)' MATRIX B',(B(I),I=1,M)
-      50  continue
      RETURN
      END
-      SUBROUTINE OVER(IP,M,N,A,Y,B)
      IMPLICIT REAL*8 (A-H,O-Z)
      DIMENSION a(250,136),b(136),y(250)
      DIMENSION IPIVOT(136),IWK(136,2)
-      DIMENSION ATA(136,136),ATY(136,1)
      DO200I=1,N
      DO200J=1,N
-      ATA(I,J)=0.
      DO200K=1,M
-      200  ATA(I,J)=ATA(I,J)+A(K,I)*A(K,J)
      DO201I=1,N
      ATY(I,1)=0.
      DO201K=1,M
-      201  ATY(I,1)=ATY(I,1)+A(K,I)*Y(K)
      MAX=136
      MDUM=1
      IOP=0
      CALL MATINV(MAX,N,ATA,MDUM,ATY,IOP,DETERM,ISCALE,IPIVOT,IWK)
      WRITE(6,*)' DETERM=',DETERM,' ISCALE=',ISCALE
      WRITE(7,*)' DETERM=',DETERM,' ISCALE=',ISCALE
-      DO101I=1,N
      B(I)=ATY(I,1)
-      101  CONTINUE
      if(ip.lt.4)goto50
      WRITE(6,*)' MATRIX B',(B(I),I=1,N)
      WRITE(7,*)' MATRIX B',(B(I),I=1,N)
-      50  continue

```

```

-      RETURN
-      END
-      subroutine statit(m,y,yhat,ng,yy,yyhat)
-      implicit real*8 (a-h,o-z)
-
-      C
-      C *****
-      C
-      C This subroutine calculates quality of approximation measures
-      C this subroutine calculates v, r2, and vg
-      C
-      C *****
-      C
-      dimension y(250),yhat(250)
-      dimension yy(2000),yyhat(2000)
-      yb=0.
-      do100id=1,m
-      yb=yb+y(id)
-      100 continue
-      yb=yb/float(m)
-      error=0.
-      do101id=1,m
-      error=error+(y(id)-yhat(id))**2
-      101 continue
-      v=sqrt(error/float(m))/yb*(100.)
-      write(7,*)' error over designs=error = ',error
-      write(7,*)' average y over design = yb =',yb
-      write(6,*)' coefficient v (as %)= ',v
-      write(7,*)' coefficient v (as %)= ',v
-      dn=0.
-      dd=0.
-      do7769id=1,m
-      dn=dn+(yhat(id)-yb)**2
-      dd=dd+(y(id)-yb)**2
-      7769 continue
-      r2=dn/dd*(100.)
-      write(6,*)' coefficient r2 (as%) = ',r2
-      write(7,*)' coefficient r2 (as%) = ',r2
-      C
-      get vg
-      perror=0.
-      yg=0.
-      do155id=1,ng
-      yg=yg+yy(id)
-      perror=pererror+(yy(id)-yyhat(id))**2
-      155 continue
-      yg=yg/float(ng)
-      vg=sqrt(pererror/float(ng))/yg*(100.)
-      write(7,*)' sum of residuals squared=pererror=',pererror
-      write(7,*)' average y over grid = yg =',yg
-      write(6,*)' coefficient vg = ',vg
-      write(7,*)' coefficient vg = ',vg
-      return
-      end
-      SUBROUTINE MATINV(MAX,N,A,M,B,IOP,DETERM,ISCALE,IPIVOT,IWK)
-      implicit real*8 (a-h,o-z)
-
-      C
-      C F1.3
-      C *****
-      C
-      C PURPOSE - MATINV INVERTS A REAL SQUARE MATRIX A.
-      C IN ADDITION THE ROUTINE SOLVES THE MATRIX
-
-      MATINV 2
-      MATINV 3
-      MATINV 4
-      MATINV 5
-      MATINV 6
-      MATINV 7

```

EQUATION $AX=B$, WHERE B IS A MATRIX OF CONSTANT VECTORS. THERE IS ALSO AN OPTION TO HAVE THE DETERMINANT EVALUATED. IF THE INVERSE IS NOT NEEDED, USE GELIM TO SOLVE A SYSTEM OF SIMULTANEOUS EQUATIONS AND DETFAC TO EVALUATE A DETERMINANT FOR SAVING TIME AND STORAGE.

USE - CALL MATINV(MAX,N,A,M,B,IOP,DETERM,ISCALE,IPIVOT,IWK)

MAX - THE MAXIMUM ORDER OF A AS STATED IN THE DIMENSION STATEMENT OF THE CALLING PROGRAM.

N - THE ORDER OF A, 1.LE.N.LE.MAX.

A - A TWO-DIMENSIONAL ARRAY OF THE COEFFICIENTS. ON RETURN TO THE CALLING PROGRAM, A INVERSE IS STORED IN A.
A MUST BE DIMENSIONED IN THE CALLING PROGRAM WITH FIRST DIMENSION MAX AND SECOND DIMENSION AT LEAST N.

M - THE NUMBER OF COLUMN VECTORS IN B.
M=0 SIGNALS THAT THE SUBROUTINE IS USED SOLELY FOR INVERSION, HOWEVER, IN THE CALL STATEMENT AN ENTRY CORRESPONDING TO B MUST BE PRESENT.

B - A TWO-DIMENSIONAL ARRAY OF THE CONSTANT VECTOR B. ON RETURN TO CALLING PROGRAM, X IS STORED IN B. B SHOULD HAVE ITS FIRST DIMENSION MAX AND ITS SECOND AT LEAST M.

IOP - COMPUTE DETERMINANT OPTION.
IOP=0 COMPUTES THE MATRIX INVERSE AND DETERMINANT.
IOP=1 COMPUTES THE MATRIX INVERSE ONLY.

DETERM- FOR IOP=0-IN CONJUNCTION WITH ISCALE REPRESENTS THE VALUE OF THE DETERMINANT OF A, $DET(A)$, AS FOLLOWS.
 $DET(A) = (DETERM) (10^{**}100 (ISCALE))$
THE COMPUTATION $DET(A)$ SHOULD NOT BE ATTEMPTED IN THE USER PROGRAM SINCE IF THE ORDER OF A IS LARGER AND/OR THE MAGNITUDE OF ITS ELEMENTS ARE LARGE (SMALL), THE $DET(A)$ CALCULATION MAY CAUSE OVERFLOW (UNDERFLOW). DETERM SET TO ZERO FOR SINGULAR MATRIX CONDITION, FOR EITHER IOP=1, OR 0. SHOULD BE CHECKED BY PROGRAMER ON RETURN TO MAIN PROGRAM.

ISCALE - A SCALE FACTOR COMPUTED BY THE SUBROUTINE TO AVOID OVERFLOW OR UNDERFLOW IN THE COMPUTATION OF THE QUANTITY, DETERM.

IPIVOT - A ONE DIMENSIONAL INTEGER ARRAY USED BY THE SUBPROGRAM TO STORE PIVOTOL INFORMATION. IT SHOULD BE DIMENSIONED AT LEAST N. IN GENERAL

MATINV 8
MATINV 9
MATINV10
MATINV11
MATINV12
MATINV13
MATINV14
MATINV15
MATINV16
MATINV17
MATINV18
MATINV19
MATINV20
MATINV21
MATINV22
MATINV23
MATINV24
MATINV25
MATINV26
MATINV27
MATINV28
MATINV29
MATINV30
MATINV31
MATINV32
MATINV33
MATINV34
MATINV35
MATINV36
MATINV37
MATINV38
MATINV39
MATINV40
MATINV41
MATINV42
MATINV43
MATINV44
MATINV45
MATINV46
MATINV47
MATINV48
MATINV49
MATINV50
MATINV51
MATINV52
MATINV53
MATINV54
MATINV55
MATINV56
MATINV57
MATINV58
MATINV59
MATINV60
MATINV61
MATINV62
MATINV63
MATINV64
MATINV65
MATINV66
MATINV67

- C	THE USER DOES NOT NEED TO MAKE USE	MATINV68
C	OF THIS ARRAY.	MATINV69
C		MATINV70
- C	IWK - A TWO-DIMENSIONAL INTEGER ARRAY OF	MATINV71
C	TEMPORARY STORAGE USED BY THE ROUTINE.	MATINV72
C	IWK SHOULD HAVE ITS FIRST DIMENSION	MATINV73
- C	MAX, AND ITS SECOND 2.	MATINV74
C		MATINV75
C	REQUIRED ROUTINES-	MATINV76
C		MATINV77
- C	REFERENCE -FOX,L, AN INTRODUCTION TO NUMERICAL	MATINV78
C	LINEAR ALGEBRA	MATINV79
C		MATINV80
- C	STORAGE - 542 OCTAL LOCATIONS	MATINV81
C		MATINV82
C	LANGUAGE -FORTRAN	MATINV83
- C	LIBRARY FUNCTIONS -ABS	MATINV84
C		MATINV85
C	RELEASED - JULY 1973	MATINV86
C		MATINV87
- C	LATEST REVISION - JULY 29, 1981	MATINV88
C	COMPUTER SCIENCES CORPORATION	MATINV89
C	HAMPTON, VA	MATINV90
- C	*****	MATINV91
C		MATINV92
	DIMENSION IPIVOT(N),A(MAX,N),B(MAX,N),IWK(MAX,2)	MATINV93
	EQUIVALENCE (IROW,JROW), (ICOLUM,JCOLUM), (AMAX, T, SWAP)	MATINV94
- C		MATINV98
C	INITIALIZATION	MATINV99
C		MATIN100
-	ISCALE=0	MATIN101
	R1=(10.0d+00)**32	MATIN102
	R2=1.0d+00/R1	MATIN103
-	DETERM=1.0d+00	MATIN104
	DO 20 J=1,N	MATIN105
	IPIVOT(J)=0	MATIN106
20	CONTINUE	MATIN107
-	DO 550 I=1,N	MATIN108
C		MATIN109
C	SEARCH FOR PIVOT ELEMENT	MATIN110
- C		MATIN111
	AMAX=0.0d+00	MATIN112
	DO 105 J=1,N	MATIN113
-	IF (IPIVOT(J)-1) 60, 105, 60	MATIN114
60	DO 100 K=1,N	MATIN115
	IF (IPIVOT(K)-1) 80, 100, 740	MATIN116
80	TMAX = ABS(A(J,K))	MATIN117
-	IF(AMAX-TMAX) 85,100,100	MATIN118
85	IROW=J	MATIN119
	ICOLUM=K	MATIN120
-	AMAX=TMAX	MATIN121
100	CONTINUE	MATIN122
105	CONTINUE	MATIN123
	IF (AMAX) 740,106,110	MATIN124
-	106 DETERM=0.0d+00	MATIN125
	ISCALE=0	MATIN126
	GO TO 740	MATIN127
-	110 IPIVOT(ICOLUM) = 1	MATIN128
C		MATIN129
C	INTERCHANGE ROWS TO PUT PIVOT ELEMENT ON DIAGONAL	MATIN130

- C			MATIN131
			MATIN132
			MATIN133
140	IF (IROW-ICOLUM) 140, 260, 140		MATIN134
	DETERM=-DETERM		MATIN135
	DO 200 L=1,N		MATIN136
	SWAP=A(IROW,L)		MATIN137
	A(IROW,L)=A(ICOLUM,L)		MATIN138
	A(ICOLUM,L)=SWAP		MATIN139
200	CONTINUE		MATIN140
	IF(M) 260, 260, 210		MATIN141
210	DO 250 L=1, M		MATIN142
	SWAP=B(IROW,L)		MATIN143
	B(IROW,L)=B(ICOLUM,L)		MATIN144
	B(ICOLUM,L)=SWAP		MATIN145
250	CONTINUE		MATIN146
260	IWK(I,1)=IROW		MATIN147
	IWK(I,2)=ICOLUM		MATIN148
	PIVOT=A(ICOLUM,ICOLUM)		MATIN149
	IF(IOP) 740,1000,321		MATIN150
C			MATIN151
C	SCALE THE DETERMINANT		MATIN152
-C			MATIN153
1000	PIVOTI=PIVOT		MATIN154
	IF(ABS(DETERM)-R1) 1030,1010,1010		MATIN155
1010	DETERM=DETERM/R1		MATIN156
	ISCALE=ISCALE+1		MATIN157
	IF(ABS(DETERM)-R1) 1060,1020,1020		MATIN158
1020	DETERM=DETERM/R1		MATIN159
	ISCALE=ISCALE+1		MATIN160
	GO TO 1060		MATIN161
1030	IF(ABS(DETERM)-R2) 1040,1040,1060		MATIN162
1040	DETERM=DETERM*R1		MATIN163
	ISCALE=ISCALE-1		MATIN164
	IF(ABS(DETERM)-R2) 1050,1050,1060		MATIN165
1050	DETERM=DETERM*R1		MATIN166
	ISCALE=ISCALE-1		MATIN167
1060	IF(ABS(PIVOTI)-R1) 1090,1070,1070		MATIN168
1070	PIVOTI=PIVOTI/R1		MATIN169
	ISCALE=ISCALE+1		MATIN170
	IF(ABS(PIVOTI)-R1) 320,1080,1080		MATIN171
1080	PIVOTI=PIVOTI/R1		MATIN172
	ISCALE=ISCALE+1		MATIN173
	GO TO 320		MATIN174
1090	IF(ABS(PIVOTI)-R2) 2000,2000,320		MATIN175
2000	PIVOTI=PIVOTI*R1		MATIN176
	ISCALE=ISCALE-1		MATIN177
	IF(ABS(PIVOTI)-R2) 2010,2010,320		MATIN178
2010	PIVOTI=PIVOTI*R1		MATIN179
	ISCALE=ISCALE-1		MATIN180
320	DETERM=DETERM*PIVOTI		MATIN181
C			MATIN182
-C	DIVIDE PIVOT ROW BY PIVOT ELEMENT		MATIN183
C			MATIN184
321	A(ICOLUM,ICOLUM)=1.0d+00		MATIN185
	DO 350 L=1,N		MATIN186
350	A(ICOLUM,L)=A(ICOLUM,L)/PIVOT		MATIN187
	IF(M) 380, 380, 360		MATIN188
360	DO 370 L=1,M		MATIN189
370	B(ICOLUM,L)=B(ICOLUM,L)/PIVOT		MATIN190
C			
C	REDUCE NON-PIVOT ROWS		

-C			MATIN191
	380	DO 550 L1=1,N	MATIN192
		IF(L1-ICOLUM) 400, 550, 400	MATIN193
-	400	T=A(L1,ICOLUM)	MATIN194
		A(L1,ICOLUM)=0.0d+00	MATIN195
		DO 450 L=1,N	MATIN196
-	450	A(L1,L)=A(L1,L)-A(ICOLUM,L)*T	MATIN197
		IF(M) 550, 550, 460	MATIN198
	460	DO 500 L=1,M	MATIN199
	500	B(L1,L)=B(L1,L)-B(ICOLUM,L)*T	MATIN200
-	550	CONTINUE	MATIN201
C			MATIN202
C		INTERCHANGE COLUMNS	MATIN203
-C			MATIN204
		DO 710 I=1,N	MATIN205
		L=N+1-I	MATIN206
		IF (IWK(L,1)-IWK(L,2)) 630,710,630	MATIN207
-	630	JROW=IWK(L,1)	MATIN208
		JCOLUM=IWK(L,2)	MATIN209
		DO 705 K=1,N	MATIN210
-		SWAP=A(K,JROW)	MATIN211
		A(K,JROW)=A(K,JCOLUM)	MATIN212
		A(K,JCOLUM)=SWAP	MATIN213
-	705	CONTINUE	MATIN214
	710	CONTINUE	MATIN215
	740	RETURN	MATIN216
		END	MATIN217
-C		ROUTINE NAME - HC318=EPSLON	EPSLON 2
C		FROM EISPACK	EPSLON 3
C			EPSLON 4
-C		-----	EPSLON 5
C			EPSLON 6
C	LATEST REVISION	- AUGUST 1,1984	EPSLON 7
-C		COMPUTER SCIENCES CORP., HAMPTON, VA.	EPSLON 8
C			EPSLON 9
C	PURPOSE	- THE FORTRAN FUNCTION EPSLON ESTIMATES UNIT	EPSLON10
C		ROUND OFF IN QUANTITIES OF SIZE X.	EPSLON11
-C			EPSLON12
C	USAGE	- VARIABLE = EPSLON(X)	EPSLON13
C			EPSLON14
-C	ARGUMENTS X	- IS A REAL INPUT VARIABLE WHICH REPRESENTS THE	EPSLON15
C		QUANTITIES OF SIZE IN WHICH UNIT ROUND OFF	EPSLON16
C		WILL BE ESTIMATED.	EPSLON17
-C			EPSLON18
C	REQUIRED ROUTINES	- NONE	EPSLON19
C			EPSLON20
C	REMARKS 1.	IT SHOULD BE NOTED THAT EPSLON IS A FUNCTION	EPSLON21
-C		DESIGNED TO BE CALLED BY ROUTINES IN THE	EPSLON22
C		EISPACK VERSION 3.	EPSLON23
C			EPSLON24
-C		THIS PROGRAM SHOULD FUNCTION PROPERLY ON ALL	EPSLON25
C		SYSTEMS SATISFYING THE FOLLOWING TWO	EPSLON26
C		ASSUMPTIONS,	EPSLON27
-C			EPSLON28
C	A.	THE BASE USED IN REPRESENTING FLOATING	EPSLON29
C		POINT NUMBERS IS NOT A POWER OF THREE.	EPSLON30
-C			EPSLON31
C	B.	THE QUANTITY A IN STATEMENT 10 IS	EPSLON32
-C		REPRESENTED TO THE ACCURACY USED IN FLOATING	EPSLON33
C		POINT VARIABLES THAT ARE STORED IN MEMORY.	EPSLON34

-C		EPSLON35
C	THE STATEMENT NUMBER 10 AND THE GO TO 10 ARE	EPSLON36
C	INTENDED TO FORCE OPTIMIZING COMPILERS TO	EPSLON37
-C	GENERATE CODE SATISFYING ASSUMPTION 2.	EPSLON38
C		EPSLON39
C	UNDER THESE ASSUMPTIONS, IT SHOULD BE TRUE	EPSLON40
-C	THAT,	EPSLON41
C		EPSLON42
C	A IS NOT EXACTLY EQUAL TO FOUR-THIRDS,	EPSLON43
C		EPSLON44
-C	B HAS A ZERO FOR ITS LAST BIT OR DIGIT,	EPSLON45
C		EPSLON46
C	C IS NOT EXACTLY EQUAL TO ONE,	EPSLON47
-C		EPSLON48
C	EPS MEASURES THE SEPARATION OF 1.0 FROM THE	EPSLON49
C	NEXT LARGER FLOATING POINT NUMBER.	EPSLON50
-C		EPSLON51
C	EXAMPLE :	EPSLON52
C	PROGRAM TR(OUTPUT,TAPE6=OUTPUT)	EPSLON53
C	REAL X	EPSLON54
-C	X = 4.	EPSLON55
C	A = EPSLON(X)	EPSLON56
C	WRITE(6,100) A	EPSLON57
-C100	FORMAT(5H0A = ,G22.14)	EPSLON58
C	STOP	EPSLON59
C	END	EPSLON60
C	OUTPUT :	EPSLON61
-CA =	.56843418860808E-13	EPSLON62
C		EPSLON63
C	-----	EPSLON64
-C	*F45V1P0*	EPSLON65
C	REAL*8 FUNCTION EPSLON (X)	EISPAK
-		EISPAK32
C	REAL*8 X	EISPAK
-	REAL*8 A,B,C,EPS	EISPAK
	A = 4.0E0/3.0E0	EISPAK35
10	B = A - 1.0E0	EISPAK36
-	C = B + B + B	EISPAK37
	EPS = ABS(C-1.0E0)	EISPAK38
-	IF (EPS .EQ. 0.0E0) GO TO 10	EISPAK39
	EPSLON = EPS*ABS(X)	EISPAK40
-	RETURN	EISPAK41
C**	THIS PROGRAM VALID ON FTN4 AND FTN5 **	EISPAK42
	END	EISPAK43
-C	ROUTINE NAME - PF260=QZHES	QZHES 2
C	FROM EISPACK	QZHES 3
C		QZHES 4
-C	-----	QZHES 5
C		QZHES 6
C	LATEST REVISION - AUGUST 1,1984	QZHES 7
-C	COMPUTER SCIENCES CORP., HAMPTON, VA.	QZHES 8
C		QZHES 9
C		QZHES 10
C	PURPOSE - THIS SUBROUTINE ACCEPTS A PAIR OF REAL	QZHES 11
-C	GENERAL MATRICES AND REDUCES ONE OF THEM TO	QZHES 12
C	UPPER HESSENBERG FORM AND THE OTHER TO UPPER	QZHES 13
C	TRIANGULAR FORM USING ORTHOGONAL	QZHES 14
-C	TRANSFORMATIONS. IT IS USUALLY FOLLOWED BY	QZHES 15
C	QZIT(PF261), QZVAL(PF262) AND, POSSIBLY,	QZHES 16
C	QZVEC(PF263).	QZHES 17

USAGE

- CALL QZHES(NM,N,A,B,MATZ,Z)

ARGUMENTS

NM

- ON INPUT NM MUST BE SET TO THE ROW DIMENSION OF TWO-DIMENSIONAL ARRAY PARAMETERS AS DECLARED IN THE CALLING PROGRAM DIMENSION STATEMENT.

N

- ON INPUT N IS THE ORDER OF THE MATRICES.

A

- ON INPUT A CONTAINS A REAL GENERAL MATRIX. MUST BE OF DIMENSION NM X N.

ON OUTPUT A HAS BEEN REDUCED TO UPPER HESSENBERG FORM. THE ELEMENTS BELOW THE FIRST SUBDIAGONAL HAVE BEEN SET TO ZERO.

B

- ON INPUT B CONTAINS A REAL GENERAL MATRIX. MUST BE OF DIMENSION NM X N.

ON OUTPUT B HAS BEEN REDUCED TO UPPER TRIANGULAR FORM. THE ELEMENTS BELOW THE MAIN DIAGONAL HAVE BEEN SET TO ZERO.

MATZ

- ON INPUT MATZ SHOULD BE SET TO .TRUE. IF THE RIGHT HAND TRANSFORMATIONS ARE TO BE ACCUMULATED FOR LATER USE IN COMPUTING EIGENVECTORS, AND TO .FALSE. OTHERWISE.

Z

- ON OUTPUT Z CONTAINS THE PRODUCT OF THE RIGHT HAND TRANSFORMATIONS IF MATZ HAS BEEN SET TO .TRUE. OTHERWISE, Z IS NOT REFERENCED. MUST BE OF DIMENSION NM X N.

REQUIRED ROUTINES - NONE

REMARKS 1. THIS SUBROUTINE IS THE FIRST STEP OF THE QZ ALGORITHM FOR SOLVING GENERALIZED MATRIX EIGENVALUE PROBLEMS, SIAM J. NUMER. ANAL. 10, 241-256(1973) BY MOLER AND STEWART.

EXAMPLE :

PROGRAM TQZHES(OUTPUT,TAPE6=OUTPUT)
DIMENSION A(5,5),Z(5,5),B(5,5)
LOGICAL MATZ

N = 5
NM = 5
MATZ = .TRUE.

DATA A /10.,2.,3.,2*1.,2.,12.,1.,2.,1.,3.,1.,11.,
* 1.,-1.,1.,2.,1.,9.,3*1.,-1.,1.,15. /

DATA B /12.,1.,-1.,2.,2*1.,14.,1.,-1.,1.,-1.,1.,
* 16.,-1.,1.,2.,-1.,-1.,12.,-1.,3*1.,-1.,11. /

CALL QZHES(NM,N,A,B,MATZ,Z)
WRITE(6,100) ((A(I,J),I=1,5),J=1,5),((B(I,J),I=1,5),J=1,5),

QZHES 18
QZHES 19
QZHES 20
QZHES 21
QZHES 22
QZHES 23
QZHES 24
QZHES 25
QZHES 26
QZHES 27
QZHES 28
QZHES 29
QZHES 30
QZHES 31
QZHES 32
QZHES 33
QZHES 34
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QZHES 68
QZHES 69
QZHES 70
QZHES 71
QZHES 72
QZHES 73
QZHES 74
QZHES 75
QZHES 76
QZHES 77

C	*	((Z(I,J),I=1,5),J=1,5)	QZHES	78			
C100	FORMAT(1H ,5H A = /5(1H ,5(G8.2,2X)/))		QZHES	79			
C	*	5H B = /5(1H ,5(G8.2,2X)/)	QZHES	80			
C	*	5H Z = /5(1H ,5(G8.2,2X)/)	QZHES	81			
C	STOP		QZHES	82			
C	END		QZHES	83			
C			QZHES	84			
C	OUTPUT :		QZHES	85			
C			QZHES	86			
C	A =		QZHES	87			
C	-9.9	4.1	0.	0.	0.	QZHES	88
C	-2.4	11.	-3.0	0.	0.	QZHES	89
C	.91	.26	-13.	3.3	0.	QZHES	90
C	-3.8	2.0	1.7	-11.	2.6	QZHES	91
C	2.7	-1.5	-.99	1.4	-11.	QZHES	92
C	B =					QZHES	93
C	-12.	0.	0.	0.	0.	QZHES	94
C	2.3	16.	0.	0.	0.	QZHES	95
C	-.34	-3.0	-12.	0.	0.	QZHES	96
C	-3.8	.80	-1.5	-10.	0.	QZHES	97
C	2.5	-1.4	-1.5	-1.5	-13.	QZHES	98
C	Z =					QZHES	99
C	1.0	0.	0.	0.	0.	QZHES	100
C	0.	.26	.95	-.14	-.70E-01	QZHES	101
C	0.	.87E-01	-.24E-01	.43	-.90	QZHES	102
C	0.	.24E-01	.16	.89	.43	QZHES	103
C	0.	-.96	.26	.22E-01	-.89E-01	QZHES	104
C						QZHES	105
C						QZHES	106
C	-----					EISP6685	
C	SUBROUTINE QZHES(NM,N,A,B,MATZ,Z)						
C	implicit real*8 (a-h,o-z)					EISP6686	
C	INTEGER I,J,K,L,N,LB,L1,NM,NK1,NM1,NM2					EISP6687	
C	REAL*8 A(NM,N),B(NM,N),Z(NM,N)					EISP66	
C	REAL*8 R,S,T,U1,U2,V1,V2,RHO					EISP66	
C	LOGICAL MATZ					EISP6690	
C	IF (.NOT. MATZ) GO TO 10					EISP6691	
C	DO 3 J = 1, N					EISP6692	
C	DO 2 I = 1, N					EISP6693	
C	Z(I,J) = 0.0E0					EISP6694	
C	2 CONTINUE					EISP6695	
C	Z(J,J) = 1.0E0					EISP6696	
C	3 CONTINUE					EISP6697	
C REDUCE B TO UPPER TRIANGULAR FORM					EISP6698	
C	10 IF (N .LE. 1) GO TO 170					EISP6699	
C	NM1 = N - 1					EISP6700	
C	DO 100 L = 1, NM1					EISP6701	
C	L1 = L + 1					EISP6702	
C	S = 0.0E0					EISP6703	
C	DO 20 I = L1, N					EISP6704	
C	S = S + ABS(B(I,L))					EISP6705	
C	20 CONTINUE					EISP6706	
C	IF (S .EQ. 0.0E0) GO TO 100					EISP6707	
C	S = S + ABS(B(L,L))					EISP6708	
C						EISP6709	
C						EISP6710	
C						EISP6711	
C						EISP6712	
C						EISP6713	
C						EISP6714	

-	R = 0.0E0	EISP6715
C		EISP6716
-	DO 25 I = L, N	EISP6717
	B(I,L) = B(I,L) / S	EISP6718
	R = R + B(I,L)**2	EISP6719
25	CONTINUE	EISP6720
-C		EISP6721
	R = SIGN(SQRT(R), B(L,L))	EISP6722
	B(L,L) = B(L,L) + R	EISP6723
	RHO = R * B(L,L)	EISP6724
-C		EISP6725
	DO 50 J = L1, N	EISP6726
	T = 0.0E0	EISP6727
-C		EISP6728
	DO 30 I = L, N	EISP6729
	T = T + B(I,L) * B(I,J)	EISP6730
30	CONTINUE	EISP6731
-C		EISP6732
	T = -T / RHO	EISP6733
C		EISP6734
-	DO 40 I = L, N	EISP6735
	B(I,J) = B(I,J) + T * B(I,L)	EISP6736
40	CONTINUE	EISP6737
-C		EISP6738
50	CONTINUE	EISP6739
C		EISP6740
-	DO 80 J = 1, N	EISP6741
	T = 0.0E0	EISP6742
C		EISP6743
	DO 60 I = L, N	EISP6744
	T = T + B(I,L) * A(I,J)	EISP6745
60	CONTINUE	EISP6746
C		EISP6747
-C		EISP6748
	T = -T / RHO	EISP6749
	DO 70 I = L, N	EISP6750
	A(I,J) = A(I,J) + T * B(I,L)	EISP6751
70	CONTINUE	EISP6752
C		EISP6753
80	CONTINUE	EISP6754
-C		EISP6755
	B(L,L) = -S * R	EISP6756
C		EISP6757
-	DO 90 I = L1, N	EISP6758
	B(I,L) = 0.0E0	EISP6759
90	CONTINUE	EISP6760
C		EISP6761
-	100 CONTINUE	EISP6762
C REDUCE A TO UPPER HESSENBERG FORM, WHILE	EISP6763
C	KEEPING B TRIANGULAR	EISP6764
-	IF (N .EQ. 2) GO TO 170	EISP6765
	NM2 = N - 2	EISP6766
C		EISP6767
-	DO 160 K = 1, NM2	EISP6768
	NK1 = NM1 - K	EISP6769
C FOR L=N-1 STEP -1 UNTIL K+1 DO --	EISP6770
	DO 150 LB = 1, NK1	EISP6771
-	L = N - LB	EISP6772
	L1 = L + 1	EISP6773
C ZERO A(L+1,K)	EISP6774
-		

-	S = ABS(A(L,K)) + ABS(A(L1,K))	EISP6775
-	IF (S .EQ. 0.0E0) GO TO 150	EISP6776
-	U1 = A(L,K) / S	EISP6777
-	U2 = A(L1,K) / S	EISP6778
-	R = SIGN(SQRT(U1*U1+U2*U2),U1)	EISP6779
-	V1 = -(U1 + R) / R	EISP6780
-	V2 = -U2 / R	EISP6781
-	U2 = V2 / V1	EISP6782
C		EISP6783
-	DO 110 J = K, N	EISP6784
-	T = A(L,J) + U2 * A(L1,J)	EISP6785
-	A(L,J) = A(L,J) + T * V1	EISP6786
-	A(L1,J) = A(L1,J) + T * V2	EISP6787
- 110	CONTINUE	EISP6788
C		EISP6789
-	A(L1,K) = 0.0E0	EISP6790
C		EISP6791
-	DO 120 J = L, N	EISP6792
-	T = B(L,J) + U2 * B(L1,J)	EISP6793
-	B(L,J) = B(L,J) + T * V1	EISP6794
-	B(L1,J) = B(L1,J) + T * V2	EISP6795
- 120	CONTINUE	EISP6796
C ZERO B(L+1,L)	EISP6797
-	S = ABS(B(L1,L1)) + ABS(B(L1,L))	EISP6798
-	IF (S .EQ. 0.0E0) GO TO 150	EISP6799
-	U1 = B(L1,L1) / S	EISP6800
-	U2 = B(L1,L) / S	EISP6801
-	R = SIGN(SQRT(U1*U1+U2*U2),U1)	EISP6802
-	V1 = -(U1 + R) / R	EISP6803
-	V2 = -U2 / R	EISP6804
-	U2 = V2 / V1	EISP6805
C		EISP6806
-	DO 130 I = 1, L1	EISP6807
-	T = B(I,L1) + U2 * B(I,L)	EISP6808
-	B(I,L1) = B(I,L1) + T * V1	EISP6809
-	B(I,L) = B(I,L) + T * V2	EISP6810
- 130	CONTINUE	EISP6811
-C		EISP6812
-	B(L1,L) = 0.0E0	EISP6813
C		EISP6814
-	DO 140 I = 1, N	EISP6815
-	T = A(I,L1) + U2 * A(I,L)	EISP6816
-	A(I,L1) = A(I,L1) + T * V1	EISP6817
-	A(I,L) = A(I,L) + T * V2	EISP6818
- 140	CONTINUE	EISP6819
C		EISP6820
-	IF (.NOT. MATZ) GO TO 150	EISP6821
-C		EISP6822
-	DO 145 I = 1, N	EISP6823
-	T = Z(I,L1) + U2 * Z(I,L)	EISP6824
-	Z(I,L1) = Z(I,L1) + T * V1	EISP6825
-	Z(I,L) = Z(I,L) + T * V2	EISP6826
- 145	CONTINUE	EISP6827
-C		EISP6828
- 150	CONTINUE	EISP6829
C		EISP6830
- 160	CONTINUE	EISP6831
-C		EISP6832
- 170	RETURN	EISP6833
-	C** THIS PROGRAM VALID ON FTN4 AND FTN5 **	EISP6834

END

ROUTINE NAME - PF261=QZIT
FROM EISPACK

EISP6835

QZIT 2

QZIT 3

QZIT 4

QZIT 5

QZIT 6

QZIT 7

QZIT 8

QZIT 9

QZIT 10

QZIT 11

QZIT 12

QZIT 13

QZIT 14

QZIT 15

QZIT 16

QZIT 17

QZIT 18

QZIT 19

QZIT 20

QZIT 21

QZIT 22

QZIT 23

QZIT 24

QZIT 25

QZIT 26

QZIT 27

QZIT 28

QZIT 29

QZIT 30

QZIT 31

QZIT 32

QZIT 33

QZIT 34

QZIT 35

QZIT 36

QZIT 37

QZIT 38

QZIT 39

QZIT 40

QZIT 41

QZIT 42

QZIT 43

QZIT 44

QZIT 45

QZIT 46

QZIT 47

QZIT 48

QZIT 49

QZIT 50

QZIT 51

QZIT 52

QZIT 53

QZIT 54

QZIT 55

QZIT 56

QZIT 57

QZIT 58

QZIT 59

QZIT 60

LATEST REVISION

- AUGUST 1, 1984
COMPUTER SCIENCES CORP., HAMPTON, VA.

PURPOSE

- THIS SUBROUTINE ACCEPTS A PAIR OF REAL
MATRICES, ONE OF THEM IN UPPER HESSENBERG
FORM AND THE OTHER IN UPPER TRIANGULAR FORM.
IT REDUCES THE HESSENBERG MATRIX TO
QUASI-TRIANGULAR FORM USING ORTHOGONAL
TRANSFORMATIONS WHILE MAINTAINING THE
TRIANGULAR FORM OF THE OTHER MATRIX. IT IS
USUALLY PRECEDED QZHES(PF260) AND FOLLOWED
BY QZVAL(PF262) AND, POSSIBLY, QZVEC(PF263).

USAGE

- CALL QZIT(NM,N,A,B,EPS1,MATZ,Z,IERR)

ARGUMENTS

NM

- ON INPUT NM MUST BE SET TO THE ROW DIMENSION
OF TWO-DIMENSIONAL ARRAY PARAMETERS AS
DECLARED IN THE CALLING PROGRAM DIMENSION
STATEMENT.

N

- ON INPUT N IS THE ORDER OF THE MATRICES.

A

- ON INPUT A CONTAINS A REAL UPPER HESSENBERG
MATRIX.
MUST BE OF DIMENSION NM X N.

ON OUTPUT A HAS BEEN REDUCED TO
QUASI-TRIANGULAR FORM. THE ELEMENTS BELOW THE
FIRST SUBDIAGONAL ARE STILL ZERO AND NO TWO
CONSECUTIVE SUBDIAGONAL ELEMENTS ARE NONZERO.

B

- ON INPUT B CONTAINS A REAL UPPER TRIANGULAR
MATRIX.
MUST BE OF DIMENSION NM X N.

ON OUTPUT B IS STILL IN UPPER TRIANGULAR
FORM, ALTHOUGH ITS ELEMENTS HAVE BEEN ALTERED.
THE LOCATION B(N,1) IS USED TO STORE EPS1
TIMES THE NORM OF B FOR LATER USE BY QZVAL
QZVAL(PF262) AND QZVEC(PF263).

EPS1

- ON INPUT EPS1 IS A TOLERANCE USED TO DETERMINE
NEGLECTIBLE ELEMENTS. EPS1 = 0.0 (OR NEGATIVE)
MAY BE INPUT, IN WHICH CASE AN ELEMENT WILL BE
NEGLECTED ONLY IF IT IS LESS THAN ROUNDOFF
ERROR TIMES THE NORM OF ITS MATRIX. IF THE
INPUT EPS1 IS POSITIVE, THEN AN ELEMENT WILL
BE CONSIDERED NEGLECTIBLE IF IT IS LESS THAN
EPS1 TIMES THE NORM OF ITS MATRIX. A POSITIVE
VALUE OF EPS1 MAY RESULT IN FASTER EXECUTION,
BUT LESS ACCURATE RESULTS.

```

MATZ      - ON INPUT MATZ SHOULD BE SET TO .TRUE.  IF THE QZIT 61
           RIGHT HAND TRANSFORMATIONS ARE TO BE QZIT 62
           ACCUMULATED FOR LATER USE IN COMPUTING QZIT 63
           EIGENVECTORS, AND TO .FALSE.  OTHERWISE. QZIT 64
           QZIT 65
           QZIT 66
Z          - ON INPUT Z CONTAINS, IF MATZ HAS BEEN SET TO QZIT 67
           .TRUE., THE TRANSFORMATION MATRIX PRODUCED IN QZIT 68
           THE REDUCTION BY QZHS(PF260), IF PERFORMED, QZIT 69
           OR ELSE THE IDENTITY MATRIX.  IF MATZ HAS BEEN QZIT 70
           SET TO .FALSE., Z IS NOT REFERENCED. QZIT 71
           MUST BE OF DIMENSION NM X N. QZIT 72
           QZIT 73
           ON OUTPUT Z CONTAINS THE PRODUCT OF THE QZIT 74
           RIGHT HAND TRANSFORMATIONS (FOR BOTH STEPS) IF QZIT 75
           MATZ HAS BEEN SET TO .TRUE.. QZIT 76
           QZIT 77
IERR      - ON OUTPUT IERR IS SET TO QZIT 78
           ZERO FOR NORMAL RETURN. QZIT 79
           J IF THE LIMIT OF 30*N ITERATIONS IS EXHAUSTED QZIT 80
           WHILE THE J-TH EIGENVALUE IS BEING SOUGHT. QZIT 81

```

REQUIRED ROUTINES - HC318=EPSLON

REMARKS 1. THIS SUBROUTINE IS THE SECOND STEP OF THE QZ ALGORITHM FOR SOLVING GENERALIZED MATRIX EIGENVALUE PROBLEMS, SIAM J. NUMER. ANAL. 10, 241-256(1973) BY MOLER AND STEWART, AS MODIFIED IN TECHNICAL NOTE NASA TN D-7305(1973) BY WARD.

EXAMPLE :

```
PROGRAM TQZIT(OUTPUT,TAPE6=OUTPUT)
DIMENSION A(5,5),B(5,5),Z(5,5)
LOGICAL MATZ
```

```
N = 5
NM = 5
MATZ = .TRUE.
EPS1 = 0.0E0
```

```
DATA A /10.,2.,3.,2*1.,2.,12.,1.,2.,1.,3.,1.,11.,
*      1.,-1.,1.,2.,1.,9.,3*1.,-1.,1.,15. /
```

```
DATA B /12.,1.,-1.,2.,2*1.,14.,1.,-1.,1.,-1.,1.,
*      16.,-1.,1.,2.,-1.,-1.,12.,-1.,3*1.,-1.,11. /
```

```
CALL QZHE5 (NM,N,A,B,MATZ,Z)
CALL QZIT (NM,N,A,B,EPS1,MATZ,Z,IERR)
```

```
WRITE(6,99) IERR
```

```
FORMAT(1H1,8H IERR = ,I4)
```

```
WRITE(6,100) ((A(I,J),I=1,5),J=1,5),((B(I,J),I=1,5),J=1,5),
*          ((Z(I,J),I=1,5),J=1,5)
```

```
FORMAT(1H ,5H A = /5(1H ,5(G8.2,2X)/))
```

$$5H \ B = /5(1H \ , 5(G8.2, 2X) /)$$
$$5H \ Z = /5(1H \ , 5(G8.2, 2X) /))$$

STOP
END

C	OUTPUT :	QZIT 121
C		QZIT 122
C	IERR = 0	QZIT 123
-C	A =	QZIT 124
C	-15. -1.3 0. 0. 0.	QZIT 125
C	1.1 7.4 0. 0. 0.	QZIT 126
-C	1.5 -1.5 -16. 0. 0.	QZIT 127
C	-2.2 .96 1.0 -10. 0.	QZIT 128
C	-2.6 -.31 1.2 1.7 -8.6	QZIT 129
C	B =	QZIT 130
-C	-9.9 0. 0. 0. .31E-12	QZIT 131
C	-.29 17. 0. 0. 0.	QZIT 132
C	1.3 -2.1 -14. 0. 0.	QZIT 133
-C	-1.9 1.7 .96 -11. 0.	QZIT 134
C	-2.6 -.32 1.3 2.1 -13.	QZIT 135
C	Z =	QZIT 136
-C	.28 -.71E-01 .16 -.24 -.91	QZIT 137
C	.52 -.24 -.66 .48 -.64E-01	QZIT 138
C	.49 .56 .49 .45 .75E-01	QZIT 139
C	-.60 .48 -.29 .44 -.38	QZIT 140
-C	-.25 -.63 .45 .57 -.94E-01	QZIT 141
C		QZIT 142
C		QZIT 143
-C	-----	EISP6836
C	SUBROUTINE QZIT(NM,N,A,B,EPS1,MATZ,Z,IERR)	
-C		EISP6837
C	implicit real*8 (a-h,o-z)	EISP6838
-	INTEGER I,J,K,L,N,EN,K1,K2,LD,LL,L1,NA,NM,ISH,ITN,ITS,KM1,LM1,	EISP6839
X	ENM2,IERR,LOR1,ENORN	EISP68
C	REAL*8 A(NM,N),B(NM,N),Z(NM,N)	EISP68
-	REAL*8 R,S,T,A1,A2,A3,EP,SH,U1,U2,U3,V1,V2,V3,ANI,A11,	EISP6842
X	A12,A21,A22,A33,A34,A43,A44,BNI,B11,B12,B22,B33,B34,	EISP6843
X	B44,EPSA,EPSB,EPS1,ANORM,BNORM,EPSLON	EISP6844
C	LOGICAL MATZ,NOTLAS	EISP6845
-C	IERR = 0	EISP6846
C COMPUTE EPSA,EPSB	EISP6847
-	ANORM = 0.0E0	EISP6848
-C	BNORM = 0.0E0	EISP6849
C		EISP6850
-	DO 30 I = 1, N	EISP6851
C	ANI = 0.0E0	EISP6852
-	IF (I .NE. 1) ANI = ABS(A(I,I-1))	EISP6853
C	BNI = 0.0E0	EISP6854
-		EISP6855
C	DO 20 J = I, N	EISP6856
-	ANI = ANI + ABS(A(I,J))	EISP6857
C	BNI = BNI + ABS(B(I,J))	EISP6858
-C	20 CONTINUE	EISP6859
C		EISP6860
-	IF (ANI .GT. ANORM) ANORM = ANI	EISP6861
C	IF (BNI .GT. BNORM) BNORM = BNI	EISP6862
-C	30 CONTINUE	EISP6863
C		EISP6864
-	IF (ANORM .EQ. 0.0E0) ANORM = 1.0E0	EISP6865
C	IF (BNORM .EQ. 0.0E0) BNORM = 1.0E0	EISP6866
-	EP = EPS1	EISP6867
C	IF (EP .GT. 0.0E0) GO TO 50	EISP6868
- USE ROUNDOFF LEVEL IF EPS1 IS ZERO	EISP6869
C	EP = EPSLON(1.0E0)	EISP6870
-C	50 EPSA = EP * ANORM	EISP6871
C	EPSB = EP * BNORM	

-C REDUCE A TO QUASI-TRIANGULAR FORM, WHILE	EISP6872
C	KEEPING B TRIANGULAR	EISP6873
	LOR1 = 1	EISP6874
-	ENORN = N	EISP6875
	EN = N	EISP6876
	ITN = 30*N	EISP6877
-C BEGIN QZ STEP	EISP6878
	60 IF (EN .LE. 2) GO TO 1001	EISP6879
	IF (.NOT. MATZ) ENORN = EN	EISP6880
	ITS = 0	EISP6881
-	NA = EN - 1	EISP6882
	ENM2 = NA - 1	EISP6883
	70 ISH = 2	EISP6884
-C CHECK FOR CONVERGENCE OR REDUCIBILITY.	EISP6885
C	FOR L=EN STEP -1 UNTIL 1 DO --	EISP6886
	DO 80 LL = 1, EN	EISP6887
-	LM1 = EN - LL	EISP6888
	L = LM1 + 1	EISP6889
	IF (L .EQ. 1) GO TO 95	EISP6890
	IF (ABS(A(L,LM1)) .LE. EPSA) GO TO 90	EISP6891
-	80 CONTINUE	EISP6892
C		EISP6893
	90 A(L,LM1) = 0.0E0	EISP6894
-	IF (L .LT. NA) GO TO 95	EISP6895
C 1-BY-1 OR 2-BY-2 BLOCK ISOLATED	EISP6896
	EN = LM1	EISP6897
	GO TO 60	EISP6898
-C CHECK FOR SMALL TOP OF B	EISP6899
	95 LD = L	EISP6900
	100 L1 = L + 1	EISP6901
-	B11 = B(L,L)	EISP6902
	IF (ABS(B11) .GT. EPSB) GO TO 120	EISP6903
	B(L,L) = 0.0E0	EISP6904
-	S = ABS(A(L,L)) + ABS(A(L1,L))	EISP6905
	U1 = A(L,L) / S	EISP6906
	U2 = A(L1,L) / S	EISP6907
-	R = SIGN(SQRT(U1*U1+U2*U2),U1)	EISP6908
	V1 = -(U1 + R) / R	EISP6909
	V2 = -U2 / R	EISP6910
	U2 = V2 / V1	EISP6911
-C		EISP6912
	DO 110 J = L, ENORN	EISP6913
	T = A(L,J) + U2 * A(L1,J)	EISP6914
-	A(L,J) = A(L,J) + T * V1	EISP6915
	A(L1,J) = A(L1,J) + T * V2	EISP6916
	T = B(L,J) + U2 * B(L1,J)	EISP6917
-	B(L,J) = B(L,J) + T * V1	EISP6918
	B(L1,J) = B(L1,J) + T * V2	EISP6919
	110 CONTINUE	EISP6920
C		EISP6921
-	IF (L .NE. 1) A(L,LM1) = -A(L,LM1)	EISP6922
	LM1 = L	EISP6923
	L = L1	EISP6924
-	GO TO 90	EISP6925
	120 A11 = A(L,L) / B11	EISP6926
	A21 = A(L1,L) / B11	EISP6927
	IF (ISH .EQ. 1) GO TO 140	EISP6928
-C ITERATION STRATEGY	EISP6929
	IF (ITN .EQ. 0) GO TO 1000	EISP6930
	IF (ITS .EQ. 10) GO TO 155	EISP6931

C DETERMINE TYPE OF SHIFT	EISP6932
	B22 = B(L1,L1)	EISP6933
	IF (ABS(B22) .LT. EPSB) B22 = EPSB	EISP6934
-	B33 = B(NA,NA)	EISP6935
	IF (ABS(B33) .LT. EPSB) B33 = EPSB	EISP6936
	B44 = B(EN,EN)	EISP6937
-	IF (ABS(B44) .LT. EPSB) B44 = EPSB	EISP6938
	A33 = A(NA,NA) / B33	EISP6939
	A34 = A(NA,EN) / B44	EISP6940
	A43 = A(EN,NA) / B33	EISP6941
-	A44 = A(EN,EN) / B44	EISP6942
	B34 = B(NA,EN) / B44	EISP6943
	T = 0.5E0 * (A43 * B34 - A33 - A44)	EISP6944
-	R = T * T + A34 * A43 - A33 * A44	EISP6945
	IF (R .LT. 0.0E0) GO TO 150	EISP6946
C DETERMINE SINGLE SHIFT ZEROth COLUMN OF A	EISP6947
	ISH = 1	EISP6948
-	R = SQRT(R)	EISP6949
	SH = -T + R	EISP6950
	S = -T - R	EISP6951
-	IF (ABS(S-A44) .LT. ABS(SH-A44)) SH = S	EISP6952
C LOOK FOR TWO CONSECUTIVE SMALL	EISP6953
C	SUB-DIAGONAL ELEMENTS OF A.	EISP6954
-C	FOR L=EN-2 STEP -1 UNTIL LD DO --	EISP6955
	DO 130 LL = LD, ENM2	EISP6956
	L = ENM2 + LD - LL	EISP6957
-	IF (L .EQ. LD) GO TO 140	EISP6958
	LM1 = L - 1	EISP6959
	L1 = L + 1	EISP6960
-	T = A(L,L)	EISP6961
	IF (ABS(B(L,L)) .GT. EPSB) T = T - SH * B(L,L)	EISP6962
	IF (ABS(A(L,LM1)) .LE. ABS(T/A(L1,L)) * EPSA) GO TO 100	EISP6963
	130 CONTINUE	EISP6964
-C		EISP6965
	140 A1 = A11 - SH	EISP6966
	A2 = A21	EISP6967
-	IF (L .NE. LD) A(L,LM1) = -A(L,LM1)	EISP6968
	GO TO 160	EISP6969
C DETERMINE DOUBLE SHIFT ZEROth COLUMN OF A	EISP6970
-	150 A12 = A(L,L1) / B22	EISP6971
	A22 = A(L1,L1) / B22	EISP6972
	B12 = B(L,L1) / B22	EISP6973
	A1 = ((A33 - A11) * (A44 - A11) - A34 * A43 + A43 * B34 * A11)	EISP6974
-	X / A21 + A12 - A11 * B12	EISP6975
	A2 = (A22 - A11) - A21 * B12 - (A33 - A11) - (A44 - A11)	EISP6976
-	X + A43 * B34	EISP6977
	A3 = A(L1+1,L1) / B22	EISP6978
-	GO TO 160	EISP6979
C AD HOC SHIFT	EISP6980
	155 A1 = 0.0E0	EISP6981
-	A2 = 1.0E0	EISP6982
	A3 = 1.1605E0	EISP6983
	160 ITS = ITS + 1	EISP6984
	ITN = ITN - 1	EISP6985
-	IF (.NOT. MATZ) LOR1 = LD	EISP6986
C MAIN LOOP	EISP6987
	DO 260 K = L, NA	EISP6988
-	NOTLAS = K .NE. NA .AND. ISH .EQ. 2	EISP6989
	K1 = K + 1	EISP6990
	K2 = K + 2	EISP6991

-	KM1 = MAX0(K-1,L)	EISP6992
	LL = MIN0(EN,K1+ISH)	EISP6993
	IF (NOTLAS) GO TO 190	EISP6994
-C ZERO A(K+1,K-1)	EISP6995
	IF (K .EQ. L) GO TO 170	EISP6996
	A1 = A(K,KM1)	EISP6997
	A2 = A(K1,KM1)	EISP6998
-	170 S = ABS(A1) + ABS(A2)	EISP6999
	IF (S .EQ. 0.0E0) GO TO 70	EISP7000
	U1 = A1 / S	EISP7001
-	U2 = A2 / S	EISP7002
	R = SIGN(SQRT(U1*U1+U2*U2),U1)	EISP7003
	V1 = -(U1 + R) / R	EISP7004
	V2 = -U2 / R	EISP7005
-	U2 = V2 / V1	EISP7006
C		EISP7007
	DO 180 J = KM1, ENORN	EISP7008
-	T = A(K,J) + U2 * A(K1,J)	EISP7009
	A(K,J) = A(K,J) + T * V1	EISP7010
	A(K1,J) = A(K1,J) + T * V2	EISP7011
-	T = B(K,J) + U2 * B(K1,J)	EISP7012
	B(K,J) = B(K,J) + T * V1	EISP7013
	B(K1,J) = B(K1,J) + T * V2	EISP7014
-	180 CONTINUE	EISP7015
-C		EISP7016
	IF (K .NE. L) A(K1,KM1) = 0.0E0	EISP7017
	GO TO 240	EISP7018
-C ZERO A(K+1,K-1) AND A(K+2,K-1)	EISP7019
-	190 IF (K .EQ. L) GO TO 200	EISP7020
	A1 = A(K,KM1)	EISP7021
	A2 = A(K1,KM1)	EISP7022
	A3 = A(K2,KM1)	EISP7023
-	200 S = ABS(A1) + ABS(A2) + ABS(A3)	EISP7024
	IF (S .EQ. 0.0E0) GO TO 260	EISP7025
	U1 = A1 / S	EISP7026
	U2 = A2 / S	EISP7027
	U3 = A3 / S	EISP7028
-	R = SIGN(SQRT(U1*U1+U2*U2+U3*U3),U1)	EISP7029
	V1 = -(U1 + R) / R	EISP7030
	V2 = -U2 / R	EISP7031
	V3 = -U3 / R	EISP7032
-	U2 = V2 / V1	EISP7033
	U3 = V3 / V1	EISP7034
C		EISP7035
-	DO 210 J = KM1, ENORN	EISP7036
	T = A(K,J) + U2 * A(K1,J) + U3 * A(K2,J)	EISP7037
	A(K,J) = A(K,J) + T * V1	EISP7038
-	A(K1,J) = A(K1,J) + T * V2	EISP7039
	A(K2,J) = A(K2,J) + T * V3	EISP7040
	T = B(K,J) + U2 * B(K1,J) + U3 * B(K2,J)	EISP7041
	B(K,J) = B(K,J) + T * V1	EISP7042
-	B(K1,J) = B(K1,J) + T * V2	EISP7043
	B(K2,J) = B(K2,J) + T * V3	EISP7044
-	210 CONTINUE	EISP7045
-C		EISP7046
	IF (K .EQ. L) GO TO 220	EISP7047
	A(K1,KM1) = 0.0E0	EISP7048
	A(K2,KM1) = 0.0E0	EISP7049
- ZERO B(K+2,K+1) AND B(K+2,K)	EISP7050
C		EISP7051
-	220 S = ABS(B(K2,K2)) + ABS(B(K2,K1)) + ABS(B(K2,K))	

-	IF (S .EQ. 0.0E0) GO TO 240	EISP7052
-	U1 = B(K2,K2) / S	EISP7053
-	U2 = B(K2,K1) / S	EISP7054
-	U3 = B(K2,K) / S	EISP7055
-	R = SIGN(SQRT(U1*U1+U2*U2+U3*U3),U1)	EISP7056
-	V1 = -(U1 + R) / R	EISP7057
-	V2 = -U2 / R	EISP7058
-	V3 = -U3 / R	EISP7059
-	U2 = V2 / V1	EISP7060
-	U3 = V3 / V1	EISP7061
-C	DO 230 I = LOR1, LL	EISP7062
-	T = A(I,K2) + U2 * A(I,K1) + U3 * A(I,K)	EISP7063
-	A(I,K2) = A(I,K2) + T * V1	EISP7064
-	A(I,K1) = A(I,K1) + T * V2	EISP7065
-	A(I,K) = A(I,K) + T * V3	EISP7066
-	T = B(I,K2) + U2 * B(I,K1) + U3 * B(I,K)	EISP7067
-	B(I,K2) = B(I,K2) + T * V1	EISP7068
-	B(I,K1) = B(I,K1) + T * V2	EISP7069
-	B(I,K) = B(I,K) + T * V3	EISP7070
-	230 CONTINUE	EISP7071
-C	B(K2,K) = 0.0E0	EISP7072
-	B(K2,K1) = 0.0E0	EISP7073
-	IF (.NOT. MATZ) GO TO 240	EISP7074
-C	DO 235 I = 1, N	EISP7075
-	T = Z(I,K2) + U2 * Z(I,K1) + U3 * Z(I,K)	EISP7076
-	Z(I,K2) = Z(I,K2) + T * V1	EISP7077
-	Z(I,K1) = Z(I,K1) + T * V2	EISP7078
-	Z(I,K) = Z(I,K) + T * V3	EISP7079
-	235 CONTINUE	EISP7080
-C ZERO B(K+1,K)	EISP7081
-	240 S = ABS(B(K1,K1)) + ABS(B(K1,K))	EISP7082
-	IF (S .EQ. 0.0E0) GO TO 260	EISP7083
-	U1 = B(K1,K1) / S	EISP7084
-	U2 = B(K1,K) / S	EISP7085
-	R = SIGN(SQRT(U1*U1+U2*U2),U1)	EISP7086
-	V1 = -(U1 + R) / R	EISP7087
-	V2 = -U2 / R	EISP7088
-	U2 = V2 / V1	EISP7089
-C	DO 250 I = LOR1, LL	EISP7090
-	T = A(I,K1) + U2 * A(I,K)	EISP7091
-	A(I,K1) = A(I,K1) + T * V1	EISP7092
-	A(I,K) = A(I,K) + T * V2	EISP7093
-	T = B(I,K1) + U2 * B(I,K)	EISP7094
-	B(I,K1) = B(I,K1) + T * V1	EISP7095
-	B(I,K) = B(I,K) + T * V2	EISP7096
-	250 CONTINUE	EISP7097
-C	B(K1,K) = 0.0E0	EISP7098
-	IF (.NOT. MATZ) GO TO 260	EISP7099
-C	DO 255 I = 1, N	EISP7100
-	T = Z(I,K1) + U2 * Z(I,K)	EISP7101
-	Z(I,K1) = Z(I,K1) + T * V1	EISP7102
-	Z(I,K) = Z(I,K) + T * V2	EISP7103
-	255 CONTINUE	EISP7104
-C		EISP7105
-		EISP7106
-		EISP7107
-		EISP7108
-		EISP7109
-		EISP7110
-		EISP7111

-	260	CONTINUE		EISP7112
C	 END QZ STEP		EISP7113
		GO TO 70		EISP7114
-C	 SET ERROR -- ALL EIGENVALUES HAVE NOT		EISP7115
C		CONVERGED AFTER 30*N ITERATIONS		EISP7116
	1000	IERR = EN		EISP7117
-C	 SAVE EPSB FOR USE BY QZVAL AND QZVEC		EISP7118
-	1001	IF (N .GT. 1) B(N,1) = EPSB		EISP7119
		RETURN		EISP7120
C**		THIS PROGRAM VALID ON FTN4 AND FTN5 **		EISP7121
-		END		EISP7122
C		ROUTINE NAME - PF262=QZVAL	QZVAL	2
C		FROM EISPACK	QZVAL	3
-C			QZVAL	4
C		-----	QZVAL	5
C			QZVAL	6
C	LATEST REVISION	- AUGUST 1,1984	QZVAL	7
-C		COMPUTER SCIENCES CORP., HAMPTON, VA.	QZVAL	8
C			QZVAL	9
C			QZVAL	10
-C	PURPOSE	- THIS SUBROUTINE ACCEPTS A PAIR OF REAL	QZVAL	11
C		MATRICES, ONE OF THEM IN QUASI-TRIANGULAR	QZVAL	12
C		FORM AND THE OTHER IN UPPER TRIANGULAR FORM.	QZVAL	13
-C		IT REDUCES THE QUASI-TRIANGULAR MATRIX	QZVAL	14
C		FURTHER, SO THAT ANY REMAINING 2-BY-2 BLOCKS	QZVAL	15
C		CORRESPOND TO PAIRS OF COMPLEX EIGENVALUES,	QZVAL	16
-C		AND RETURNS QUANTITIES WHOSE RATIOS GIVE THE	QZVAL	17
C		GENERALIZED EIGENVALUES. IT IS USUALLY	QZVAL	18
C		PRECEDED BY QZHES(PF260) AND QZIT(PF261) AND	QZVAL	19
-C		MAY BE FOLLOWED BY QZVEC(PF263).	QZVAL	20
C			QZVAL	21
C			QZVAL	22
C	USAGE	- CALL QZVAL(NM,N,A,B,ALFR,ALFI,BETA,MATZ,Z)	QZVAL	23
-C			QZVAL	24
C	ARGUMENTS	NM	QZVAL	25
C		- ON INPUT NM MUST BE SET TO THE ROW DIMENSION	QZVAL	26
-C		OF TWO-DIMENSIONAL ARRAY PARAMETERS AS	QZVAL	27
C		DECLARED IN THE CALLING PROGRAM DIMENSION	QZVAL	28
C		STATEMENT.	QZVAL	29
-C		N	QZVAL	30
C		- ON INPUT N IS THE ORDER OF THE MATRICES.	QZVAL	31
C		A	QZVAL	32
-C		- ON INPUT A CONTAINS A REAL UPPER QUASI-	QZVAL	33
C		TRIANGULAR MATRIX.	QZVAL	34
C		MUST BE OF DIMENSION NM X N.	QZVAL	35
-C		ON OUTPUT A HAS BEEN REDUCED FURTHER TO A	QZVAL	36
C		QUASI-TRIANGULAR MATRIX IN WHICH ALL NONZERO	QZVAL	37
-C		SUBDIAGONAL ELEMENTS CORRESPOND TO PAIRS OF	QZVAL	38
C		COMPLEX EIGENVALUES.	QZVAL	39
C			QZVAL	40
-C		B	QZVAL	41
C		- ON INPUT B CONTAINS A REAL UPPER TRIANGULAR	QZVAL	42
C		MATRIX.	QZVAL	43
-C		MUST BE OF DIMENSION NM X N.	QZVAL	44
C		IN ADDITION, LOCATION B(N,1) CONTAINS THE	QZVAL	45
-C		TOLERANCE QUANTITY (EPSB) COMPUTED AND SAVED	QZVAL	46
C		IN QZIT(PF261).	QZVAL	47
-C		ON OUTPUT B IS STILL IN UPPER TRIANGULAR	QZVAL	48
C		FORM, ALTHOUGH ITS ELEMENTS HAVE BEEN ALTERED.	QZVAL	49
C		B(N,1) IS UNALTERED.	QZVAL	50

		QZVAL 51
		QZVAL 52
ALFR	- ON OUTPUT ALFR CONTAINS THE REAL PART OF THE	QZVAL 53
	DIAGONAL ELEMENTS OF THE TRIANGULAR MATRIX	QZVAL 54
	THAT WOULD BE OBTAINED IF A WERE REDUCED	QZVAL 55
	COMPLETELY TO TRIANGULAR FORM BY UNITARY	QZVAL 56
	TRANSFORMATIONS. NON-ZERO VALUES OF ALFI	QZVAL 57
	OCCUR IN PAIRS, THE FIRST MEMBER POSITIVE AND	QZVAL 58
	THE SECOND NEGATIVE.	QZVAL 59
	MUST BE OF DIMENSION N.	QZVAL 60
		QZVAL 61
ALFI	- ON OUTPUT ALFI CONTAINS THE IMAGINARY PART	QZVAL 62
	OF THE DIAGONAL ELEMENTS OF OF THE TRIANGULAR	QZVAL 63
	MATRIX THAT WOULD BE OBTAINED IF A WERE	QZVAL 64
	REDUCED COMPLETELY TO TRIANGULAR FORM BY	QZVAL 65
	UNITARY TRANSFORMATIONS. NON-ZERO VALUES	QZVAL 66
	OF ALFI OCCUR IN PAIRS, THE FIRST MEMBER	QZVAL 67
	POSITIVE AND THE SECOND NEGATIVE.	QZVAL 68
	MUST BE OF DIMENSION N.	QZVAL 69
		QZVAL 70
BETA	- ON OUTPUT BETA CONTAINS THE DIAGONAL ELEMENTS	QZVAL 71
	OF THE CORRESPONDING B, NORMALIZED TO BE REAL	QZVAL 72
	AND NON-NEGATIVE. THE GENERALIZED EIGENVALUES	QZVAL 73
	ARE THEN THE RATIOS ((ALFR+I*ALFI)/BETA).	QZVAL 74
	MUST BE OF DIMENSION N.	QZVAL 75
		QZVAL 76
		QZVAL 77
MATZ	- ON INPUT MATZ SHOULD BE SET TO .TRUE. IF	QZVAL 78
	THE RIGHT HAND TRANSFORMATIONS ARE TO BE	QZVAL 79
	ACCUMULATED FOR LATER USE IN COMPUTING	QZVAL 80
	EIGENVECTORS, AND TO .FALSE. OTHERWISE.	QZVAL 81
		QZVAL 82
Z	- ON INPUT Z CONTAINS, IF MATZ HAS BEEN SET	QZVAL 83
	TO .TRUE., THE TRANSFORMATION MATRIX PRODUCED	QZVAL 84
	IN THE REDUCTIONS BY QZHES(PF260) AND QZIT	QZVAL 85
	(PF261) IF PERFORMED, OR ELSE THE IDENTITY	QZVAL 86
	MATRIX. IF MATZ HAS BEEN SET TO .FALSE., Z	QZVAL 87
	IS NOT REFERENCED.	QZVAL 88
	MUST BE OF DIMENSION NM X N.	QZVAL 89
		QZVAL 90
	ON OUTPUT Z CONTAINS THE PRODUCT OF THE	QZVAL 91
	RIGHT HAND TRANSFORMATIONS (FOR ALL THREE	QZVAL 92
	STEPS) IF MATZ HAS BEEN SET TO .TRUE.	QZVAL 93
		QZVAL 94
REQUIRED ROUTINES	- NONE	QZVAL 95
		QZVAL 96
REMARKS	1. THIS SUBROUTINE IS THE THIRD STEP OF THE QZ	QZVAL 97
	ALGORITHM FOR SOLVING GENERALIZED MATRIX	QZVAL 98
	EIGENVALUE PROBLEMS, SIAM J. NUMER. ANAL. 10,	QZVAL 99
	241-256(1973) BY MOLER AND STEWART.	QZVAL100
EXAMPLE :		QZVAL101
	PROGRAM TQZVAL(OUTPUT,TAPE6=OUTPUT)	QZVAL102
	DIMENSION A(5,5),B(5,5),ALFR(5),ALFI(5),BETA(5),Z(5,5)	QZVAL103
	LOGICAL MATZ	QZVAL104
		QZVAL105
		QZVAL106
	N = 5	QZVAL107
	NM = 5	QZVAL108
	MATZ = .TRUE.	QZVAL109
	EPS1 = 0.0E0	QZVAL110

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-C
C      DATA A /10.,2.,3.,2*1.,2.,12.,1.,2.,1.,3.,1.,11.,
C      *      1.,-1.,1.,2.,1.,9.,3*1.,-1.,1.,15. /
-C
C      DATA B /12.,1.,-1.,2.,2*1.,14.,1.,-1.,1.,-1.,1.,
C      *      16.,-1.,1.,2.,-1.,-1.,12.,-1.,3*1.,-1.,11. /
-C
C      CALL QZHES(NM,N,A,B,MATZ,Z)
C      CALL QZIT(NM,N,A,B,EPS1,MATZ,Z,IERR)
C      CALL QZVAL(NM,N,A,B,ALFR,ALFI,BETA,MATZ,Z)
-C
C      WRITE(6,99) IERR
C      WRITE(6,100) ALFR,ALFI,BETA,((Z(I,J),I=1,5),J=1,5)
C99  FORMAT(1H1,8H IERR = ,I4)
-C100 FORMAT(1H ,8H ALFR = /1H ,5(G8.2,2X) /
C      *      8H ALFI = /1H ,5(G8.2,2X) /
C      *      8H BETA = /1H ,5(G8.2,2X) /
C      *      5H Z = /5(1H ,5(G8.2,2X) /))
-C
C      STOP
C      END
C
-C      OUTPUT :
C
C      IERR =      0
-C      ALFR =
C      15.          7.2          16.          10.          8.6
C      ALFI =
C      0.           0.           0.           0.           0.
-C      BETA =
C      9.9          17.          14.          11.          13.
C      Z =
-C      .24          -.54E-01   .21          -.27          -.91
C      -.54          .25         .65          -.46          .13
C      .49          .56         .49          .45          .75E-01
-C      -.60          .48         -.29         .44          -.38
C      -.25          -.63         .45         .57          -.94E-01
C
C-----
C      SUBROUTINE QZVAL(NM,N,A,B,ALFR,ALFI,BETA,MATZ,Z)
C
C      implicit real*8 (a-h,o-z)
C      INTEGER I,J,N,EN,NA,NM,NN,ISW
C      REAL*8 A(NM,N),B(NM,N),ALFR(N),ALFI(N),BETA(N),Z(NM,N)
C      REAL*8 C,D,E,R,S,T,AN,A1,A2,BN,CQ,CZ,DI,DR,EI,TI,TR,U1,
X      U2,V1,V2,A1I,A1I,A12,A2I,A21,A22,B11,B12,B22,SQI,SQR,
X      SSI,SSR,SZI,SZR,A11I,A11R,A12I,A12R,A22I,A22R,EPSB
C      LOGICAL MATZ
C      EPSB = B(N,1)
C      ISW = 1
C      ..... FIND EIGENVALUES OF QUASI-TRIANGULAR MATRICES.
C      FOR EN=N STEP -1 UNTIL 1 DO -- .....
C
C      DO 510 NN = 1, N
C          EN = N + 1 - NN
C          NA = EN - 1
C          IF (ISW .EQ. 2) GO TO 505
C          IF (EN .EQ. 1) GO TO 410
C          IF (A(EN,NA) .NE. 0.0E0) GO TO 420
C      ..... 1-BY-1 BLOCK, ONE REAL ROOT .....
C      410  ALFR(EN) = A(EN,EN)
C          IF (B(EN,EN) .LT. 0.0E0) ALFR(EN) = -ALFR(EN)
C          BETA(EN) = ABS(B(EN,EN))

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QZVAL111
 QZVAL112
 QZVAL113
 QZVAL114
 QZVAL115
 QZVAL116
 QZVAL117
 QZVAL118
 QZVAL119
 QZVAL120
 QZVAL121
 QZVAL122
 QZVAL123
 QZVAL124
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 QZVAL144
 QZVAL145
 QZVAL146
 QZVAL147
 EISP7123
 EISP7124
 EISP7125
 EISP7126
 EISP7127
 EISP7128
 EISP7129
 EISP7130
 EISP7131
 EISP7132
 EISP7133
 EISP7134
 EISP7135
 EISP7136
 EISP7137
 EISP7138
 EISP7139
 EISP7140
 EISP7141
 EISP7142
 EISP7143
 EISP7144

-		ALFI(EN) = 0.0E0	EISP7145
		GO TO 510	EISP7146
C	 2-BY-2 BLOCK	EISP7147
-	420	IF (ABS(B(NA,NA)) .LE. EPSB) GO TO 455	EISP7148
		IF (ABS(B(EN,EN)) .GT. EPSB) GO TO 430	EISP7149
		A1 = A(EN,EN)	EISP7150
		A2 = A(EN,NA)	EISP7151
		BN = 0.0E0	EISP7152
		GO TO 435	EISP7153
	430	AN = ABS(A(NA,NA)) + ABS(A(NA,EN)) + ABS(A(EN,NA))	EISP7154
-	X	+ ABS(A(EN,EN))	EISP7155
		BN = ABS(B(NA,NA)) + ABS(B(NA,EN)) + ABS(B(EN,EN))	EISP7156
		A11 = A(NA,NA) / AN	EISP7157
		A12 = A(NA,EN) / AN	EISP7158
		A21 = A(EN,NA) / AN	EISP7159
		A22 = A(EN,EN) / AN	EISP7160
		B11 = B(NA,NA) / BN	EISP7161
		B12 = B(NA,EN) / BN	EISP7162
		B22 = B(EN,EN) / BN	EISP7163
		E = A11 / B11	EISP7164
		EI = A22 / B22	EISP7165
		S = A21 / (B11 * B22)	EISP7166
		T = (A22 - E * B22) / B22	EISP7167
		IF (ABS(E) .LE. ABS(EI)) GO TO 431	EISP7168
		E = EI	EISP7169
		T = (A11 - E * B11) / B11	EISP7170
	431	C = 0.5E0 * (T - S * B12)	EISP7171
-		D = C * C + S * (A12 - E * B12)	EISP7172
		IF (D .LT. 0.0E0) GO TO 480	EISP7173
C	 TWO REAL ROOTS.	EISP7174
-C		ZERO BOTH A(EN,NA) AND B(EN,NA)	EISP7175
		E = E + (C + SIGN(SQRT(D),C))	EISP7176
		A11 = A11 - E * B11	EISP7177
		A12 = A12 - E * B12	EISP7178
		A22 = A22 - E * B22	EISP7179
		IF (ABS(A11) + ABS(A12) .LT.	EISP7180
	X	ABS(A21) + ABS(A22)) GO TO 432	EISP7181
		A1 = A12	EISP7182
		A2 = A11	EISP7183
		GO TO 435	EISP7184
-	432	A1 = A22	EISP7185
		A2 = A21	EISP7186
C	 CHOOSE AND APPLY REAL Z	EISP7187
-	435	S = ABS(A1) + ABS(A2)	EISP7188
		U1 = A1 / S	EISP7189
		U2 = A2 / S	EISP7190
		R = SIGN(SQRT(U1*U1+U2*U2),U1)	EISP7191
		V1 = -(U1 + R) / R	EISP7192
		V2 = -U2 / R	EISP7193
		U2 = V2 / V1	EISP7194
-C			EISP7195
		DO 440 I = 1, EN	EISP7196
		T = A(I,EN) + U2 * A(I,NA)	EISP7197
		A(I,EN) = A(I,EN) + T * V1	EISP7198
		A(I,NA) = A(I,NA) + T * V2	EISP7199
		T = B(I,EN) + U2 * B(I,NA)	EISP7200
		B(I,EN) = B(I,EN) + T * V1	EISP7201
		B(I,NA) = B(I,NA) + T * V2	EISP7202
	440	CONTINUE	EISP7203
C			EISP7204

-	IF (.NOT. MATZ) GO TO 450	EISP7205
C		EISP7206
-	DO 445 I = 1, N	EISP7207
-	T = Z(I,EN) + U2 * Z(I,NA)	EISP7208
-	Z(I,EN) = Z(I,EN) + T * V1	EISP7209
-	Z(I,NA) = Z(I,NA) + T * V2	EISP7210
-	445 CONTINUE	EISP7211
C		EISP7212
-	450 IF (BN .EQ. 0.0E0) GO TO 475	EISP7213
-	IF (AN .LT. ABS(E) * BN) GO TO 455	EISP7214
-	A1 = B(NA,NA)	EISP7215
-	A2 = B(EN,NA)	EISP7216
-	GO TO 460	EISP7217
-	455 A1 = A(NA,NA)	EISP7218
-	A2 = A(EN,NA)	EISP7219
C CHOOSE AND APPLY REAL Q	EISP7220
-	460 S = ABS(A1) + ABS(A2)	EISP7221
-	IF (S .EQ. 0.0E0) GO TO 475	EISP7222
-	U1 = A1 / S	EISP7223
-	U2 = A2 / S	EISP7224
-	R = SIGN(SQRT(U1*U1+U2*U2),U1)	EISP7225
-	V1 = -(U1 + R) / R	EISP7226
-	V2 = -U2 / R	EISP7227
-	U2 = V2 / V1	EISP7228
C		EISP7229
-	DO 470 J = NA, N	EISP7230
-	T = A(NA,J) + U2 * A(EN,J)	EISP7231
-	A(NA,J) = A(NA,J) + T * V1	EISP7232
-	A(EN,J) = A(EN,J) + T * V2	EISP7233
-	T = B(NA,J) + U2 * B(EN,J)	EISP7234
-	B(NA,J) = B(NA,J) + T * V1	EISP7235
-	B(EN,J) = B(EN,J) + T * V2	EISP7236
-	470 CONTINUE	EISP7237
C		EISP7238
-	475 A(EN,NA) = 0.0E0	EISP7239
-	B(EN,NA) = 0.0E0	EISP7240
-	ALFR(NA) = A(NA,NA)	EISP7241
-	ALFR(EN) = A(EN,EN)	EISP7242
-	IF (B(NA,NA) .LT. 0.0E0) ALFR(NA) = -ALFR(NA)	EISP7243
-	IF (B(EN,EN) .LT. 0.0E0) ALFR(EN) = -ALFR(EN)	EISP7244
-	BETA(NA) = ABS(B(NA,NA))	EISP7245
-	BETA(EN) = ABS(B(EN,EN))	EISP7246
-	ALFI(EN) = 0.0E0	EISP7247
-	ALFI(NA) = 0.0E0	EISP7248
-	GO TO 505	EISP7249
C TWO COMPLEX ROOTS	EISP7250
-	480 E = E + C	EISP7251
-	EI = SQRT(-D)	EISP7252
-	A11R = A11 - E * B11	EISP7253
-	A11I = EI * B11	EISP7254
-	A12R = A12 - E * B12	EISP7255
-	A12I = EI * B12	EISP7256
-	A22R = A22 - E * B22	EISP7257
-	A22I = EI * B22	EISP7258
-	IF (ABS(A11R) + ABS(A11I) + ABS(A12R) + ABS(A12I) .LT.	EISP7259
X	ABS(A21) + ABS(A22R) + ABS(A22I)) GO TO 482	EISP7260
-	A1 = A12R	EISP7261
-	A1I = A12I	EISP7262
-	A2 = -A11R	EISP7263
-	A2I = -A11I	EISP7264

-	GO TO 485	EISP7265
482	A1 = A22R	EISP7266
	A1I = A22I	EISP7267
-	A2 = -A21	EISP7268
	A2I = 0.0E0	EISP7269
C CHOOSE COMPLEX Z	EISP7270
-	485 CZ = SQRT(A1*A1+A1I*A1I)	EISP7271
	IF (CZ .EQ. 0.0E0) GO TO 487	EISP7272
	SZR = (A1 * A2 + A1I * A2I) / CZ	EISP7273
-	SZI = (A1 * A2I - A1I * A2) / CZ	EISP7274
	R = SQRT(CZ*CZ+SZR*SZR+SZI*SZI)	EISP7275
-	CZ = CZ / R	EISP7276
	SZR = SZR / R	EISP7277
-	SZI = SZI / R	EISP7278
	GO TO 490	EISP7279
487	SZR = 1.0E0	EISP7280
	SZI = 0.0E0	EISP7281
-	490 IF (AN .LT. (ABS(E) + EI) * BN) GO TO 492	EISP7282
	A1 = CZ * B11 + SZR * B12	EISP7283
	A1I = SZI * B12	EISP7284
-	A2 = SZR * B22	EISP7285
	A2I = SZI * B22	EISP7286
	GO TO 495	EISP7287
-	492 A1 = CZ * A11 + SZR * A12	EISP7288
	A1I = SZI * A12	EISP7289
	A2 = CZ * A21 + SZR * A22	EISP7290
	A2I = SZI * A22	EISP7291
-C CHOOSE COMPLEX Q	EISP7292
-	495 CQ = SQRT(A1*A1+A1I*A1I)	EISP7293
	IF (CQ .EQ. 0.0E0) GO TO 497	EISP7294
-	SQR = (A1 * A2 + A1I * A2I) / CQ	EISP7295
	SQI = (A1 * A2I - A1I * A2) / CQ	EISP7296
-	R = SQRT(CQ*CQ+SQR*SQR+SQI*SQI)	EISP7297
	CQ = CQ / R	EISP7298
-	SQR = SQR / R	EISP7299
	SQI = SQI / R	EISP7300
	GO TO 500	EISP7301
-	497 SQR = 1.0E0	EISP7302
	SQI = 0.0E0	EISP7303
C COMPUTE DIAGONAL ELEMENTS THAT WOULD RESULT	EISP7304
-C	IF TRANSFORMATIONS WERE APPLIED	EISP7305
-	500 SSR = SQR * SZR + SQI * SZI	EISP7306
	SSI = SQR * SZI - SQI * SZR	EISP7307
	I = 1	EISP7308
-	TR = CQ * CZ * A11 + CQ * SZR * A12 + SQR * CZ * A21	EISP7309
X	+ SSR * A22	EISP7310
-	TI = CQ * SZI * A12 - SQI * CZ * A21 + SSI * A22	EISP7311
-	DR = CQ * CZ * B11 + CQ * SZR * B12 + SSR * B22	EISP7312
	DI = CQ * SZI * B12 + SSI * B22	EISP7313
	GO TO 503	EISP7314
-	502 I = 2	EISP7315
	TR = SSR * A11 - SQR * CZ * A12 - CQ * SZR * A21	EISP7316
X	+ CQ * CZ * A22	EISP7317
-	TI = -SSI * A11 - SQI * CZ * A12 + CQ * SZI * A21	EISP7318
-	DR = SSR * B11 - SQR * CZ * B12 + CQ * CZ * B22	EISP7319
	DI = -SSI * B11 - SQI * CZ * B12	EISP7320
-	503 T = TI * DR - TR * DI	EISP7321
	J = NA	EISP7322
	IF (T .LT. 0.0E0) J = EN	EISP7323
	R = SQRT(DR*DR+DI*DI)	EISP7324

BETA(J) = BN * R		EISP7325
ALFR(J) = AN * (TR * DR + TI * DI) / R		EISP7326
ALFI(J) = AN * T / R		EISP7327
IF (I .EQ. 1) GO TO 502		EISP7328
505 ISW = 3 - ISW		EISP7329
510 CONTINUE		EISP7330
B(N,1) = EPSB		EISP7331
C		EISP7332
RETURN		EISP7333
C** THIS PROGRAM VALID ON FTN4 AND FTN5 **		EISP7334
END		EISP7335
C ROUTINE NAME - PF263=QZVEC		QZVEC 2
C FROM EISPACK		QZVEC 3
C		QZVEC 4
C		QZVEC 5
C-----		QZVEC 6
C		QZVEC 7
C LATEST REVISION - AUGUST 1,1984		QZVEC 8
C COMPUTER SCIENCES CORP., HAMPTON, VA.		QZVEC 9
C		QZVEC 10
C		QZVEC 11
C PURPOSE - THIS SUBROUTINE ACCEPTS A PAIR OF REAL		QZVEC 12
C MATRICES, ONE OF THEM IN QUASI-TRIANGULAR		QZVEC 13
C FORM (IN WHICH EACH 2-BY-2 BLOCK CORRESPONDS		QZVEC 14
C TO A PAIR OF COMPLEX EIGENVALUES) AND THE		QZVEC 15
C OTHER IN UPPER TRIANGULAR FORM. IT COMPUTES		QZVEC 16
C THE EIGENVECTORS OF THE TRIANGULAR PROBLEM		QZVEC 17
C AND TRANSFORMS THE RESULTS BACK TO THE		QZVEC 18
C ORIGINAL COORDINATE SYSTEM. IT IS USUALLY		QZVEC 19
C PRECEDED BY QZHES(PF260), QZIT(PF261), AND		QZVEC 20
C QZVAL(PF262).		QZVEC 21
C		QZVEC 22
C USAGE - CALL QZVEC(NM,N,A,B,ALFR,ALFI,BETA,Z)		QZVEC 23
C		QZVEC 24
C ARGUMENTS NM - ON INPUT NM MUST BE SET TO THE ROW DIMENSION		QZVEC 25
C OF TWO-DIMENSIONAL ARRAY PARAMETERS AS		QZVEC 26
C DECLARED IN THE CALLING PROGRAM DIMENSION		QZVEC 27
C STATEMENT.		QZVEC 28
C		QZVEC 29
C N - ON INPUT N IS THE ORDER OF THE MATRICES.		QZVEC 30
C		QZVEC 31
C A - ON INPUT A CONTAINS A REAL UPPER QUASI-		QZVEC 32
C TRIANGULAR MATRIX.		QZVEC 33
C MUST BE OF DIMENSION NM X N.		QZVEC 34
C		QZVEC 35
C - ON OUTPUT A IS UNALTERED. ITS SUBDIAGONAL		QZVEC 36
C ELEMENTS PROVIDE INFORMATION ABOUT THE STORAGE		QZVEC 37
C OF THE COMPLEX EIGENVECTORS.		QZVEC 38
C		QZVEC 39
C B - ON INPUT B CONTAINS A REAL UPPER TRIANGULAR		QZVEC 40
C MATRIX. IN ADDITION, LOCATION B(N,1) CONTAINS		QZVEC 41
C THE TOLERANCE QUANTITY (EPSB) COMPUTED AND		QZVEC 42
C SAVED IN QZIT(PF261).		QZVEC 43
C MUST BE OF DIMENSION NM X N.		QZVEC 44
C		QZVEC 45
C ON OUTPUT B HAS BEEN DESTROYED.		QZVEC 46
C		QZVEC 47
C ALFR - ON INPUT ALFR IS A VECTOR SUCH THAT THE		QZVEC 48
C RATIOS ((ALFR+I*ALFI)/BETA) ARE THE		QZVEC 49
C GENERALIZED EIGENVALUES. THEY ARE USUALLY		QZVEC 50

OBTAINED FROM QZVAL(PF262).
MUST BE OF DIMENSION N.

ALFI - ON INPUT ALFI IS A VECTOR SUCH THAT THE RATIOS
((ALFR+I*ALFI)/BETA) ARE THE GENERALIZED
EIGENVALUES. THEY ARE USUALLY OBTAINED FROM
QZVAL(PF262).
MUST BE OF DIMENSION N.

BETA - ON INPUT BETA IS A VECTOR SUCH THAT THE RATIOS
((ALFR+I*ALFI)/BETA) ARE THE GENERALIZED
EIGENVALUES. THEY ARE USUALLY OBTAINED FROM
QZVAL(PF262).
MUST BE OF DIMENSION N.

Z - ON INPUT Z CONTAINS THE TRANSFORMATION MATRIX
PRODUCED IN THE REDUCTIONS BY QZHES(PF260),
QZIT(PF261), AND QZVAL(PF262), IF PERFORMED.
IF THE EIGENVECTORS OF THE TRIANGULAR PROBLEM
ARE DESIRED, Z MUST CONTAIN THE IDENTITY
MATRIX.
MUST BE OF DIMENSION NM X N.

ON OUTPUT Z CONTAINS THE REAL AND IMAGINARY
PARTS OF THE EIGENVECTORS. IF ALFI(I) .EQ.
0.0, THE I-TH EIGENVALUE IS REAL AND THE I-TH
COLUMN OF Z CONTAINS ITS EIGENVECTOR. IF
ALFI(I) .NE. 0.0, THE I-TH EIGENVALUE IS
COMPLEX. IF ALFI(I) .GT. 0.0, THE EIGENVALUE
IS THE FIRST OF A COMPLEX PAIR AND THE I-TH
AND (I+1)-TH COLUMNS OF Z CONTAIN ITS EIGEN-
VECTOR. IF ALFI(I) .LT. 0.0, THE EIGEN-
VALUE IS THE SECOND OF A COMPLEX PAIR AND THE
(I-1)-TH AND I-TH COLUMNS OF Z CONTAIN THE
CONJUGATE OF ITS EIGENVECTOR. EACH EIGEN-
VECTOR IS NORMALIZED SO THAT THE MODULUS
OF ITS LARGEST COMPONENT IS 1.0 .

REQUIRED ROUTINES - NONE

REMARKS 1. THIS SUBROUTINE IS THE OPTIONAL FOURTH STEP
OF THE QZ ALGORITHM FOR SOLVING GENERALIZED
MATRIX EIGENVALUE PROBLEMS, SIAM J. NUMER.
ANAL. 10, 241-256(1973) BY MOLER AND STEWART.

EXAMPLE :

PROGRAM TQZVEC(OUTPUT,TAPE6=OUTPUT)
DIMENSION A(5,5),B(5,5),ALFR(5),ALFI(5),BETA(5),Z(5,5)
LOGICAL MATZ

N = 5
NM = 5
MATZ = .TRUE.
EPS1 = 0.0E0

DATA A /10.,2.,3.,2*1.,2.,12.,1.,2.,1.,3.,1.,11.,
* 1.,-1.,1.,2.,1.,9.,3*1.,-1.,1.,15. /

DATA B /12.,1.,-1.,2.,2*1.,14.,1.,-1.,1.,-1.,1.,

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C      *      16.,-1.,1.,2.,-1.,-1.,12.,-1.,3*1.,-1.,11.  /
C
C      CALL QZHES(NM,N,A,B,MATZ,Z)
C      CALL QZIT(NM,N,A,B,EPS1,MATZ,Z,IERR)
C      CALL QZVAL(NM,N,A,B,ALFR,ALFI,BETA,MATZ,Z)
C      CALL QZVEC(NM,N,A,B,ALFR,ALFI,BETA,Z)
C      WRITE(6,99) IERR
C      WRITE(6,100) ((Z(I,J),I=1,5),J=1,5)
C99    FORMAT(1H1,7HIERR = ,I4)
C100   FORMAT(5H Z = /5(1H ,5(G8.2,2X)/))
C      STOP
C      END
C
C      OUTPUT :
C
C      IERR =      0
C      Z =
C      .26      -.59E-01      .23      -.30      -1.0
C      -.85      .39      1.0      -.69      .26
C      1.0      1.0      .85      .88      .54E-01
C      -1.0      .83      -.39      .72      -.46
C      -.45      -.84      .65      1.0      -.19E-01
C
C-----
C      SUBROUTINE QZVEC(NM,N,A,B,ALFR,ALFI,BETA,Z)
C
C      implicit real*8 (a-h,o-z)
C      INTEGER I,J,K,M,N,EN,II,JJ,NA,NM,NN,ISW,ENM2
C      REAL*8 A(NM,N),B(NM,N),ALFR(N),ALFI(N),BETA(N),Z(NM,N)
C      REAL*8 D,Q,R,S,T,W,X,Y,DI,DR,RA,RR,SA,TI,TR,T1,T2,W1,X1,
X      ZZ,Z1,ALFM,ALMI,ALMR,BETM,EPSB
C      EPSB = B(N,1)
C      ISW = 1
C      ..... FOR EN=N STEP -1 UNTIL 1 DO -- .....
C      DO 800 NN = 1, N
C          EN = N + 1 - NN
C          NA = EN - 1
C          IF (ISW .EQ. 2) GO TO 795
C          IF (ALFI(EN) .NE. 0.0E0) GO TO 710
C      ..... REAL VECTOR .....
C          M = EN
C          B(EN,EN) = 1.0E0
C          IF (NA .EQ. 0) GO TO 800
C          ALFM = ALFR(M)
C          BETM = BETA(M)
C      ..... FOR I=EN-1 STEP -1 UNTIL 1 DO -- .....
C          DO 700 II = 1, NA
C              I = EN - II
C              W = BETM * A(I,I) - ALFM * B(I,I)
C              R = 0.0E0
C
C          DO 610 J = M, EN
C      610      R = R + (BETM * A(I,J) - ALFM * B(I,J)) * B(J,EN)
C
C          IF (I .EQ. 1 .OR. ISW .EQ. 2) GO TO 630
C          IF (BETM * A(I,I-1) .EQ. 0.0E0) GO TO 630
C          ZZ = W
C          S = R
C          GO TO 690
C      630      M = I

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QZVEC111
 QZVEC112
 QZVEC113
 QZVEC114
 QZVEC115
 QZVEC116
 QZVEC117
 QZVEC118
 QZVEC119
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 QZVEC121
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 EISP7367
 EISP7368
 EISP7369
 EISP7370

	IF (ISW .EQ. 2) GO TO 640	EISP7371
C REAL 1-BY-1 BLOCK	EISP7372
	T = W	EISP7373
	IF (W .EQ. 0.0E0) T = EPSB	EISP7374
	B(I,EN) = -R / T	EISP7375
	GO TO 700	EISP7376
C REAL 2-BY-2 BLOCK	EISP7377
640	X = BETM * A(I,I+1) - ALFM * B(I,I+1)	EISP7378
	Y = BETM * A(I+1,I)	EISP7379
	Q = W * ZZ - X * Y	EISP7380
	T = (X * S - ZZ * R) / Q	EISP7381
	B(I,EN) = T	EISP7382
	IF (ABS(X) .LE. ABS(ZZ)) GO TO 650	EISP7383
	B(I+1,EN) = (-R - W * T) / X	EISP7384
	GO TO 690	EISP7385
650	B(I+1,EN) = (-S - Y * T) / ZZ	EISP7386
690	ISW = 3 - ISW	EISP7387
700	CONTINUE	EISP7388
C END REAL VECTOR	EISP7389
	GO TO 800	EISP7390
C COMPLEX VECTOR	EISP7391
710	M = NA	EISP7392
	ALMR = ALFR(M)	EISP7393
	ALMI = ALFI(M)	EISP7394
	BETM = BETA(M)	EISP7395
C LAST VECTOR COMPONENT CHOSEN IMAGINARY SO THAT	EISP7396
C	EIGENVECTOR MATRIX IS TRIANGULAR	EISP7397
	Y = BETM * A(EN,NA)	EISP7398
	B(NA,NA) = -ALMI * B(EN,EN) / Y	EISP7399
	B(NA,EN) = (ALMR * B(EN,EN) - BETM * A(EN,EN)) / Y	EISP7400
	B(EN,NA) = 0.0E0	EISP7401
	B(EN,EN) = 1.0E0	EISP7402
	ENM2 = NA - 1	EISP7403
	IF (ENM2 .EQ. 0) GO TO 795	EISP7404
C FOR I=EN-2 STEP -1 UNTIL 1 DO --	EISP7405
	DO 790 II = 1, ENM2	EISP7406
	I = NA - II	EISP7407
	W = BETM * A(I,I) - ALMR * B(I,I)	EISP7408
	W1 = -ALMI * B(I,I)	EISP7409
	RA = 0.0E0	EISP7410
	SA = 0.0E0	EISP7411
C		EISP7412
	DO 760 J = M, EN	EISP7413
	X = BETM * A(I,J) - ALMR * B(I,J)	EISP7414
	X1 = -ALMI * B(I,J)	EISP7415
	RA = RA + X * B(J,NA) - X1 * B(J,EN)	EISP7416
	SA = SA + X * B(J,EN) + X1 * B(J,NA)	EISP7417
760	CONTINUE	EISP7418
C		EISP7419
	IF (I .EQ. 1 .OR. ISW .EQ. 2) GO TO 770	EISP7420
	IF (BETM * A(I,I-1) .EQ. 0.0E0) GO TO 770	EISP7421
	ZZ = W	EISP7422
	Z1 = W1	EISP7423
	R = RA	EISP7424
	S = SA	EISP7425
	ISW = 2	EISP7426
	GO TO 790	EISP7427
770	M = I	EISP7428
	IF (ISW .EQ. 2) GO TO 780	EISP7429
C COMPLEX 1-BY-1 BLOCK	EISP7430

-	TR = -RA	EISP7431
	TI = -SA	EISP7432
773	DR = W	EISP7433
-	DI = W1	EISP7434
C COMPLEX DIVIDE (T1,T2) = (TR,TI) / (DR,DI)	EISP7435
775	IF (ABS(DI) .GT. ABS(DR)) GO TO 777	EISP7436
-	RR = DI / DR	EISP7437
	D = DR + DI * RR	EISP7438
	T1 = (TR + TI * RR) / D	EISP7439
	T2 = (TI - TR * RR) / D	EISP7440
-	GO TO (787,782), ISW	EISP7441
	CALL GOTOER	EISP7442
777	RR = DR / DI	EISP7443
-	D = DR * RR + DI	EISP7444
	T1 = (TR * RR + TI) / D	EISP7445
	T2 = (TI * RR - TR) / D	EISP7446
-	GO TO (787,782), ISW	EISP7447
	CALL GOTOER	EISP7448
C COMPLEX 2-BY-2 BLOCK	EISP7449
780	X = BETM * A(I,I+1) - ALMR * B(I,I+1)	EISP7450
-	X1 = -ALMI * B(I,I+1)	EISP7451
	Y = BETM * A(I+1,I)	EISP7452
	TR = Y * RA - W * R + W1 * S	EISP7453
-	TI = Y * SA - W * S - W1 * R	EISP7454
	DR = W * ZZ - W1 * Z1 - X * Y	EISP7455
	DI = W * Z1 + W1 * ZZ - X1 * Y	EISP7456
-	IF (DR .EQ. 0.0E0 .AND. DI .EQ. 0.0E0) DR = EPSB	EISP7457
	GO TO 775	EISP7458
782	B(I+1,NA) = T1	EISP7459
	B(I+1,EN) = T2	EISP7460
-	ISW = 1	EISP7461
	IF (ABS(Y) .GT. ABS(W) + ABS(W1)) GO TO 785	EISP7462
-	TR = -RA - X * B(I+1,NA) + X1 * B(I+1,EN)	EISP7463
	TI = -SA - X * B(I+1,EN) - X1 * B(I+1,NA)	EISP7464
-	GO TO 773	EISP7465
785	T1 = (-R - ZZ * B(I+1,NA) + Z1 * B(I+1,EN)) / Y	EISP7466
	T2 = (-S - ZZ * B(I+1,EN) - Z1 * B(I+1,NA)) / Y	EISP7467
-	787 B(I,NA) = T1	EISP7468
	B(I,EN) = T2	EISP7469
790	CONTINUE	EISP7470
-C END COMPLEX VECTOR	EISP7471
795	ISW = 3 - ISW	EISP7472
800	CONTINUE	EISP7473
C END BACK SUBSTITUTION.	EISP7474
-C	TRANSFORM TO ORIGINAL COORDINATE SYSTEM.	EISP7475
C	FOR J=N STEP -1 UNTIL 1 DO --	EISP7476
	DO 880 JJ = 1, N	EISP7477
-	J = N + 1 - JJ	EISP7478
C		EISP7479
	DO 880 I = 1, N	EISP7480
-	ZZ = 0.0E0	EISP7481
C		EISP7482
	DO 860 K = 1, J	EISP7483
860	ZZ = ZZ + Z(I,K) * B(K,J)	EISP7484
-C		EISP7485
	Z(I,J) = ZZ	EISP7486
880	CONTINUE	EISP7487
-C NORMALIZE SO THAT MODULUS OF LARGEST	EISP7488
C	COMPONENT OF EACH VECTOR IS 1.	EISP7489
C	(ISW IS 1 INITIALLY FROM BEFORE)	EISP7490

DO 950 J = 1, N		EISP7491
D = 0.0E0		EISP7492
IF (ISW .EQ. 2) GO TO 920		EISP7493
IF (ALFI(J) .NE. 0.0E0) GO TO 945		EISP7494
		EISP7495
DO 890 I = 1, N		EISP7496
IF (ABS(Z(I,J)) .GT. D) D = ABS(Z(I,J))		EISP7497
CONTINUE		EISP7498
		EISP7499
DO 900 I = 1, N		EISP7500
Z(I,J) = Z(I,J) / D		EISP7501
		EISP7502
GO TO 950		EISP7503
		EISP7504
DO 920 I = 1, N		EISP7505
R = ABS(Z(I,J-1)) + ABS(Z(I,J))		EISP7506
IF (R .NE. 0.0E0) R = R * SQRT((Z(I,J-1)/R)**2		EISP7507
	+(Z(I,J)/R)**2)	EISP7508
IF (R .GT. D) D = R		EISP7509
CONTINUE		EISP7510
		EISP7511
DO 940 I = 1, N		EISP7512
Z(I,J-1) = Z(I,J-1) / D		EISP7513
Z(I,J) = Z(I,J) / D		EISP7514
CONTINUE		EISP7515
		EISP7516
945 ISW = 3 - ISW		EISP7517
950 CONTINUE		EISP7518
		EISP7519
RETURN		EISP7520
END		EISP7521
ROUTINE NAME	- PF266=RGG	RGG 2
FROM EISPACK		RGG 3
		RGG 4
		RGG 5
		RGG 6
		RGG 7
LATEST REVISION	- AUGUST 1,1984	RGG 8
	COMPUTER SCIENCES CORP., HAMPTON, VA.	RGG 9
		RGG 10
PURPOSE	- THIS SUBROUTINE CALLS THE RECOMMENDED	RGG 11
	SEQUENCE OF SUBROUTINES FROM THE EIGENSYSTEM	RGG 12
	SUBROUTINE PACKAGE (EISPACK) TO FIND THE	RGG 13
	EIGENVALUES AND EIGENVECTORS (IF DESIRED) FOR	RGG 14
	THE REAL GENERAL GENERALIZED EIGENPROBLEM AX	RGG 15
	= (LAMBDA)BX.	RGG 16
		RGG 17
		RGG 18
USAGE	- CALL RGG(NM,N,A,B,ALFR,ALFI,BETA,MATZ,Z,IERR)	RGG 19
		RGG 20
ARGUMENTS	NM	RGG 21
	- ON INPUT NM MUST BE SET TO THE ROW DIMENSION	RGG 22
	OF THE TWO-DIMENSIONAL ARRAY PARAMETERS AS	RGG 23
	DECLARED IN THE CALLING PROGRAM DIMENSION	RGG 24
	STATEMENT.	RGG 25
		RGG 26
	N	RGG 27
	- ON INPUT N IS THE ORDER OF THE MATRICES A	RGG 28
	AND B.	RGG 29
		RGG 30
	A	RGG 30
	- ON INPUT A CONTAINS A REAL GENERAL MATRIX.	
	MUST BE OF DIMENSION NM X N.	

		RGG	31
B	- ON INPUT B CONTAINS A REAL GENERAL MATRIX. MUST BE OF DIMENSION NM X N.	RGG	32
		RGG	33
		RGG	34
ALFR	- ON OUTPUT ALFR CONTAINS THE REAL PART OF THE NUMERATORS OF THE EIGENVALUES. MUST BE OF DIMENSION N.	RGG	35
		RGG	36
		RGG	37
		RGG	38
ALFI	- ON OUTPUT ALFI CONTAINS THE IMAGINARY PART OF THE NUMERATORS OF THE EIGENVALUES. MUST BE OF DIMENSION N.	RGG	39
		RGG	40
		RGG	41
		RGG	42
BETA	- ON OUTPUT BETA CONTAINS THE DENOMINATORS OF THE EIGENVALUES, WHICH ARE THUS GIVEN BY THE RATIOS (ALFR+I*ALFI)/BETA. COMPLEX CONJUGATE PAIRS OF EIGENVALUES APPEAR CONSECUTIVELY WITH THE EIGENVALUE HAVING THE POSITIVE IMAGINARY PART FIRST. MUST BE OF DIMENSION N.	RGG	43
		RGG	44
		RGG	45
		RGG	46
		RGG	47
		RGG	48
		RGG	49
		RGG	50
MATZ	- ON INPUT MATZ IS AN INTEGER VARIABLE SET EQUAL TO ZERO IF ONLY EIGENVALUES ARE DESIRED. OTHERWISE IT IS SET TO ANY NON-ZERO INTEGER FOR BOTH EIGENVALUES AND EIGENVECTORS.	RGG	51
		RGG	52
		RGG	53
		RGG	54
		RGG	55
		RGG	56
Z	- ON OUTPUT Z CONTAINS THE REAL AND IMAGINARY PARTS OF THE EIGENVECTORS IF MATZ IS NOT ZERO. IF THE J-TH EIGENVALUE IS REAL, THE J-TH COLUMN OF Z CONTAINS ITS EIGENVECTOR. IF THE J-TH EIGENVALUE IS COMPLEX WITH POSITIVE IMAGINARY PART, THE J-TH AND (J+1)-TH COLUMNS OF Z CONTAIN THE REAL AND IMAGINARY PARTS OF ITS EIGENVECTOR. THE CONJUGATE OF THIS VECTOR IS THE EIGENVECTOR FOR THE CONJUGATE EIGENVALUE. MUST BE OF DIMENSION NM X N.	RGG	57
		RGG	58
		RGG	59
		RGG	60
		RGG	61
		RGG	62
		RGG	63
		RGG	64
		RGG	65
		RGG	66
		RGG	67
		RGG	68
		RGG	69
IERR	- ON OUTPUT IERR IS AN INTEGER OUTPUT VARIABLE SET EQUAL TO AN ERROR COMPLETION CODE DESCRIBED IN THE DOCUMENTATION FOR QZIT PF261). THE NORMAL COMPLETION CODE IS ZERO.	RGG	70
		RGG	71
		RGG	72
		RGG	73
		RGG	74
		RGG	75
REQUIRED ROUTINES	- PF260=QZHES, PF261=QZIT, PF262=QZVAL, PF263=QZVEC HC318=EPSLON	RGG	76
		RGG	77
		RGG	78
REMARKS 1.	REFERENCES	RGG	79
		RGG	80
	FROM THE EISPACK PACKAGE OF EIGENSYSTEM ROUTINES.	RGG	81
		RGG	82
2.	SUBROUTINE RGG IS A DRIVER ROUTINE WHICH CALLS ROUTINES QZHES(PF260), QZIT(PF261), QZVAL(PF262), AND QZVEC(PF263).	RGG	83
		RGG	84
		RGG	85
		RGG	86
	QZHES(PF260) ACCEPTS A PAIR OF REAL GENERAL MATRICES AND REDUCES ONE OF THEM TO UPPER HESSENBERG FORM AND THE OTHER TO UPPER TRIANGULAR FORM USING ORTHOGONAL TRANSFORMATIONS.	RGG	87
		RGG	88
		RGG	89
		RGG	90

	RGG	91
QZIT(PF261) ACCEPTS A PAIR OF REAL MATRICES, ONE OF	RGG	92
THEM IN UPPER HESSENBERG FORM AND THE OTHER IN UPPER	RGG	93
TRIANGULAR FORM. IT REDUCES THE HESSENBERG MATRIX TO	RGG	94
QUASI-TRIANGULAR FORM USING ORTHOGONAL TRANSFORMATIONS	RGG	95
WHILE MAINTAINING THE TRIANGULAR FORM OF THE OTHER	RGG	96
MATRIX.	RGG	97
	RGG	98
QZVAL(PF262) ACCEPTS A PAIR OF REAL MATRICES, ONE OF	RGG	99
THEM IN QUASI-TRIANGULAR FORM AND THE OTHER IN UPPER	RGG	100
TRIANGULAR FORM. IT REDUCES THE QUASI-TRIANGULAR	RGG	101
MATRIX FURTHER, SO THAT ANY REMAINING 2-BY-2 BLOCKS	RGG	102
CORRESPOND TO PAIRS OF COMPLEX EIGENVALUES, AND RETURNS	RGG	103
QUANTITIES WHOSE RATIOS GIVE THE GENERALIZED	RGG	104
EIGENVALUES.	RGG	105
	RGG	106
QZVEC(PF263) ACCEPTS A PAIR OF REAL MATRICES, ONE OF	RGG	107
THEM IN QUASI-TRIANGULAR FORM (IN WHICH EACH 2-BY-2	RGG	108
BLOCK CORRESPONDS TO A PAIR OF COMPLEX EIGENVALUES) AND	RGG	109
THE OTHER IN UPPER TRIANGULAR FORM. IT COMPUTES THE	RGG	110
EIGENVECTORS OF THE TRIANGULAR PROBLEM AND TRANSFORMS	RGG	111
THE RESULTS BACK TO THE ORIGINAL COORDINATE SYSTEM.	RGG	112

EXAMPLE :

PROGRAM TRGG(OUTPUT,TAPE6=OUTPUT)

DIMENSION A(5,5),B(5,5),ALFR(5),ALFI(5),BETA(5),Z(5,5)

N = 5

NM = 5

MATZ = 1

DATA A /10.,2.,3.,2*1.,2.,12.,1.,2.,1.,3.,1.,11.,
* 1.,-1.,1.,2.,1.,9.,3*1.,-1.,1.,15. /

DATA B /12.,1.,-1.,2.,2*1.,14.,1.,-1.,1.,-1.,1.,
* 16.,-1.,1.,2.,-1.,-1.,12.,-1.,3*1.,-1.,11. /

CALL RGG(NM,N,A,B,ALFR,ALFI,BETA,MATZ,Z,IERR)

WRITE(6,99) IERR

WRITE(6,100) ALFR,ALFI,BETA,((Z(I,J),I=1,5),J=1,5)

C99 FORMAT(1H1,7HIERR = ,I4)

C100 FORMAT(1H0,7HALFR = /1H ,5(G8.2,2X) /

* 8H0ALFI = /1H ,5(G8.2,2X) /

* 8H0BETA = /1H ,5(G8.2,2X) /

* 5H0Z = /5(1H ,5(G8.2,2X) /)

STOP

END

OUTPUT :

IERR = 0

ALFR =

15. 7.2 16. 10. 8.6

ALFI =

0. 0. 0. 0. 0.

-C	BETA =					RGG	151
C	9.9	17.	14.	11.	13.	RGG	152
C	Z =					RGG	153
-C	.26	-.59E-01	.23	-.30	-1.0	RGG	154
C	-.85	.39	1.0	-.69	.26	RGG	155
C	1.0	1.0	.85	.88	.54E-01	RGG	156
-C	-1.0	.83	-.39	.72	-.46	RGG	157
C	-.45	-.84	.65	1.0	-.19E-01	RGG	158
C						RGG	159
C	-----					RGG	160
-	SUBROUTINE diverg(NM,N,A,B,ALFR,ALFI,BETA,MATZ,Z,IERR)						EISP7
C							
-	implicit real*8 (a-h,o-z)					EISP7613	
-	INTEGER N,NM,IERR,MATZ					EISP7614	
-	REAL*8 A(NM,N),B(NM,N),ALFR(N),ALFI(N),BETA(N),Z(NM,N)					EISP7615	
-	LOGICAL TF					EISP7616	
-	zero = 0.0e+00						
-	IF (N .LE. NM) GO TO 10					EISP7617	
-	IERR = 10 * N					EISP7618	
-	GO TO 50					EISP7619	
-C	10 IF (MATZ .NE. 0) GO TO 20					EISP7620	
C FIND EIGENVALUES ONLY					EISP7621	
-	TF = .FALSE.					EISP7622	
-	CALL QZHES(NM,N,A,B,TF,Z)					EISP7623	
-	CALL QZIT(NM,N,A,B,zero,TF,Z,IERR)					EISP7624	
-	CALL QZVAL(NM,N,A,B,ALFR,ALFI,BETA,TF,Z)					EISP7625	
-	GO TO 50					EISP7626	
C FIND BOTH EIGENVALUES AND EIGENVECTORS					EISP7627	
-	20 TF = .TRUE.					EISP7628	
-	CALL QZHES(NM,N,A,B,TF,Z)					EISP7629	
-	CALL QZIT(NM,N,A,B,zero,TF,Z,IERR)					EISP7630	
-	CALL QZVAL(NM,N,A,B,ALFR,ALFI,BETA,TF,Z)					EISP7631	
-	IF (IERR .NE. 0) GO TO 50					EISP7632	
-	CALL QZVEC(NM,N,A,B,ALFR,ALFI,BETA,Z)					EISP7633	
-	50 RETURN					EISP7634	
-	C** THIS PROGRAM VALID ON FTN4 AND FTN5 **					EISP7635	
-	END					EISP7636	
-	subroutine gotoer						
-	write(6,10)						
-	10 format('there is an error in calculating subroutine')						
-	return						
-	end					EISP7637	
-C							